

Simple Optimization Problems

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Computational Mathematics for Learning and Data Analysis
Master in Computer Science – University of Pisa

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Outline

Optimization Problems

Optimization is difficult

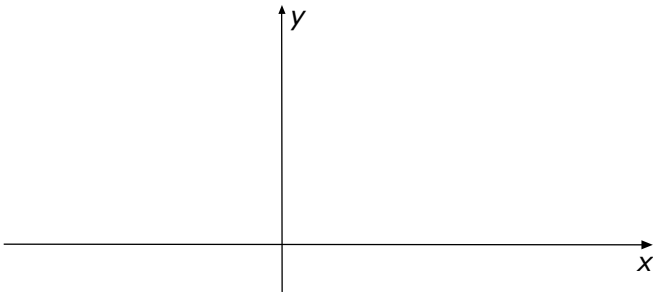
Simple Functions, Univariate case

Simple Functions, Multivariate case

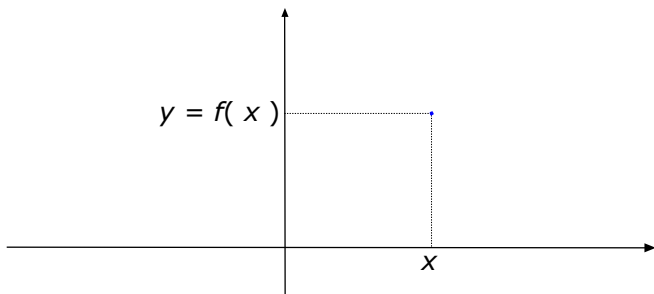
Multivariate Quadratic case: Gradient Method

Wrap up & References

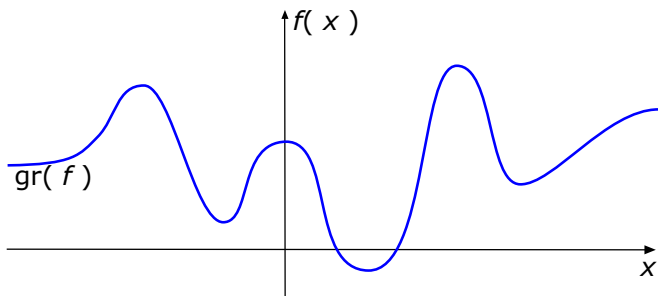
Solutions



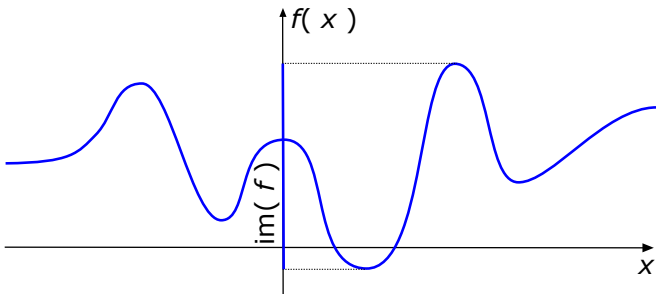
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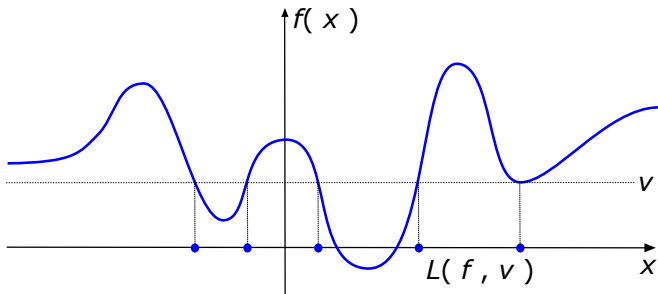
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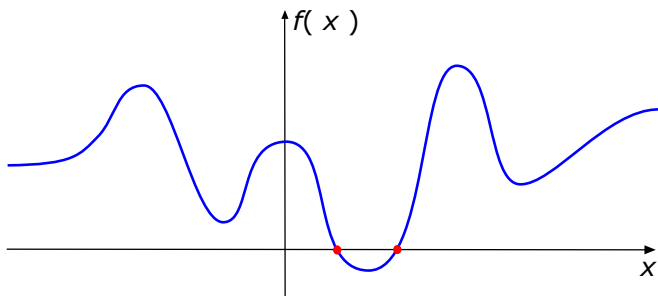
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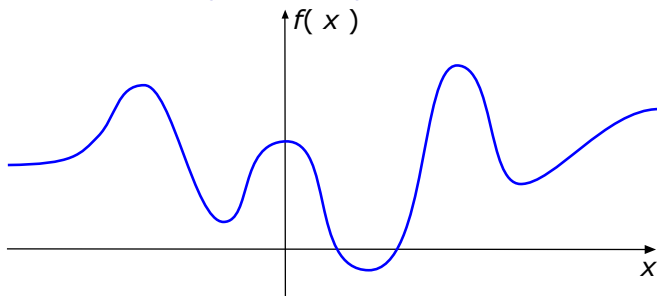
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i.e., projection of $\text{gr}(f)$ on output space (a.k.a. co-domain)



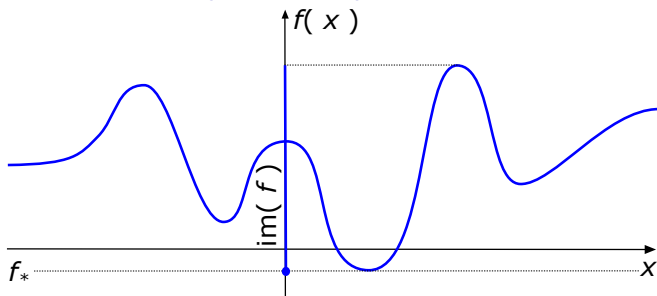
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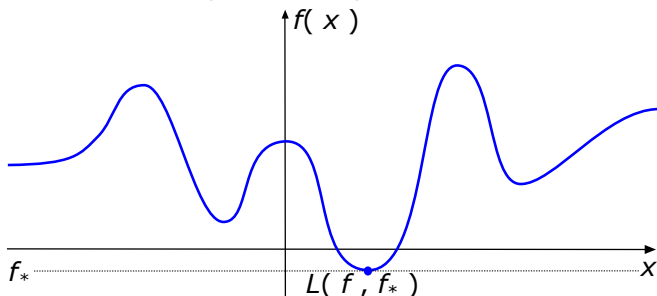
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(**roots of f** = $L(f, 0)$ = level set at value 0)



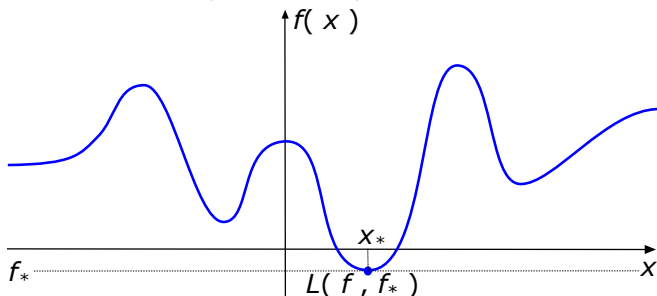
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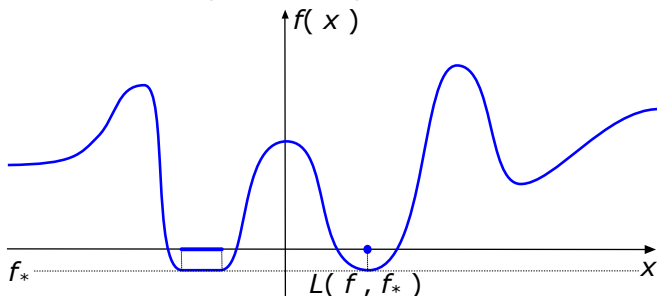
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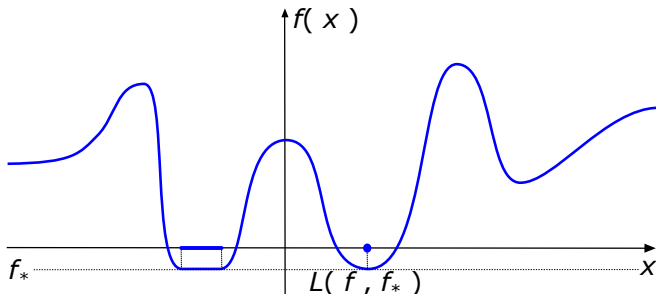
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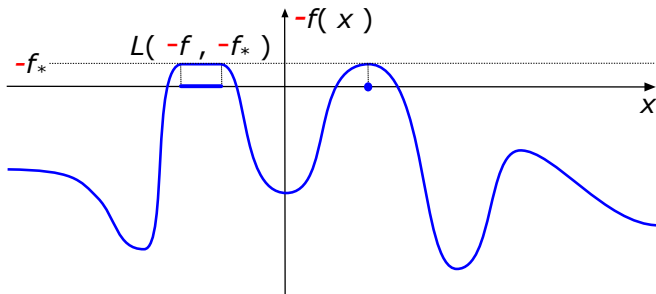
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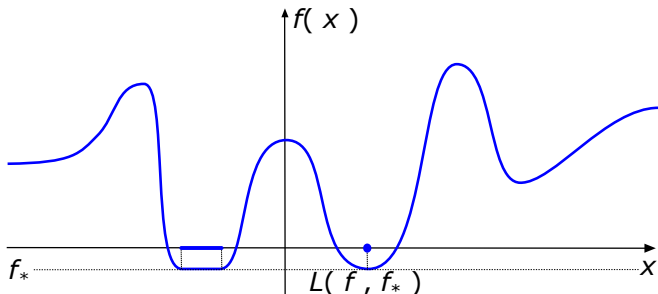
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- ▶ x_* s.t. $f_* = f(x_*) \leq f(x) \quad \forall x \in \mathbb{R}$ optimal solution (if \exists , which it may not)
- ▶ x_* may not be unique: $\exists x' \neq x_* \in L(f, f_*) = X_*$ set of optimal solutions, but we don't care (mostly): all optimal solutions equivalent "in the eyes of f "



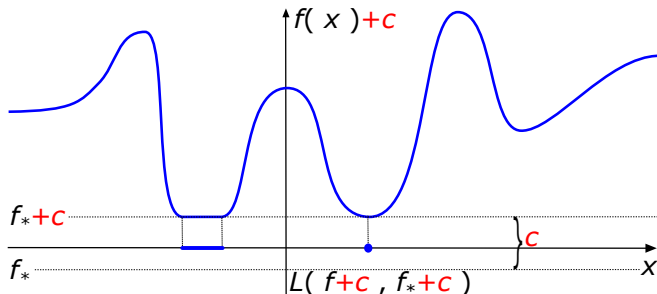
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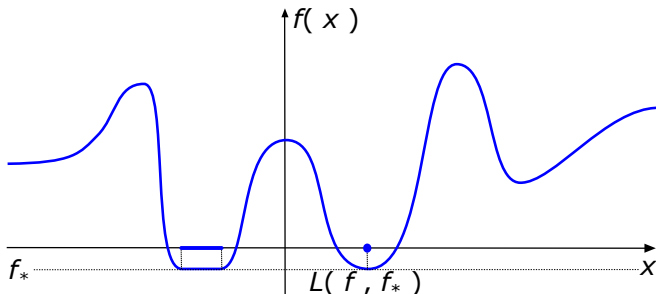
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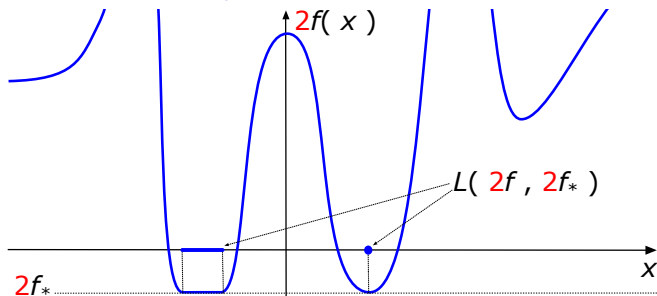
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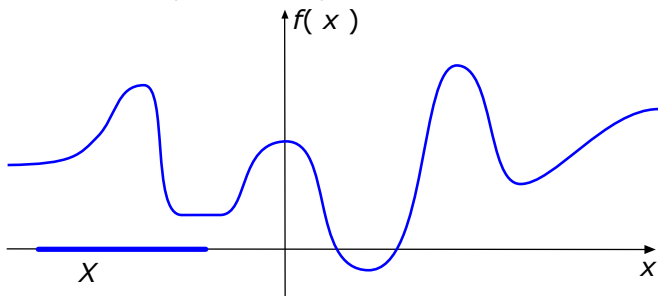
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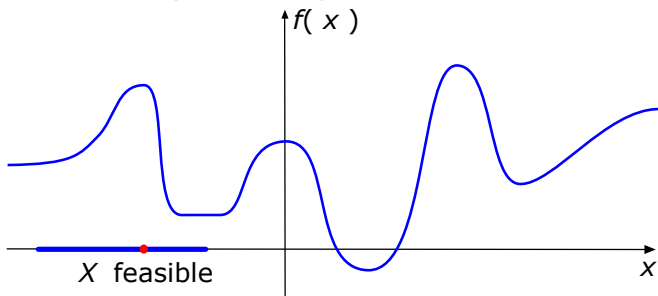
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(Univariate) Constrained optimization problem

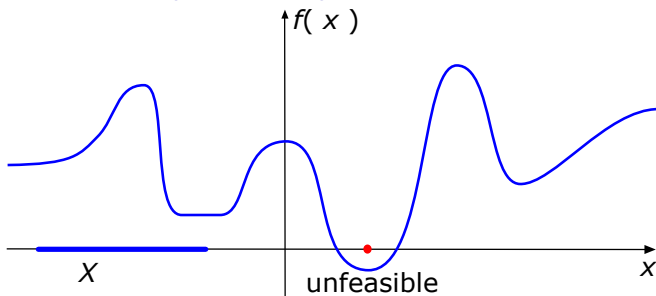
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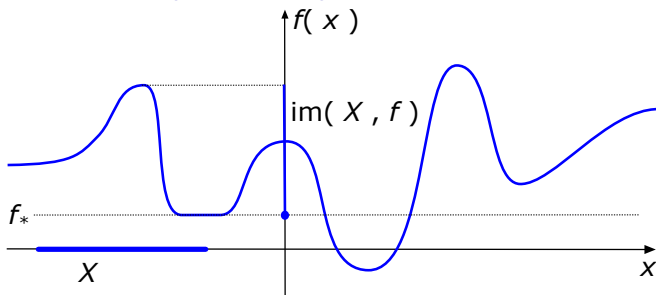
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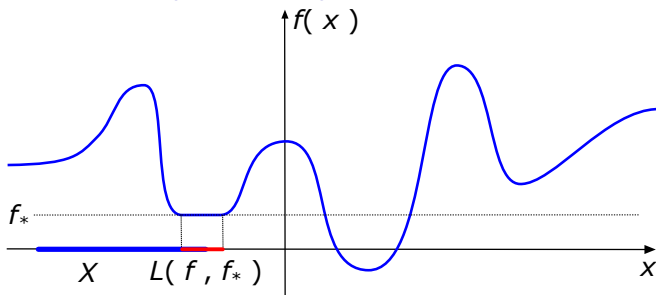
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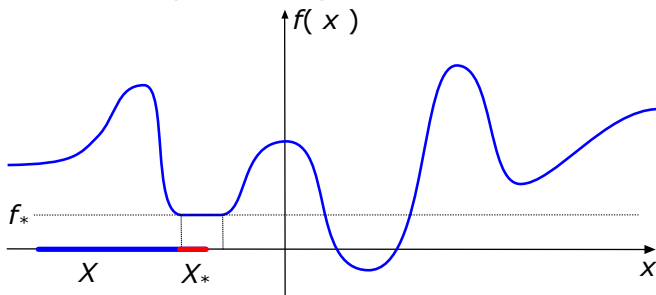
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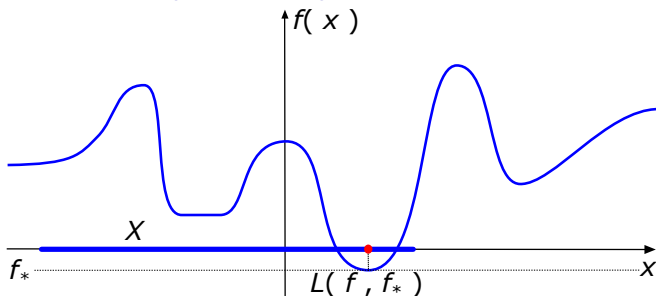
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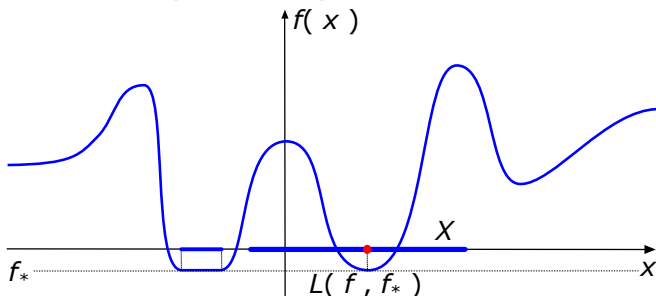
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- ▶ $X_* = L(f, f_*)$



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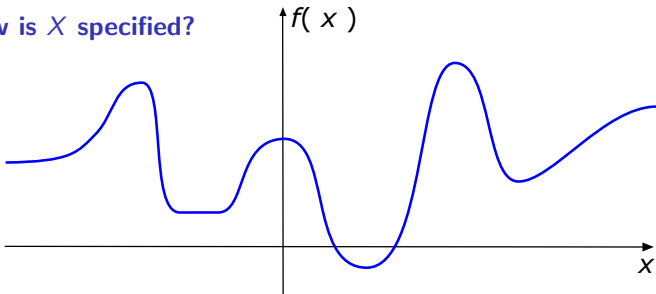
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- ▶ $X_* = L(f, f_*) \cap X$: set of best **feasible solutions**
- ▶ X can be “**useless**” (X_* same) or **partly so** (f_* same) \implies
makes sense to **study the unconstrained case $X = \mathbb{R}$ first**

Anyhow, how is X specified?

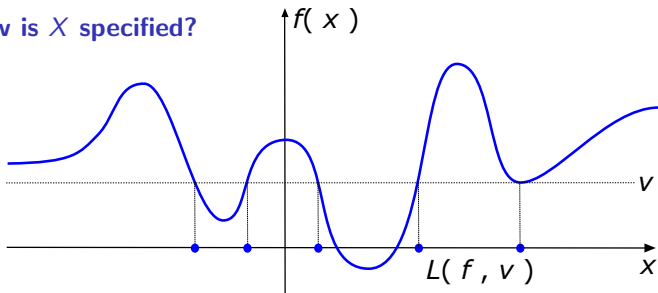
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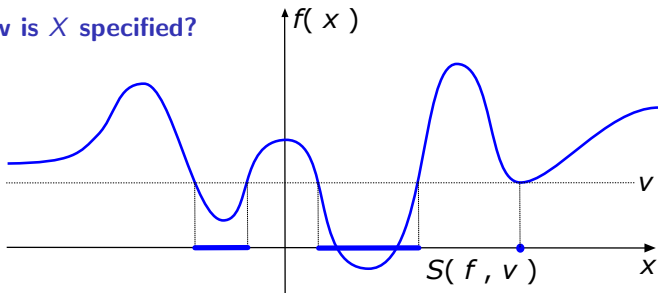
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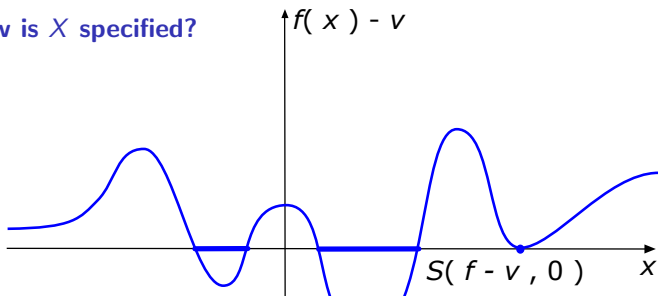
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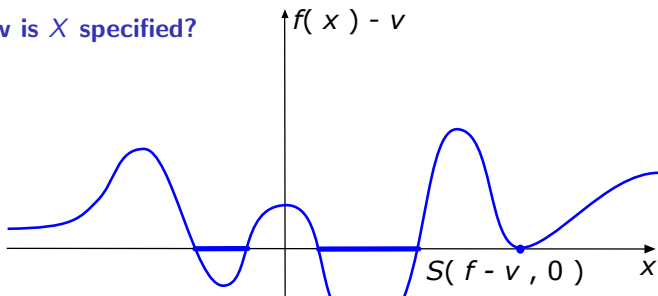
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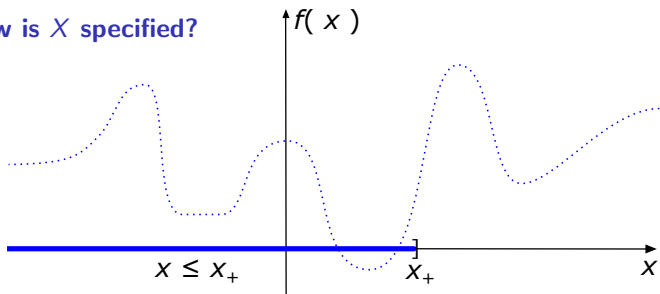
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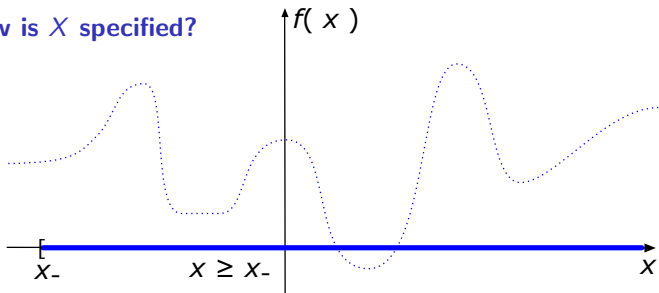
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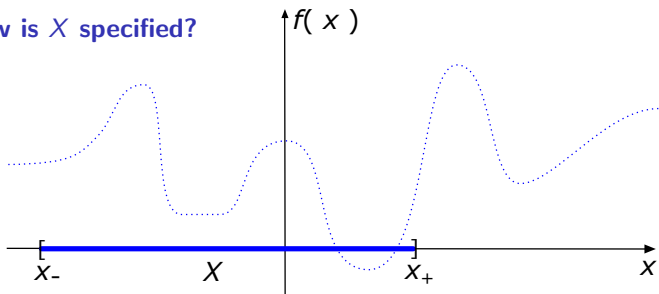
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- ▶ Standard ways: equality constraint $g(x) = v \equiv X = \text{level set } L(g, v)$, inequality constraint $g(x) \leq v \equiv \text{sublevel set } S(g, v) = \{x : g(x) \leq v\}$
- ▶ For convenience “ v hidden in f ” $\implies f(x) \leq 0$, $f(x) = 0$
- ▶ What if one rather wants $g(x) \geq 0$? Simply $-g(x) \leq 0$
- ▶ Usually multiple constraints: “ $g_1(x) \leq 0$, $g_2(x) \leq 0$ ” \equiv logical conjunction (“first condition and second condition”) \equiv intersection of (sub)level sets
- ▶ Simple and common: bounds $x \leq x_+$ (up) / $x \geq x_-$ (dn), boxes $x_- \leq x \leq x_+$

Outline

Optimization Problems

Optimization is difficult

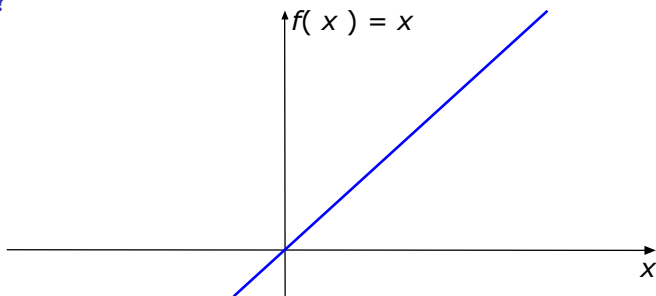
Simple Functions, Univariate case

Simple Functions, Multivariate case

Multivariate Quadratic case: Gradient Method

Wrap up & References

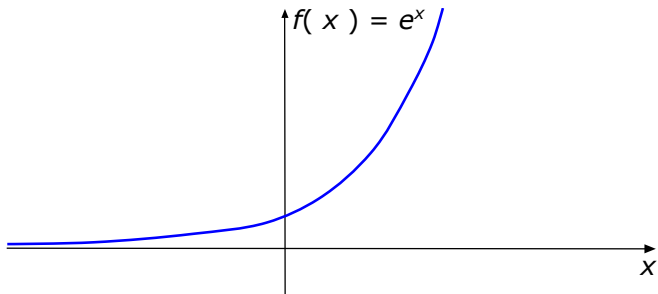
Solutions



- ▶ f has no minimum, (P) unbounded (below): $f_* = \nu(P) = -\infty$
- ▶ Just a convenient shorthand for $\forall t \in \mathbb{R} \exists x \in \mathbb{R}$ s.t. $f(x) \leq t$
i.e., “there is no (finite) lower bound on $\text{im}(f)$ ”
- ▶ Solving (P) actually (at least) two different things:
 - ▶ finding x_* and proving it is optimal (how??)
 - ▶ constructively proving f unbounded below (how??)
- ▶ Hardly ever happens in learning since $\mathcal{L}(w) \geq 0$
- ▶ Nontrivial and important in optimization (tied with duality, nonemptiness, ...)

What if $f_* \exists$ but $x_* \nexists$?

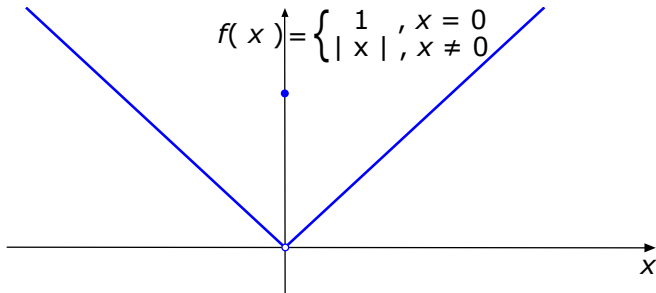
7



- ▶ $\text{im}(f)$ is bounded below but has no minimum
- ▶ Either “naturally”

What if $f_* \exists$ but $x_* \nexists$?

7



- ▶ $\text{im}(f)$ is bounded below but has no minimum
- ▶ Either “naturally” or “forcibly”
- ▶ $\inf\{f(x) : x \in \mathbb{R}\} \exists$, but $\min\{f(x) : x \in \mathbb{R}\} \nexists$
- ▶ Arguably $f_* = \inf\{f(x) : x \in \mathbb{R}\}$, but $\nexists x_*$ s.t. $f_* = f(x_*)$
- ▶ $\text{im}(f)$ is open, does not contain its boundary (will see)

- ▶ \mathbb{R} totally ordered $\implies \forall x, y \in \mathbb{R}$, at least one among $x \leq y$, $y \leq x$ holds
- ▶ $S \subseteq \mathbb{R}$, $\underline{s} = \inf S \iff \underline{s} \leq s \ \forall s \in S \ \wedge \ \forall t > \underline{s} \exists s \in S \text{ s.t. } s \leq t$
 $\bar{s} = \sup S \iff \bar{s} \geq s \ \forall s \in S \ \wedge \ \forall t < \bar{s} \exists s \in S \text{ s.t. } s \geq t$
- ▶ $\underline{s} \in S \implies \underline{s} = \min S$, $\bar{s} \in S \implies \bar{s} = \max S$
- ▶ Issues: i) $\inf S / \sup S$ may not \exists in \mathbb{R} , ii) $\inf S / \sup S$ may not $\in S$
- ▶ Should write “ $\inf\{f(x) \dots\}$ ”, but we want (approximately) optimal solutions
- ▶ Set of extended reals: $\overline{\mathbb{R}} = \{-\infty\} \cup \mathbb{R} \cup \{+\infty\}$ (usually just \mathbb{R})
- ▶ For all $S \subseteq \mathbb{R}$, $\exists \sup / \inf S \in \overline{\mathbb{R}}$
- ▶ $\inf S = -\infty \iff \forall t \in \mathbb{R} \exists s \in S \text{ s.t. } s \leq t$
 $\sup S = +\infty \iff \forall t \in \mathbb{R} \exists s \in S \text{ s.t. } s \geq t$
 just a convenient shorthand for “there is no (finite) inf / sup”
- ▶ $\inf \emptyset = \infty$, $\sup \emptyset = -\infty$

- ▶ Several ways to ensure this never happens (hypotheses on f , X)
- ▶ On **computers** “ $x \in \mathbb{R}$ ” typically is “ $x \in \mathbb{Q}$ ” with **up to 16 digits precision**
 \implies approximation errors unavoidable anyway
- ▶ Exact algebraic computation **may be possible** (if f is algebraic, which it may be **not**) but anyway usually **too slow**
- ▶ In fact **learning going the opposite way** (float, half, FP8, ...)
- ▶ Anyway, **finding the exact x_* impossible in general** [4, p. 408]
- ▶ For **any fixed $\varepsilon > 0$** , plenty of **ε -approximate solutions** (ε -optima):
$$x_\varepsilon \in \mathbb{R} \text{ s.t. } f_* \leq f(x_\varepsilon) \leq f_* + \varepsilon$$

“as close to the optimal solution (value) as you want”
- ▶ **Cost of solution algorithms typically depends on ε** (sometimes **very badly**)
- ▶ And ε can't really become very small anyway (see above)

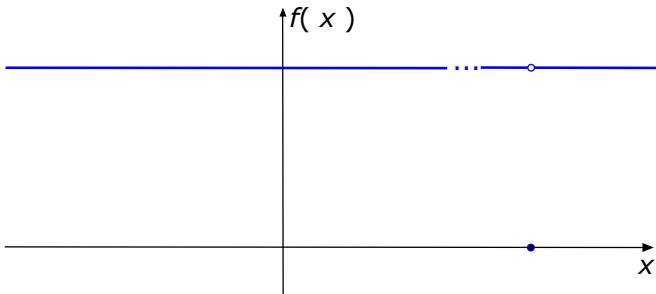
- ▶ Absolute gap: $A(x) = f(x) - f_*$ (≥ 0)
- ▶ Relative gap: $R(x) = (f(x) - f_*) / |f_*| = A(x) / |f_*|$ (≥ 0)
- ▶ Why $R(x)$? Because $\forall \alpha > 0 \quad (P) \equiv (P_\alpha) \min\{\alpha f(x) : x \in \mathbb{R}\}$
 $\nu(P_\alpha) = \alpha f_* = \alpha \nu(P) \implies$ same $R(x)$ (scale invariant), different $A(x)$

Exercise: $R(x)$ ill-defined if $f_* = 0$, propose solutions & justify them (change f_*)

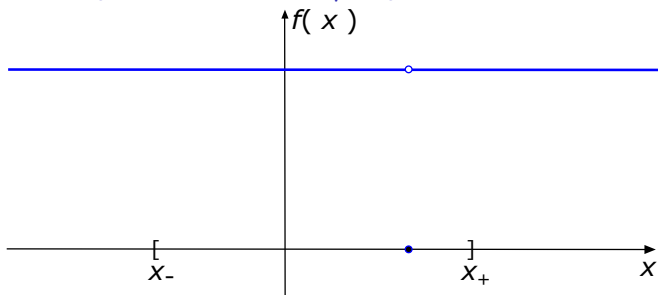
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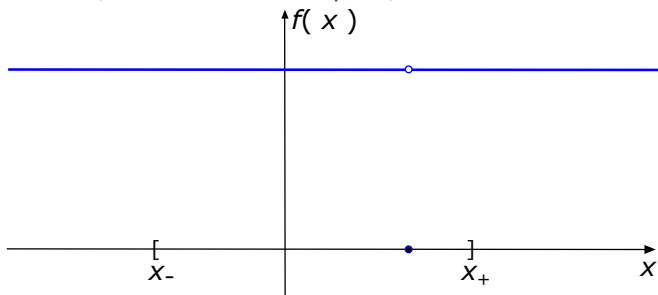
- ▶ (Approximately) solve (P) : fix ε , find x s.t. either $A(x) \leq \varepsilon$ or $R(x) \leq \varepsilon$
- ▶ Issue: computing $A(x)$ or $R(x)$ requires f_* which is typically unknown
- ▶ Could argue this is “the issue” in optimization: compute (an estimate of) f_*
- ▶ Sometimes \approx known in learning ($f_* \approx 0$ in NN, but not in SVM)
- ▶ A real issue only if global optimum x_* needed, hence not always



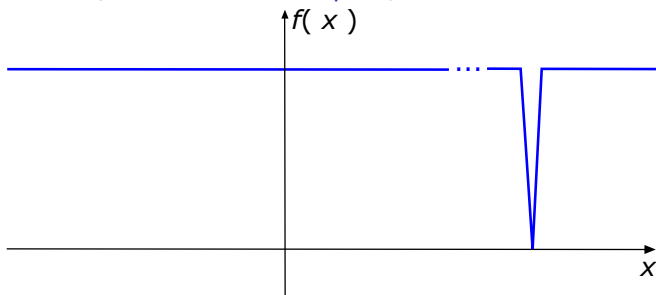
- Impossible because isolated minima can be anywhere [4, p. 408]



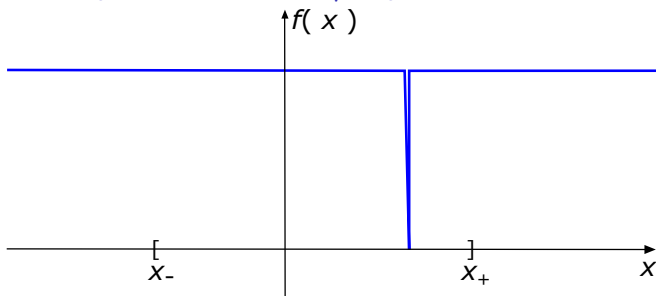
- ▶ **Impossible** because **isolated minima can be anywhere** [4, p. 408]
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- ▶ **No**: still **uncountably many** points to try



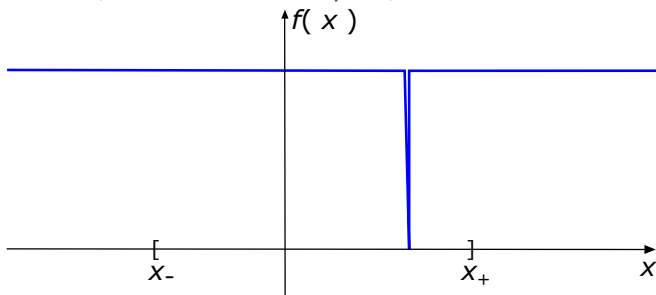
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- ▶ Is it because f “jumps”? **No**, f can have **isolated \downarrow spikes** anywhere ... even on $X = [x_-, x_+]$ as **spikes can be arbitrarily narrow**
- ▶ To make (even approximate) optimization even possible, f must be “nice”
- ▶ Let's start with the **nicest possible ones** where optimization is (\approx) trivial

Outline

Optimization Problems

Optimization is difficult

Simple Functions, Univariate case

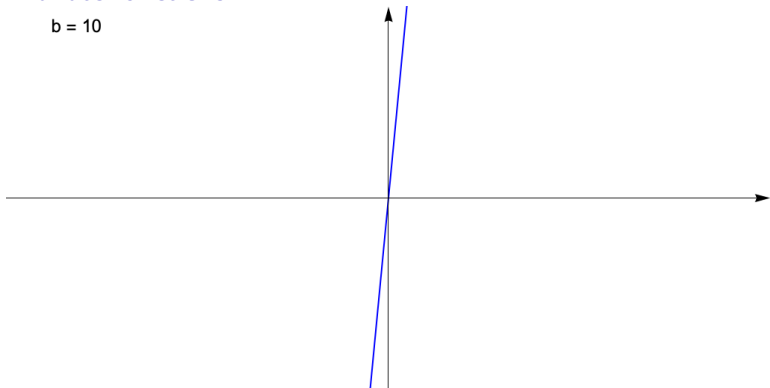
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Wrap up & References

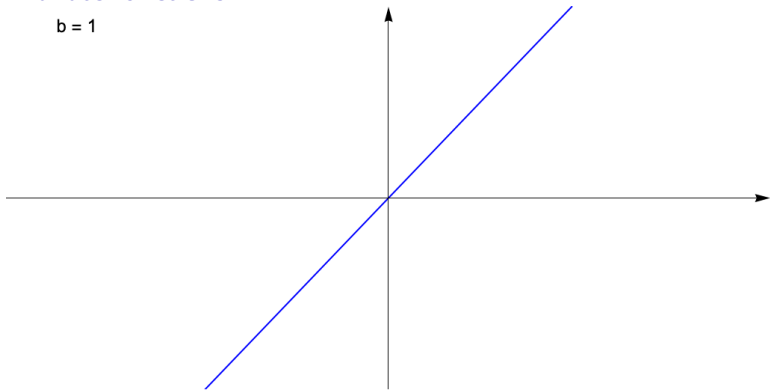
Solutions

$b = 10$



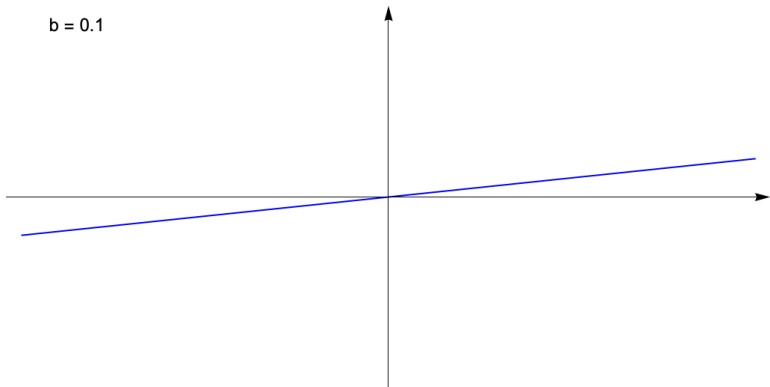
- ▶ The simplest possible function: $f(x) = bx$ (linear), fixed $b \in \mathbb{R}$
- ▶ As many different functions as real numbers (bijection)
- ▶ $b > 0 \equiv$ increasing: $x > z \implies f(x) > f(z)$

$$b = 1$$



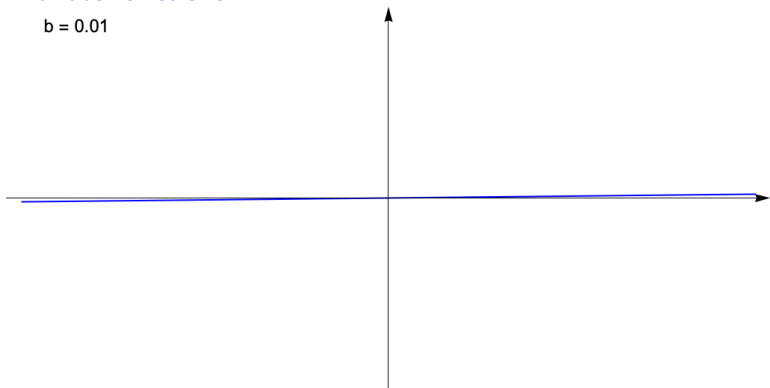
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$b = 0.1$



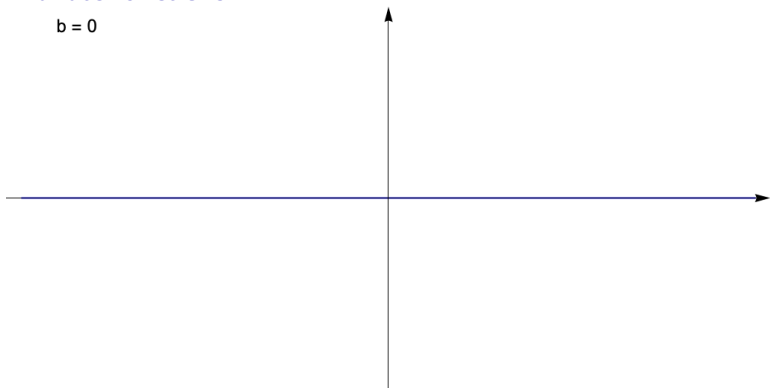
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$b = 0.01$



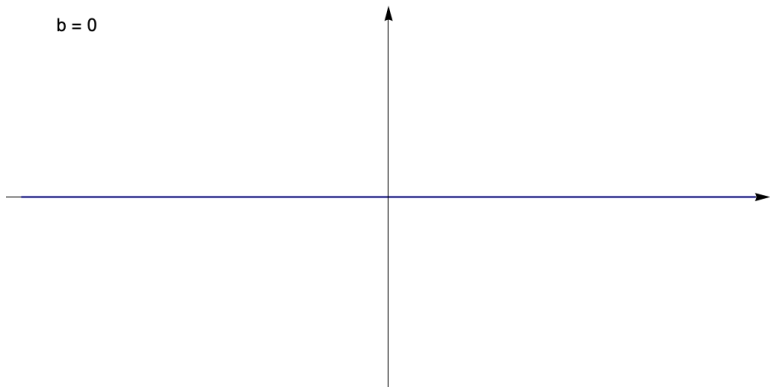
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$$b = 0$$



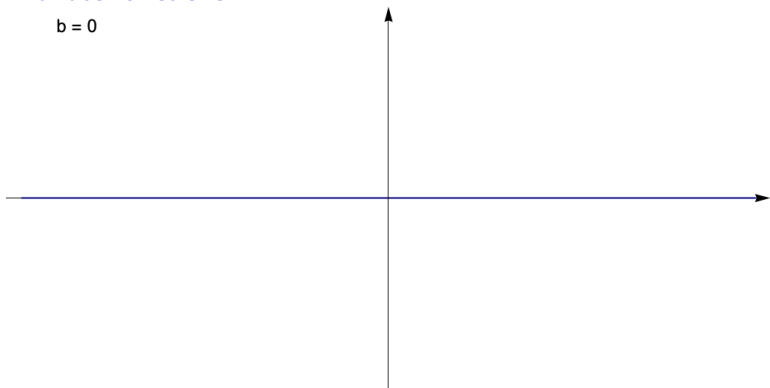
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$$b = 0$$



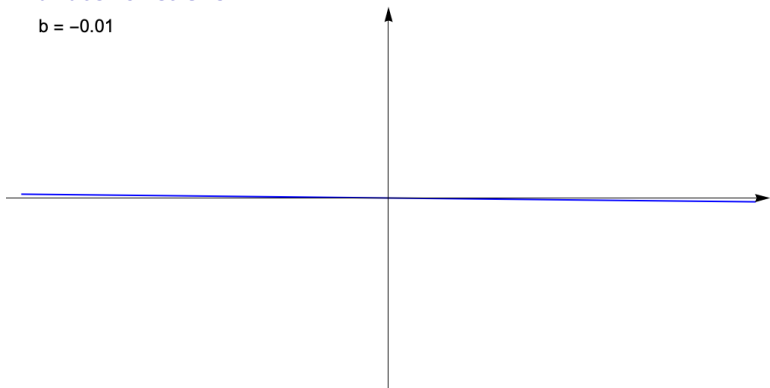
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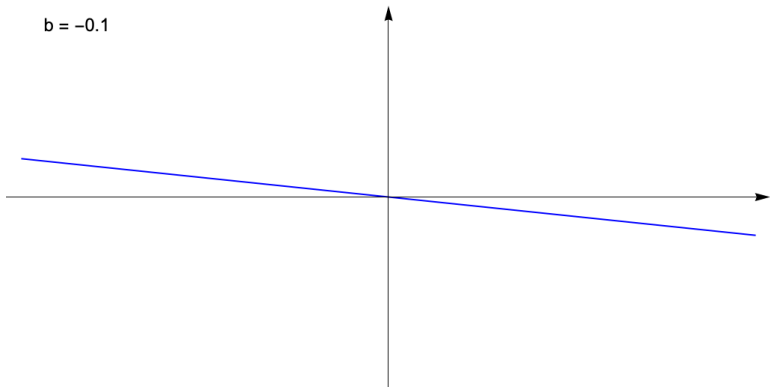
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- ▶ $b = 0 \equiv$ constant: $x > z \implies f(x) = f(z)$

$$b = -0.01$$

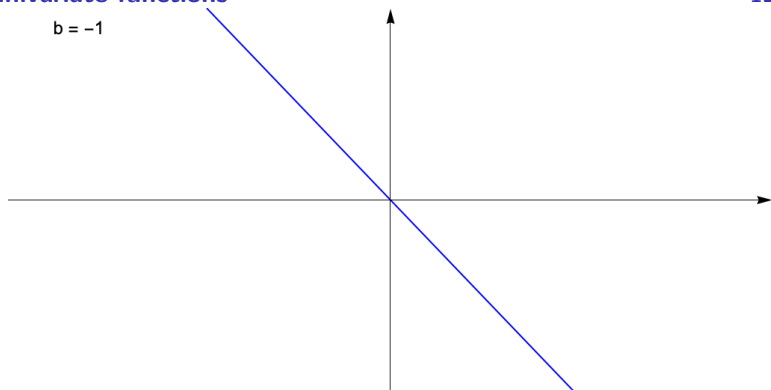


- ▶ The simplest possible function: $f(x) = bx$ (linear), fixed $b \in \mathbb{R}$
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- ▶ $b < 0 \equiv$ decreasing: $x > z \implies f(x) < f(z)$

$$b = -0.1$$

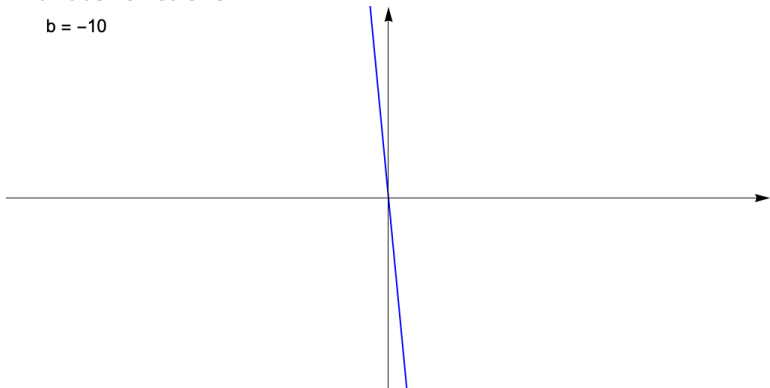


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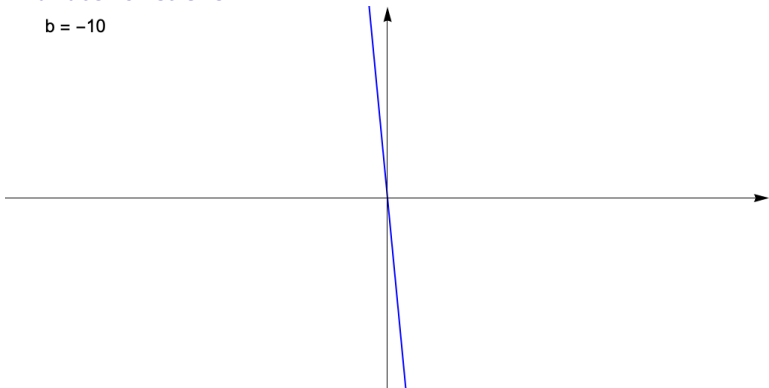
$$b = -10$$



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Exercise: Formally prove the stated properties

$$b = -10$$



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Exercise: Formally prove the stated properties

- ▶ $b =$ linear coefficient = slope: the larger $|b|$, the steeper the line

▶ Too easy: $\min = -\infty$, $\max = +\infty$ unless $b = 0 \implies \min = \max = 0$

▶ More interesting: box-constrained optimization

$$(P) \min\{f(x) : x \in [x_-, x_+]\}$$

with $-\infty \leq x_- \leq x_+ \leq +\infty \equiv X$ possibly (half-)infinite interval

▶ Constraints often useful, (finite) box constraints (very simple) especially so

▶ $b > 0 \implies \operatorname{argmin} = x_-$, $\min = f(x_-)$, $\operatorname{argmax} = x_+$, $\max = f(x_+)$

▶ “Works” even if $x_- = -\infty$ and/or $x_+ = +\infty$, as $b \cdot (\pm\infty) = \pm\infty$

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Exercise: Formally prove the result, state & prove cases $b < 0$ and $b = 0$

▶ Closed formula $O(1)$, don't get used to it

▶ Yet solving simple problems the basis of solving complicated ones

- ▶ Could have used $X = (x_-, x_+) = \{x \in \mathbb{R} : x_- < x < x_+\}$ (open interval)?

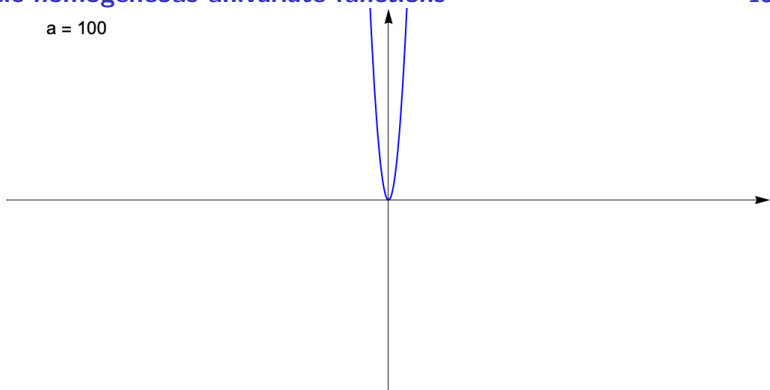
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 \implies cannot be chosen/measured to ∞ precision
(Planck scale, Heisenberg's Uncertainty Principle, ...)

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- ▶ It is a problem for algorithms? In theory yes, in practice hardly:
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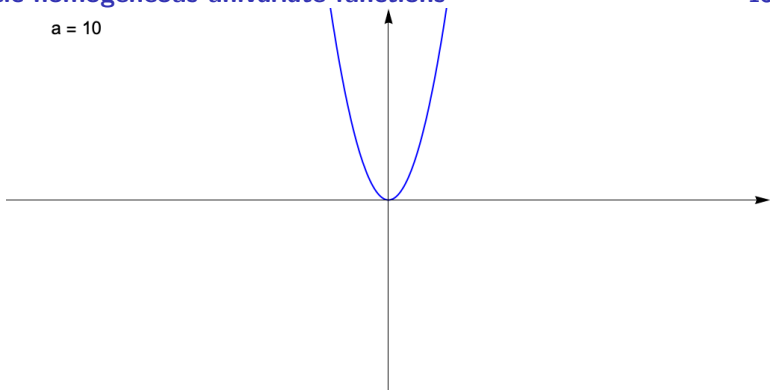
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- ▶ It is a problem for algorithms? In theory yes, in practice hardly: again, plenty of ε -optimal solutions however chosen $\varepsilon > 0$
- ▶ Does it make any sense at all? Hardly: if x_-, x_+ “can't be touched”, use $X = [x_- + \varepsilon_-, x_+ - \varepsilon_+]$ for appropriately chosen ε_{\pm}
- ▶ All in all? Just use closed intervals and be done with it
- ▶ Will generalise to “just use closed sets and be done with it”

$a = 100$

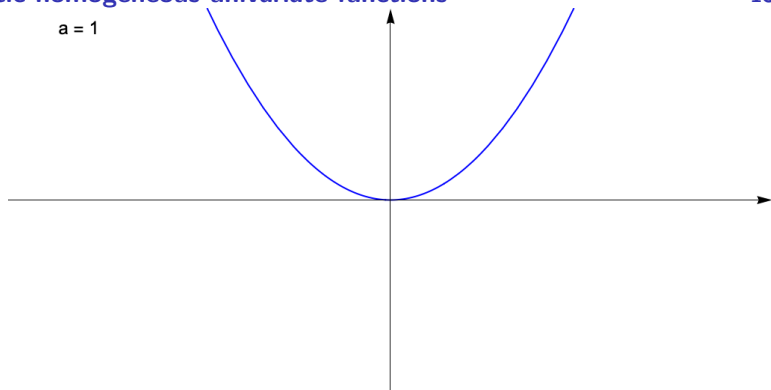


- ▶ Next simplest function: $f(x) = ax^2$ (homogeneous quadratic), fixed $a \in \mathbb{R}$
- ▶ As many different functions as real numbers (bijection)
- ▶ $a > 0 \equiv$ decreasing for $x \leq 0$, increasing for $x \geq 0$

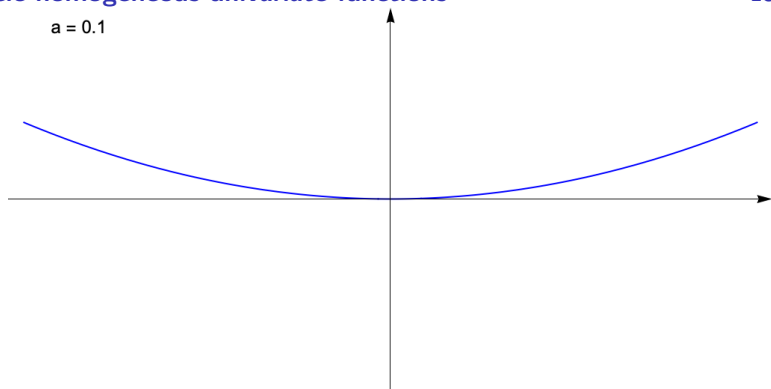
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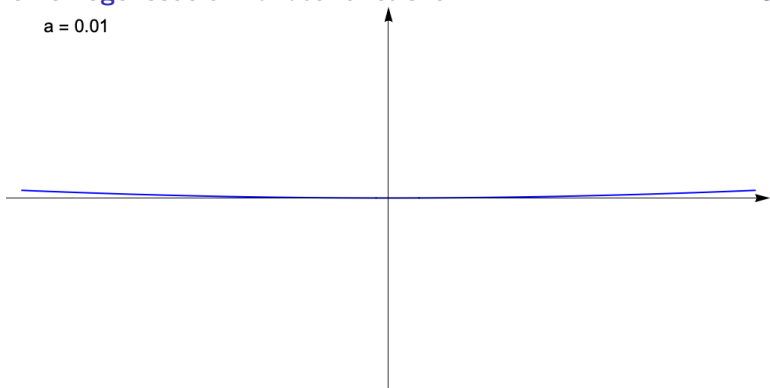


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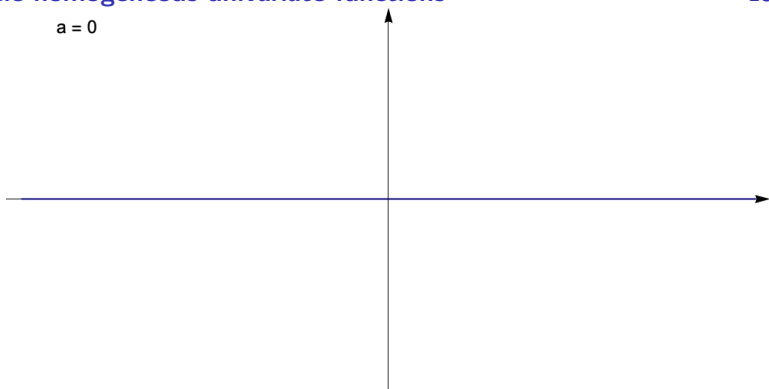


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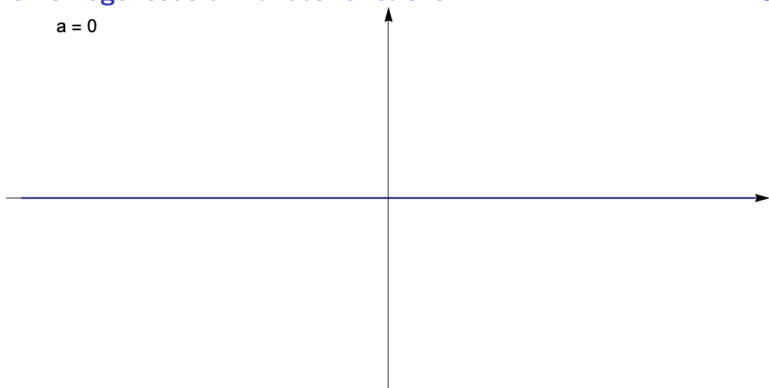
$a = 0.01$



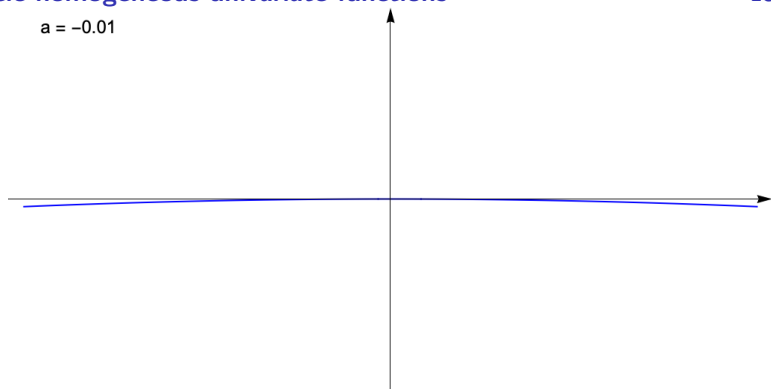
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- ▶ Next simplest function: $f(x) = ax^2$ (homogeneous quadratic), fixed $a \in \mathbb{R}$
- ▶ As many different functions as real numbers (bijection)
- ▶ $a = 0 \equiv$ nonincreasing for $x \leq 0$, nondecreasing for $x \geq 0$ and

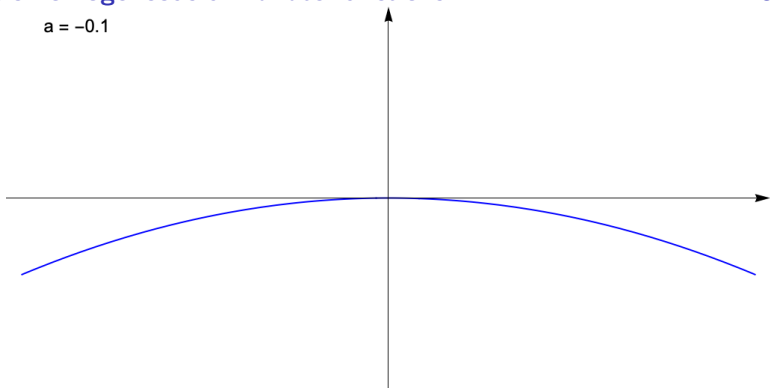


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- ▶ As many different functions as real numbers (bijection)
- ▶ $a = 0 \equiv$ nondecreasing for $x \leq 0$, nonincreasing for $x \geq 0$ (constant)

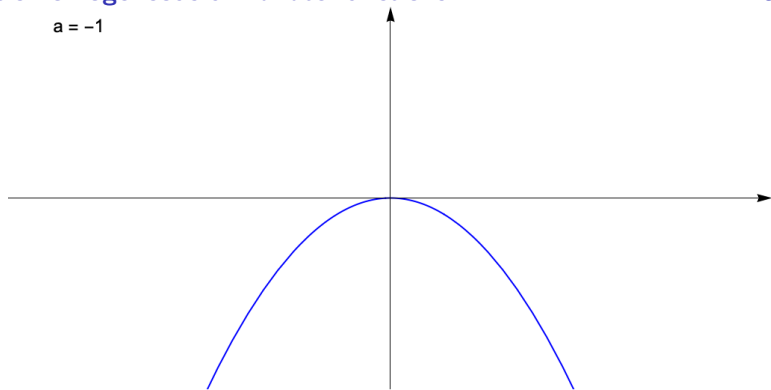


- ▶ Next simplest function: $f(x) = ax^2$ (homogeneous quadratic), fixed $a \in \mathbb{R}$
- ▶ As many different functions as real numbers (bijection)
- ▶ $a < 0 \equiv$ increasing for $x \leq 0$, decreasing for $x \geq 0$

$a = -0.1$

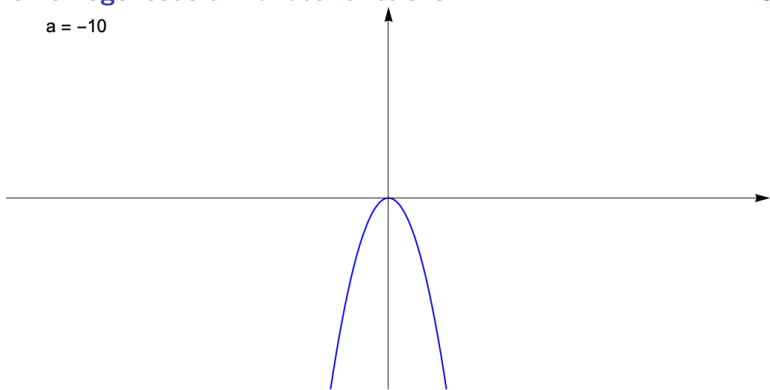


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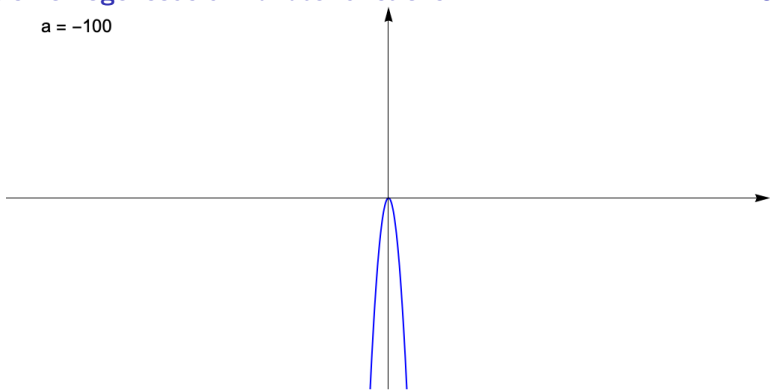
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$$a = -10$$



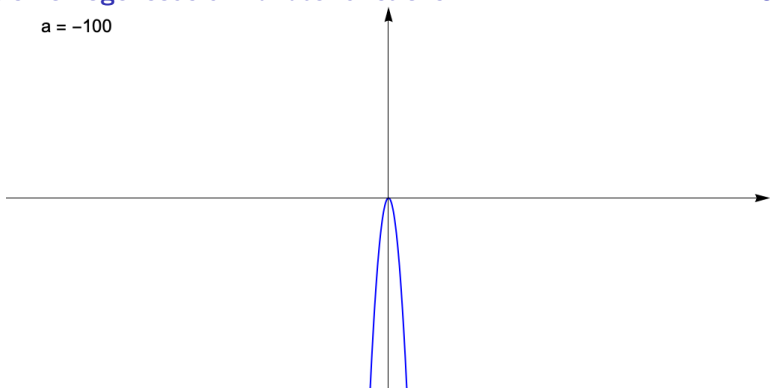
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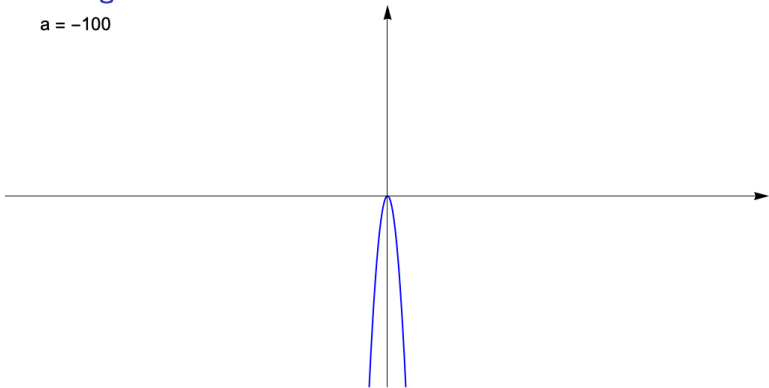
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Exercise: Formally prove the stated properties

$$a = -100$$



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Exercise: Formally prove the stated properties

- ▶ $a =$ quadratic coefficient = **curvature**: the larger $|a|$, the steeper the parabola

- ▶ Clearly depends (and symmetric) on **sign of a** :
 - ▶ $a > 0 \implies \min = \operatorname{argmin} = 0, \max = +\infty, \operatorname{argmax} = \pm\infty$
 - ▶ $a < 0 \implies \max = \operatorname{argmax} = 0, \min = -\infty, \operatorname{argmin} = \pm\infty$
- ▶ Box-constrained optimization on (closed) $X = [x_-, x_+]$ more interesting
- ▶ $a > 0 \implies$ **three cases**
 - ▶ $x_+ < 0 \implies \operatorname{argmin} = x_+, \operatorname{argmax} = x_-$
 - ▶ $x_- > 0 \implies \operatorname{argmin} = x_-, \operatorname{argmax} = x_+$
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- ▶ Again closed formula $O(1)$, don't get used to it
- ▶ $\max\{f(x)\}$ and $\min\{f(x)\}$ somewhat \neq (cf. last case), will see **much** more

- ▶ Next step: $f(x) = ax^2 + bx$ (non-homogeneous quadratic), fixed $(a, b) \in \mathbb{R}^2$
- ▶ As many different functions as **pairs of** real numbers (bijection)
- ▶ Basically, a homogeneous quadratic + a linear
- ▶ However, $\min\{ax^2 + bx\} \neq \min\{ax^2\} + \min\{bx\}$
- ▶ 0 clearly always a root, but in general **not** the argmin / argmax
- ▶ Powerful general concept: if $f(x)$ is “**too complicated**”, make it “**simpler**”
- ▶ Sometimes this can be done by **changing the space of variables** (reformulation)
- ▶ In this case: **change the input space so that it becomes homogeneous**
- ▶ Clearly only needed if **both $a \neq 0$ and $b \neq 0$**

- ▶ Fundamental trick: $\bar{x} = -b/2a$ (because I say so), $z = x - \bar{x} \equiv x = z + \bar{x}$
- ▶ The z -space is the x -space where the origin is moved to \bar{x}
- ▶ Just algebra: $f(x) = a(z + \bar{x})^2 + b(z + \bar{x}) = az^2 + 2az\bar{x} + a\bar{x}^2 + bz + b\bar{x}$
 $= az^2 + (2a\bar{x} + b)z + [a\bar{x}^2 + b\bar{x}] = az^2 + f(\bar{x}) = g(z) \quad [2a\bar{x} + b = 0]$
- ▶ Translated by \bar{x} horizontally (and by $f(\bar{x})$ vertically), $f(x)$ is homogeneous
- ▶ Its argmin / argmax (depending on sign of a) is $z = 0 \equiv x = \bar{x}$
- ▶ Then, just [Optimizing a quadratic homogeneous function](#) for $g(z)$
- ▶ Yet again, closed formula $O(1)$, don't get used to it

Exercise: Flesh out the details: describe all cases in terms of f and x

Exercise: Discuss the position of \bar{x} and the roots of f depending on a, b

Outline

Optimization Problems

Optimization is difficult

Simple Functions, Univariate case

Simple Functions, Multivariate case

Multivariate Quadratic case: Gradient Method

Wrap up & References

Solutions

- ▶ Next crucial step: $f : \mathbb{R}^n \rightarrow \mathbb{R}$, i.e., $f(x_1, x_2, \dots, x_n) = f(x)$
with $x = [x_i]_{i=1}^n = [x_1, x_2, \dots, x_n] \in \mathbb{R}^n$
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- ▶ Even picturing things is more complex and **requires appropriate tools**

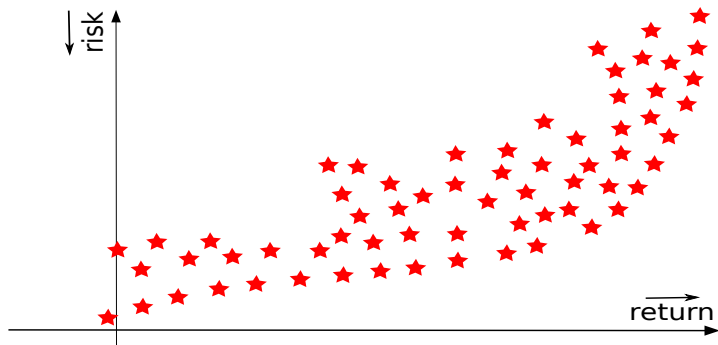
- ▶ Already “ $f : X \rightarrow \mathbb{R}$ ” a rather strong assumption:
can “express all the value of any $x \in X$ with a single number” \implies
given x' and x'' I can always tell which one I like best (\mathbb{R} has total order)

- ▶ Often there would be more than one objective:

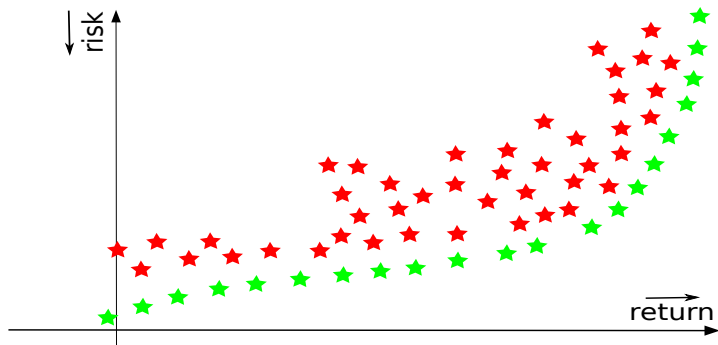
$$(P) \quad \min \{ [f_1(x), f_2(x), \dots] : x \in X \}$$

with f_1, f_2, \dots contrasting and/or with incomparable units (apples vs. oranges)

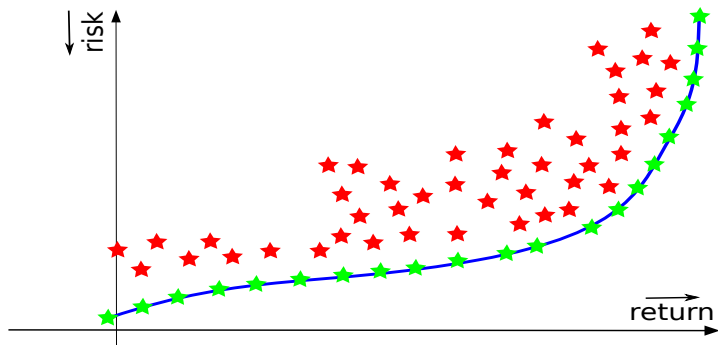
- ▶ car cost vs. flashiness vs. km/l vs. # seats vs. trunk space ...
- ▶ loss function $\mathcal{L}(w)$ vs. regularity $R(w)$ in ML
- ▶ ...
- ▶ Vector-valued (a.k.a. multi-objective) optimization: $f : X \rightarrow \mathbb{R}^k$ with $k > 1$
- ▶ Textbook example: portfolio selection problem
 - ▶ X = set of financial instruments portfolios available to buy
 - ▶ $f_1(x)$ = expected return of portfolio x (€)
 - ▶ $f_2(x)$ = risk of portfolio x not achieving the expected return (% , CVAR, ...)



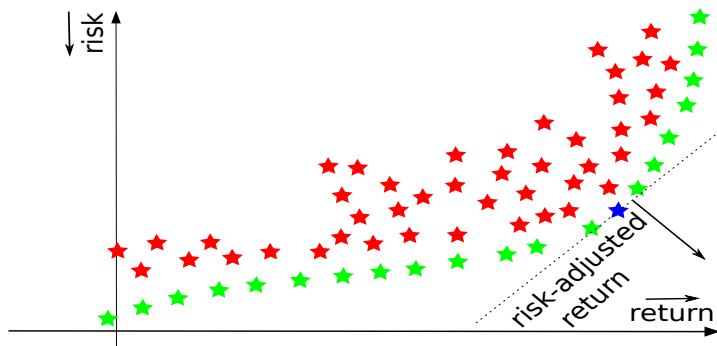
- ▶ \mathbb{R}^k with $k > 1$ has no total order \implies no “best” solution, only



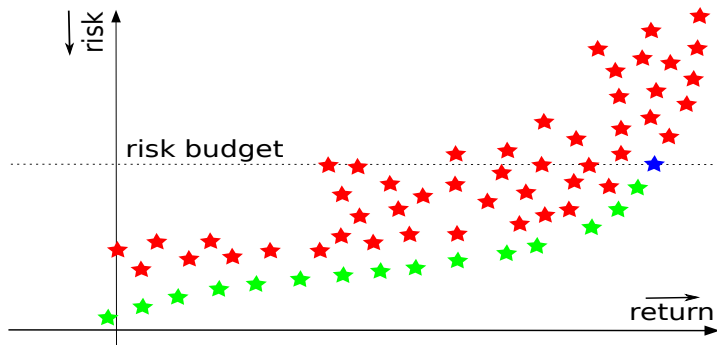
- ▶ \mathbb{R}^k with $k > 1$ has **no total order** \implies no “best” solution, only **non-dominated** ones on the



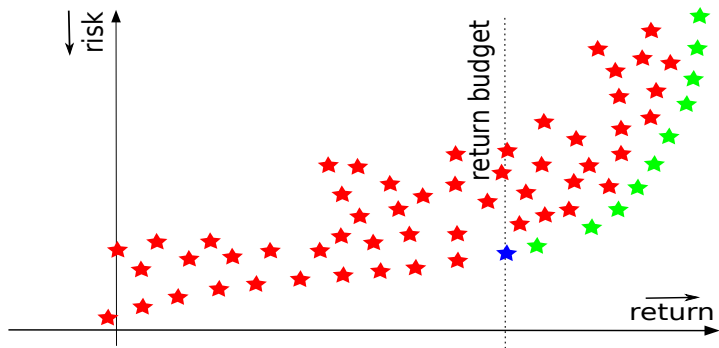
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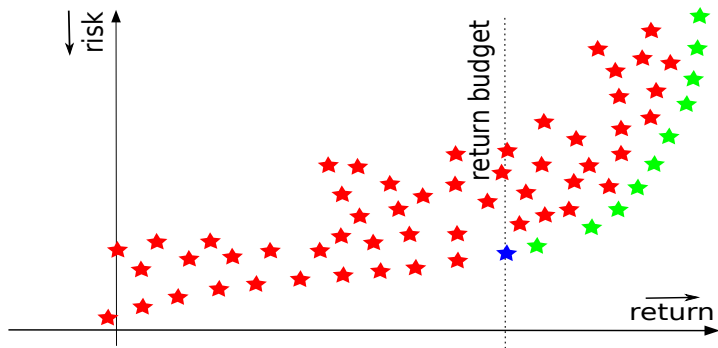
- ▶ \mathbb{R}^k with $k > 1$ has **no total order** \implies
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a.k.a. **scalarization** $\min \{ f_1(x) + \alpha f_2(x) : x \in X \}$ (**which α ??**)



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- ▶ Two practical solutions: maximize return with **budget on maximum risk**,
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- ▶ All **a bit fuzzy**, but it's the nature of the beast
- ▶ We always **assume this done** if necessary at **modelling stage** (regularization, **grid search** used to divine **hyperparameters** α, β_1, β_2)

- (Euclidean) scalar product of $x \in \mathbb{R}^n$ and $z \in \mathbb{R}^n$:

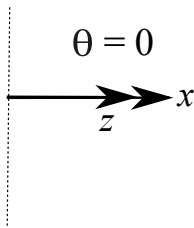
$$\langle x, z \rangle = \sum_{i=1}^n x_i z_i = x_1 z_1 + \cdots + x_n z_n$$

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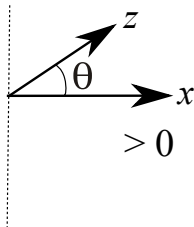


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 $\langle x, z \rangle > 0 \equiv$ “ x and z point in the **same** direction”

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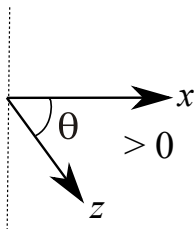


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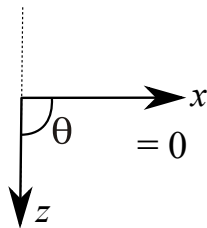


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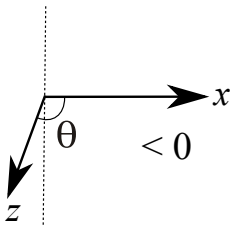
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$$\langle x, z \rangle = 0 \equiv x \perp z \text{ (orthogonal)}$$

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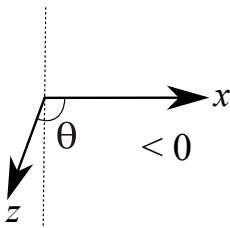


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- ▶ Cauchy-Schwarz inequality: $|\langle x, z \rangle| \leq \|x\| \|z\| \quad \forall x, z$

- ▶ (Euclidean) **distance** between x and $z =$ norm of x when z is the origin:

$$d(x, z) := \|x - z\| = \sqrt{(x_1 - z_1)^2 + \dots + (x_n - z_n)^2}$$

- ▶ **Ball**, center $x \in \mathbb{R}^n$, radius $r > 0$: $\mathcal{B}(x, r) = \{z \in \mathbb{R}^n : \|z - x\| \leq r\}$

- ▶ $\mathbb{R}^n \in$ vector space \equiv closed under sum and scalar multiplication

$$x + z = [x_1 + z_1, \dots, x_n + z_n] , \quad \alpha x = [\alpha x_1, \dots, \alpha x_n]$$

- ▶ **Finite-dimensional** vector space: $\{u^i\}_{i=1}^n$ **finite base** s.t. $\forall x \in \mathbb{R}^n \exists \alpha_1, \dots, \alpha_n$
s.t. $x = \alpha_1 u^1 + \dots + \alpha_n u^n$ (canonical base: $u_i^i = 1, u_h^i = 0$ for $h \neq i, \alpha_i = x_i$)

- ▶ Not all vector spaces are finite-dimensional (function spaces, ...)

- ▶ Properties \equiv **definition of scalar product**:

1. $\langle x, z \rangle = \langle z, x \rangle \quad \forall x, z \in \mathbb{R}^n$ (symmetry)
2. $\langle x, x \rangle \geq 0 \quad \forall x \in \mathbb{R}^n, \quad \langle x, x \rangle = 0 \iff x = 0$
3. $\langle \alpha x, z \rangle = \alpha \langle x, z \rangle \quad \forall x \in \mathbb{R}^n, \alpha \in \mathbb{R}$
4. $\langle x + w, z \rangle = \langle x, z \rangle + \langle w, z \rangle \quad \forall x, w, z \in \mathbb{R}^n$

- ▶ \exists **other scalar products** that make sense in other spaces
(matrices, integrable functions, random variables, ...)

- ▶ Not just theoretical stuff (cf. **kernel in SVM**)

► Properties \equiv definition of norm:

1. $\|x\| \geq 0 \quad \forall x \in \mathbb{R}^n$, $\|x\| = 0 \iff x = 0$

2. $\|\alpha x\| = |\alpha| \|x\| \quad \forall x \in \mathbb{R}^n$, $\alpha \in \mathbb{R}$

3. $\|x + z\| \leq \|x\| + \|z\| \quad \forall x, z \in \mathbb{R}^n$ (triangle inequality)

► $\|x + z\|^2 = \|x\|^2 + \|z\|^2 + 2\langle x, z \rangle$ (only Euclidean norm)

► $2\|x\|^2 + 2\|z\|^2 = \|x + z\|^2 + \|x - z\|^2$ (Parallelogram Law)

► Properties \equiv definition of distance:

1. $d(x, z) \geq 0 \quad \forall x, z \in \mathbb{R}^n$, $d(x, z) = 0 \iff x = z$

2. $d(\alpha x, 0) = |\alpha| d(x, 0) \quad \forall x \in \mathbb{R}^n$, $\alpha \in \mathbb{R}$

3. $d(x, w) \leq d(x, z) + d(z, w) \quad \forall x, w, z \in \mathbb{R}^n$ (triangle inequality)

► $\|\cdot\|$ defines $\mathcal{B}(\cdot, \cdot) \equiv$ the topology of the vector space:
what is next to what (will be useful later on)

- ▶ $\text{gr}(f) \in \mathbb{R}^{n+1}$, impossible if $n > 3$ ($n = 3$ hard already)
- ▶ $L(f, \cdot) \in \mathbb{R}^n$, impossible if $n > 4$ ($n = 4$ hard already)

- ▶ $\text{gr}(f) \in \mathbb{R}^{n+1}$, impossible if $n > 3$ ($n = 3$ hard already)
- ▶ $L(f, \cdot) \in \mathbb{R}^n$, impossible if $n > 4$ ($n = 4$ hard already)
- ▶ General n , $f : \mathbb{R}^n \rightarrow \mathbb{R}$, $x \in \mathbb{R}^n$ (origin), $d \in \mathbb{R}^n$ (direction):
 $\varphi_{x,d}(\alpha) = f(x + \alpha d) : \mathbb{R} \rightarrow \mathbb{R}$ tomography of f from x along d
- ▶ $\text{gr}(\varphi_{x,d})$ can always be pictured, but too many of them: which x, d ?
- ▶ $\|d\|$ only changes the scale: $\varphi_{x,\beta d}(\alpha) = \varphi_{x,d}(\beta\alpha)$ (check) \implies
 often (but not always) convenient to use normalised direction ($\|d\| = 1$)
- ▶ Simplest case: restriction along i -th coordinate ($\|u^i\| = 1$)
 $f_x^i(\alpha) = f(x_1, \dots, x_{i-1}, \alpha, x_{i+1}, \dots, x_n) \equiv \varphi_{[x_1, \dots, x_{i-1}, 0, x_{i+1}, \dots, x_n], u^i}(\alpha)$
- ▶ For small n can “look at all d ”
- ▶ Otherwise, find the specific d that “shows what you want to see”
- ▶ When x and d clear from context (will happen a lot), just $\varphi(\alpha)$

- ▶ Linear function: $f(x) = \langle b, x \rangle = \sum_{i=1}^n b_i x_i$, fixed $b \in \mathbb{R}^n$
- ▶ Linear \equiv i. $f(\gamma x) = \gamma f(x)$, ii. $f(x+z) = f(x) + f(z) \quad \forall x, \gamma, z$

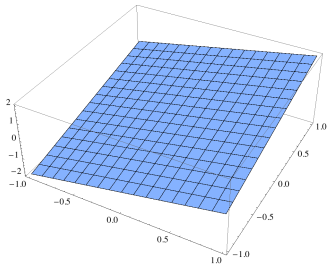
Exercise: Linear \implies i) + ii) trivial, prove \iff ; extends to **affine** ($\dots + c$)?

- ▶ $\langle b, x \rangle = \sum_{i=1}^n [f_i(x_i) = b_i x_i]$, sum of n univariate linear functions

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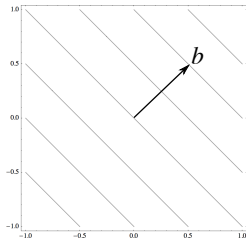
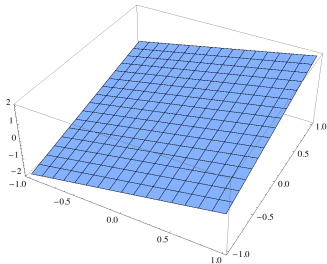


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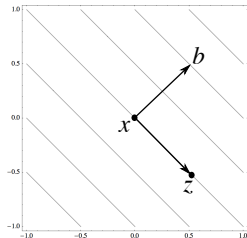
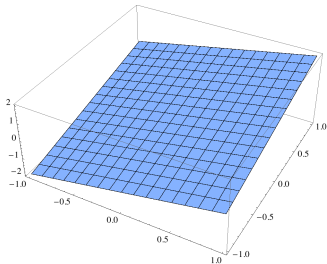


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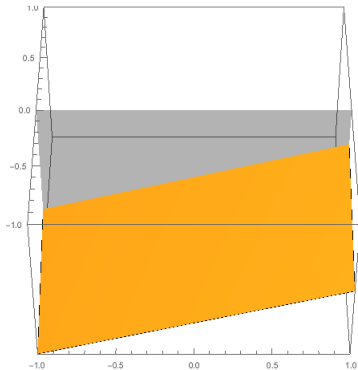
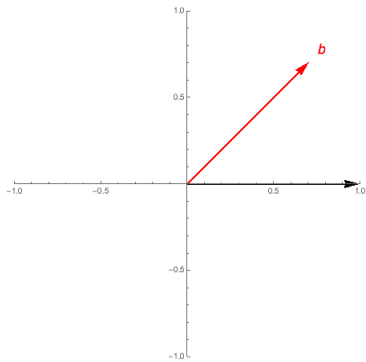
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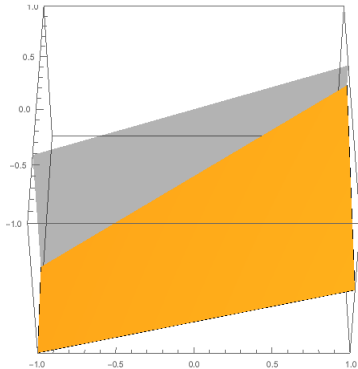
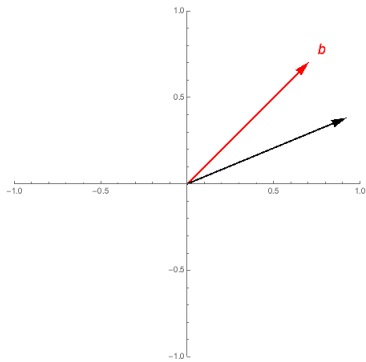
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 $f(x) = f(z) \equiv \langle b, x \rangle = \langle b, z \rangle \equiv \langle b, z - x \rangle = 0 \equiv b \perp z - x$

- ▶ $f(x) = \langle b, x \rangle$, $x = 0$, $\|d\| = 1$: $\varphi(\alpha) = \alpha \langle b, d \rangle = \alpha \|b\| \cos(\theta)$

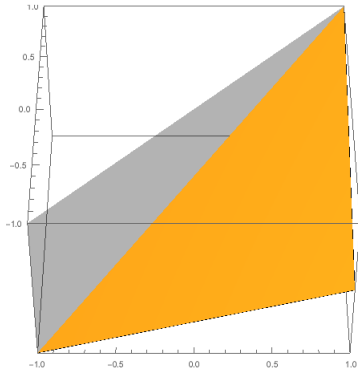
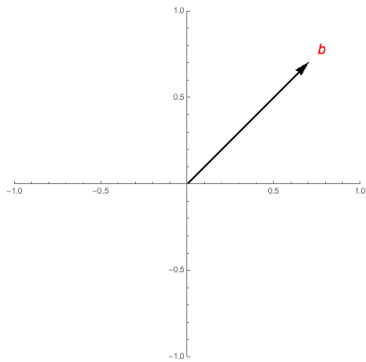
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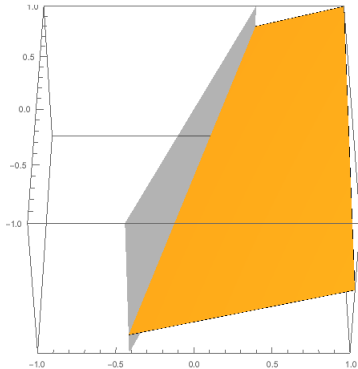
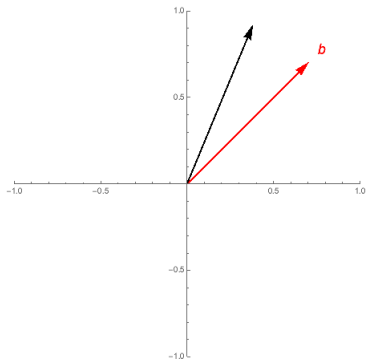
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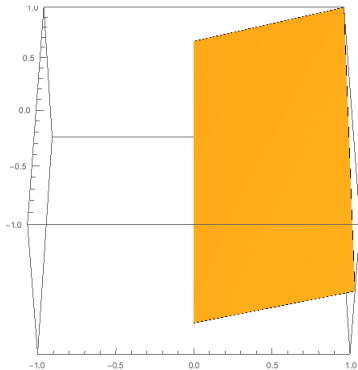
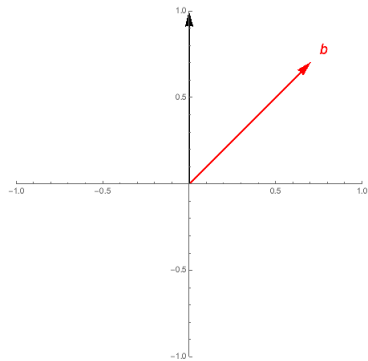
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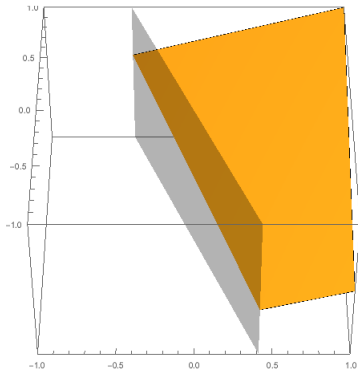
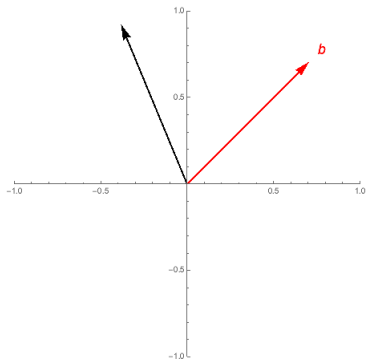
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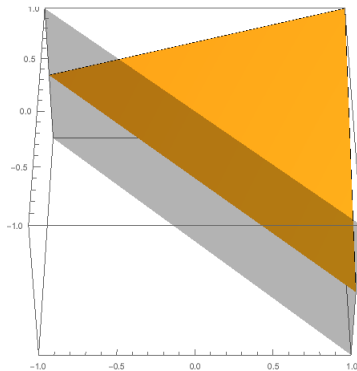
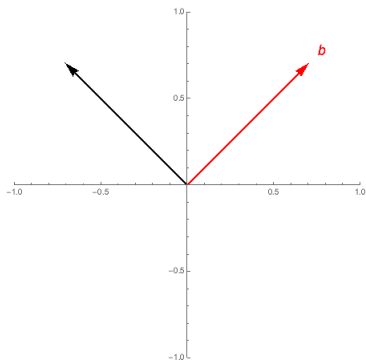
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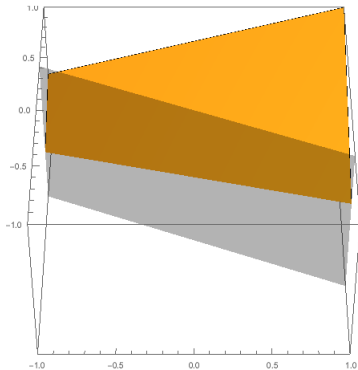
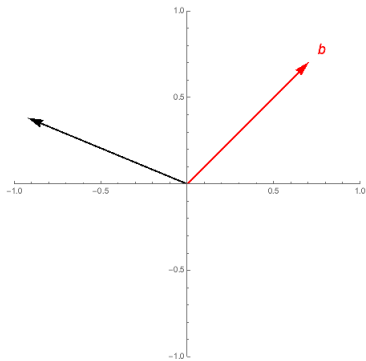
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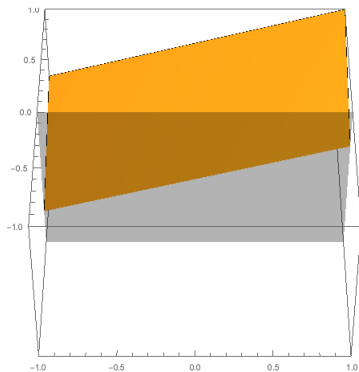
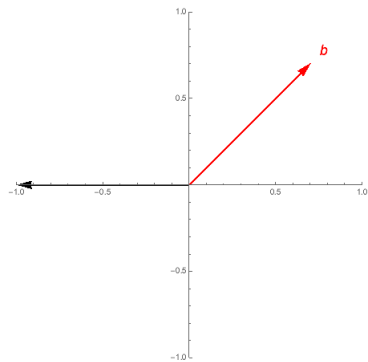
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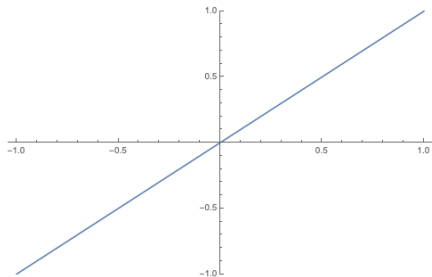
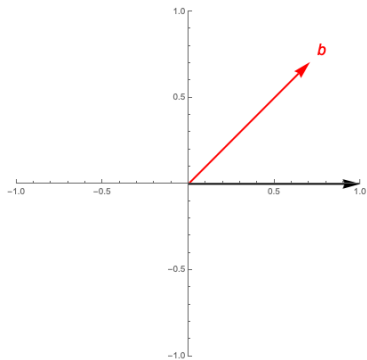
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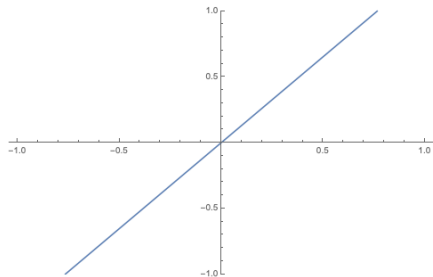
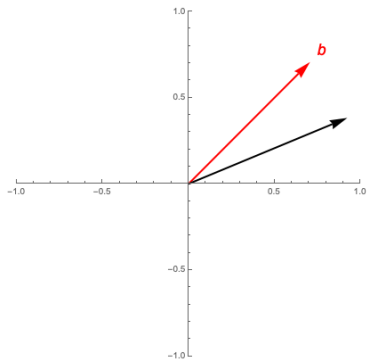


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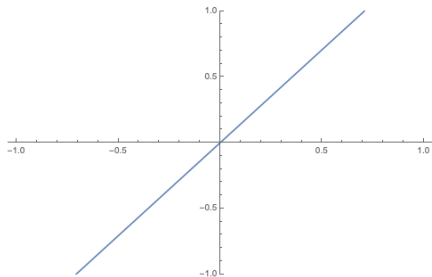
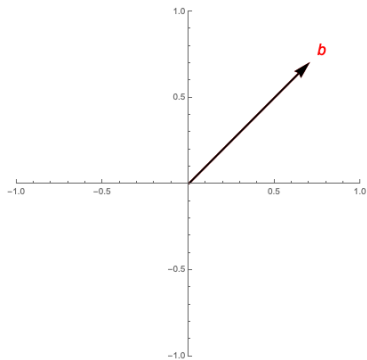
- Increasing if “ b same direction as d ”,

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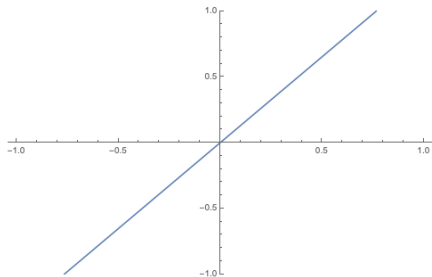
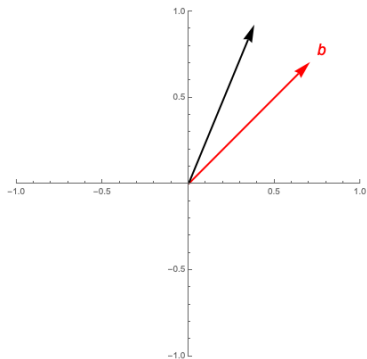
- Increasing if “ b same direction as d ”, “more collinear” \implies steeper

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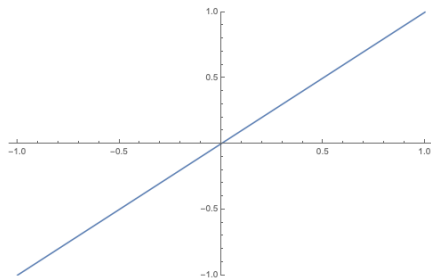
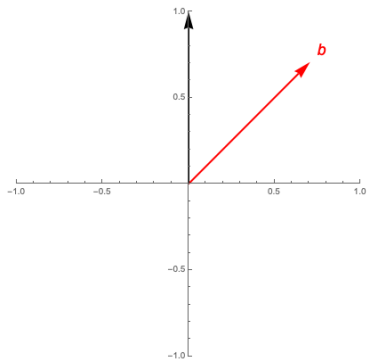
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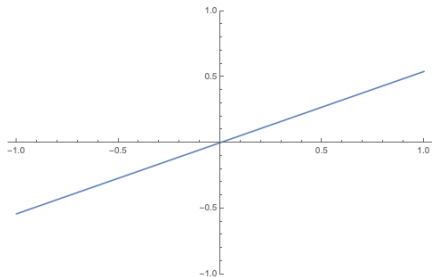
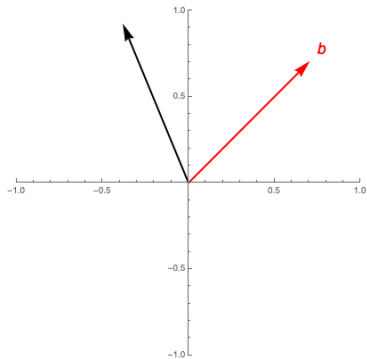
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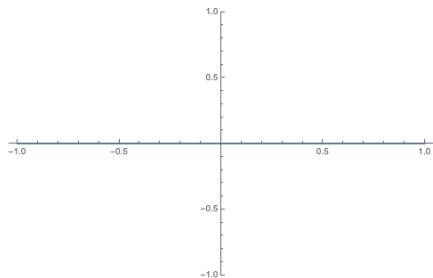
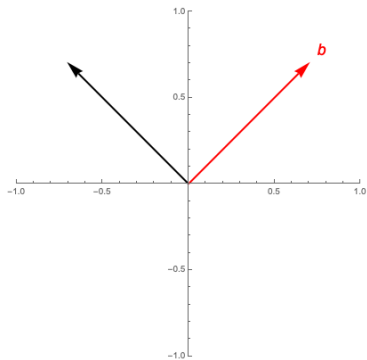
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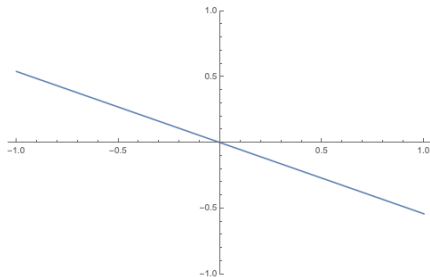
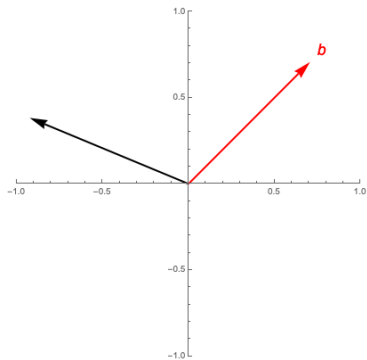
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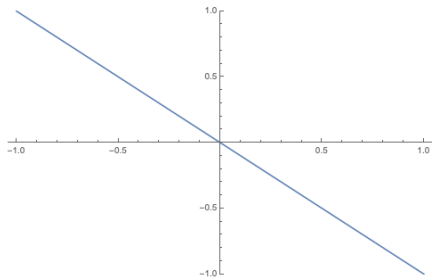
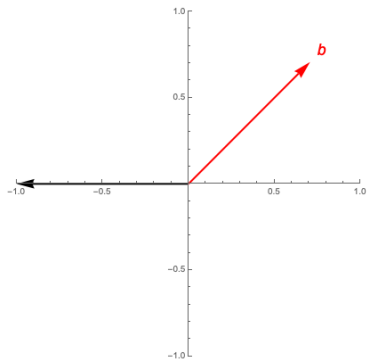
- “Flat” if $d \perp b$

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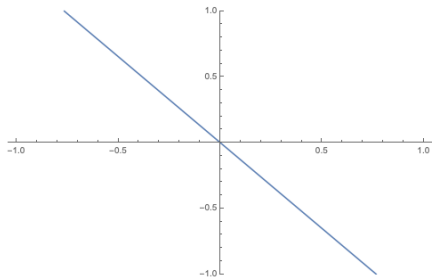
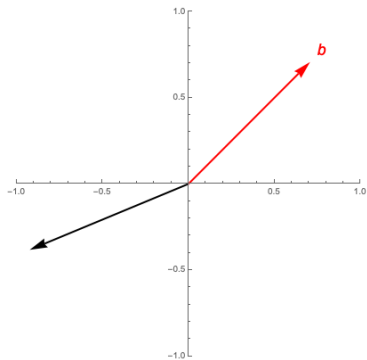
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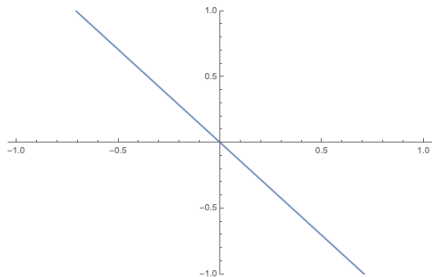
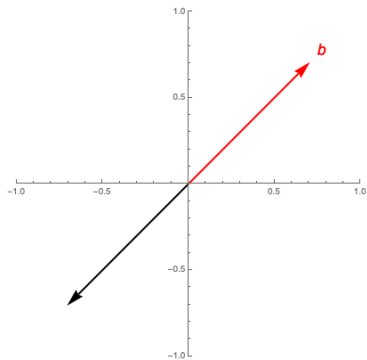
- Decreasing if “ b opposite direction as d ”, “more collinear” \implies steeper

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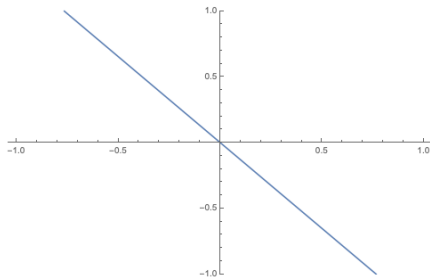
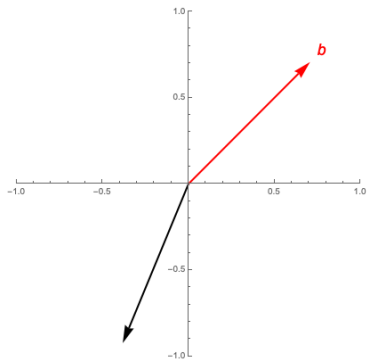
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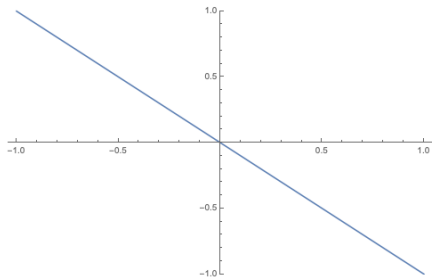
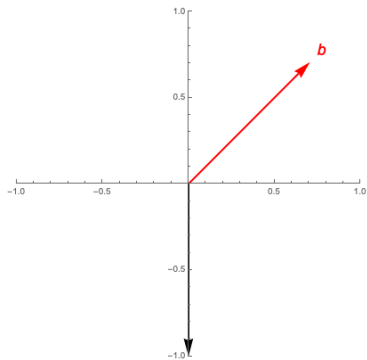
- Decreasing if “ b opposite direction as d ”, **collinear** \implies **steepest** (negative)

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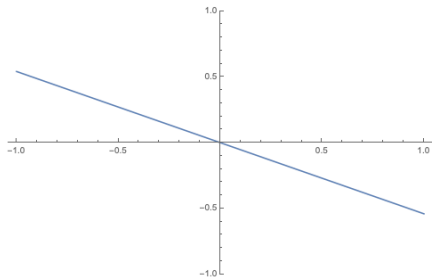
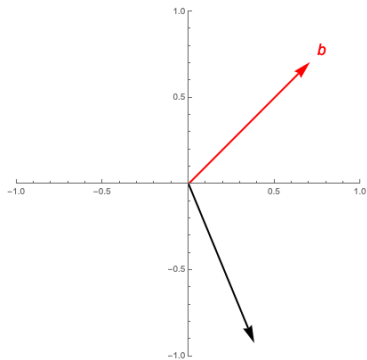
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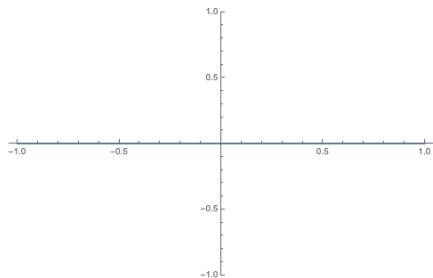
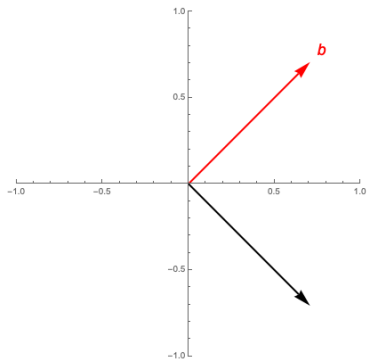
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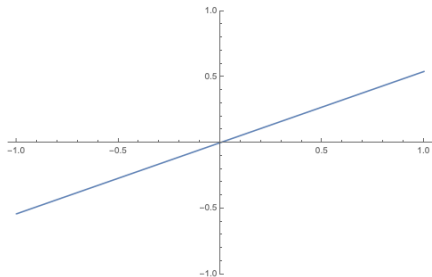
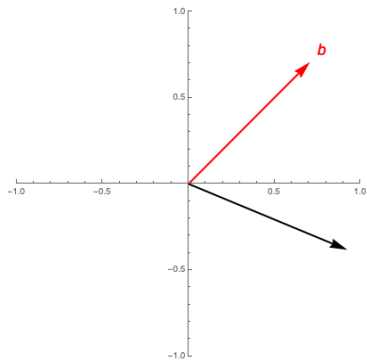
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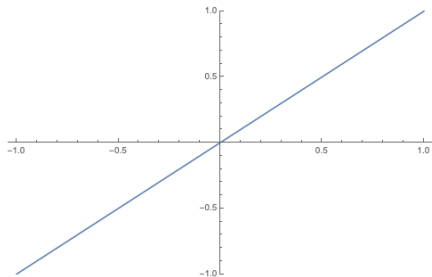
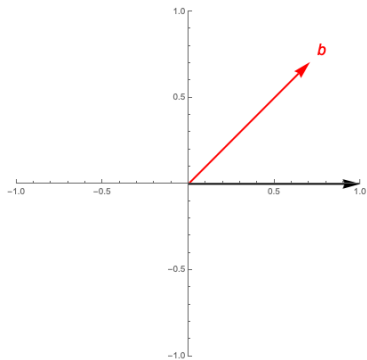
- “Flat” if $d \perp b$

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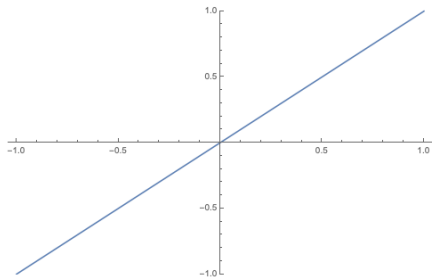
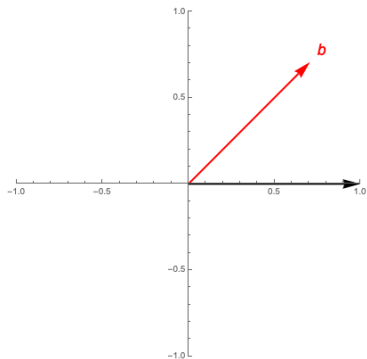
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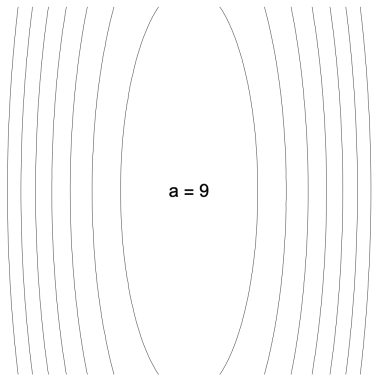
- ▶ Increasing if “ b in the same direction as d ”
- ▶ $f_* = \min\{f(x)\} = -\infty$ except if $b = 0$, in which case $f_* = 0$ (same for max)
- ▶ $\min\{f(x) : x \in X\}$, X hyperrectangle, [▶ Optimizing a linear function](#) (same for max)
 n independent problems, as nothing links x_i and x_j for $i \neq j$
- ▶ n closed formulæ $O(1)$ each, almost the last time

- ▶ Separable (non-homogeneous) quadratic function:

$$f(x) = \sum_{i=1}^n [f_i(x_i) = a_i x_i^2 + b_i x_i], \text{ fixed } (a, b) \in \mathbb{R}^{2n}$$

= sum of n univariate quadratic (non-homogeneous) functions

- ▶ $f(x) = \|x\|^2 = \sum_{i=1}^n x_i^2$ an important special case



- ▶ $f(x_1, x_2) = ax_1^2 + x_2^2 [+0x_1 + 0x_2]$

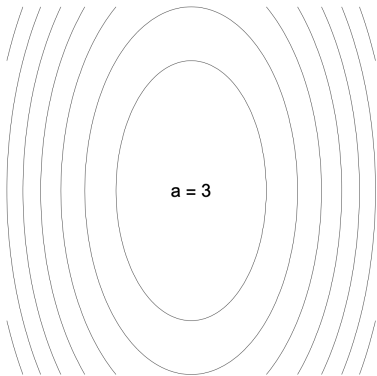
- ▶ Contour plots for different values of a

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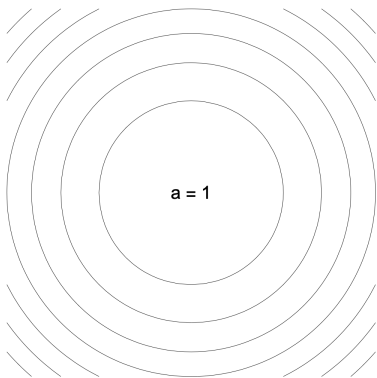
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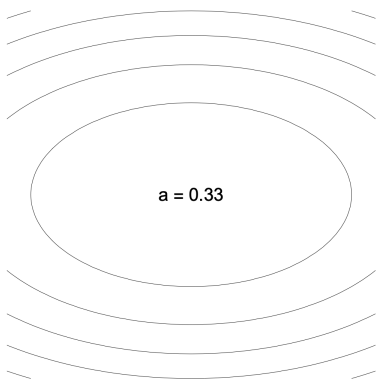
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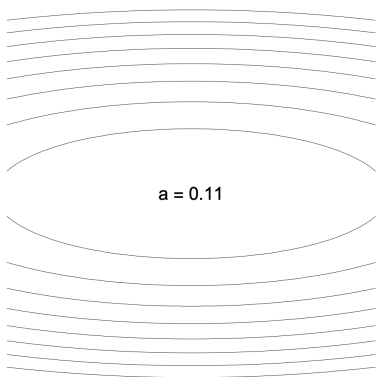
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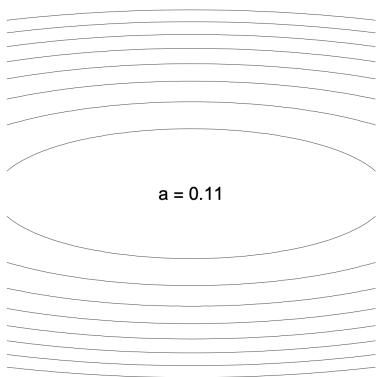
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- ▶ **Not a general quadratic** function, coming right next

- ▶ **Nonseparable** homogeneous quadratic function: fixed $Q \in \mathbb{R}^{n \times n}$ (n $Q_i \in \mathbb{R}^n$)

$$f(x) = \frac{1}{2}x^T Qx = \frac{1}{2} \left[\sum_{i=1}^n Q_{ii}x_i^2 + \sum_{i=1}^n \sum_{j=1, j \neq i}^n Q_{ij}x_i x_j \right]$$

- ▶ **Not** linear: $f(x+z) = \frac{1}{2}(x+z)^T Q(x+z) = f(x) + f(z) + z^T Qx$

- ▶ W.l.o.g. Q symmetric:

$$x^T Qx = [(x^T Qx) + (x^T Qx)^T] / 2 = x^T [(Q + Q^T) / 2]x$$

- ▶ f symmetric: $f(x) = f(-x) \implies$ "centred in $x = 0$ "

- ▶ Tomography: $\varphi(\alpha) = f(\alpha d) = \frac{1}{2}\alpha^2(d^T Qd) \implies$

homogeneous quadratic univariate, sign and steepness depend on $d^T Qd$

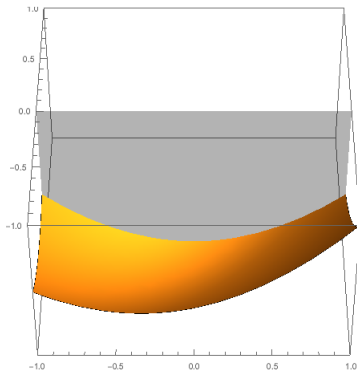
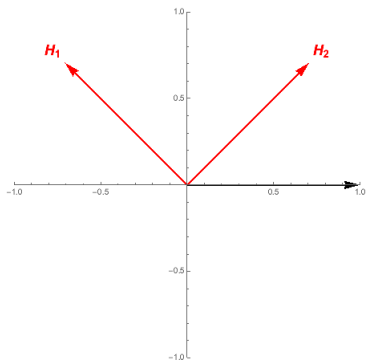
- ▶ Need to know about signs of $d^T Qd$ when d changes: (multi)linear algebra

- ▶ Crucial stuff: spectral decomposition, eigenvalues, eigenvectors of Q

- ▶ $Q \in \mathbb{R}^{n \times n}$, $v \in \mathbb{R}^n$, $\lambda \in \mathbb{R}$ s.t. $Qv = \lambda v$: v eigenvector of Q , λ eigenvalue
- ▶ v eigenvector $\equiv Qv = \lambda v \equiv Q(-v) = \lambda(-v) \equiv -v$ eigenvector
- ▶ Q symmetric \implies has n distinct eigenvectors H_1, H_2, \dots, H_n and n (not necessarily distinct) corresponding real eigenvalues $\lambda_1, \lambda_2, \dots, \lambda_n$
- ▶ Eigenvectors can always be taken orthonormal: $H_i \perp H_j$ for $i \neq j$, $\|H_i\| = 1 \implies$ linearly independent (check) \implies a(n orthonormal) basis of \mathbb{R}^n
- ▶ Spectral decomposition: $H = [H_1, \dots, H_n] \in \mathbb{R}^{n \times n}$, $\Lambda = \text{diag}(\lambda_1, \dots, \lambda_n)$
 $Q = H\Lambda H^T = \lambda_1 H_1 H_1^T + \dots + \lambda_n H_n H_n^T$ (check)
- ▶ Notation: $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$ ($\lambda_1 = \max$, $\lambda_n = \min$)
- ▶ Variational characterization of eigenvalues:
 $\lambda_1 = \max\{d^T Qd / d^T d : d \neq 0\} = \max\{d^T Qd : \|d\| = 1\}$
 $\lambda_n = \min\{d^T Qd / d^T d : d \neq 0\} = \min\{d^T Qd : \|d\| = 1\}$
- ▶ $Q \succ 0$ = positive definite if $\lambda_i > 0 \forall i \equiv \lambda_n > 0 \equiv d^T Qd > 0 \forall d \neq 0$
 $Q \succeq 0$ = positive semi-definite if $\lambda_i \geq 0 \forall i \equiv \lambda_n \geq 0 \equiv d^T Qd \geq 0 \forall d \neq 0$
 negative definite (\prec), semi-definite (\preceq), indefinite (\times) obvious

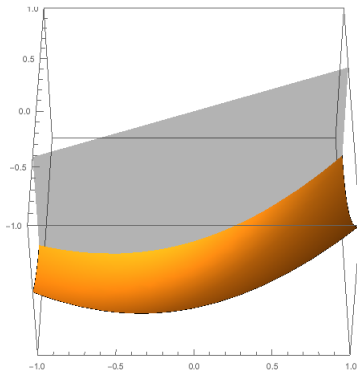
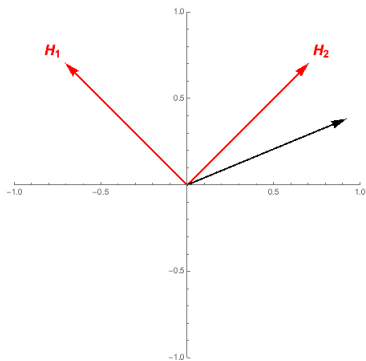
► Fundamental relation: $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$ (check)

► $Q = \begin{bmatrix} 6 & -2 \\ -2 & 6 \end{bmatrix} \succ 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} 8 \\ 4 \end{bmatrix}$



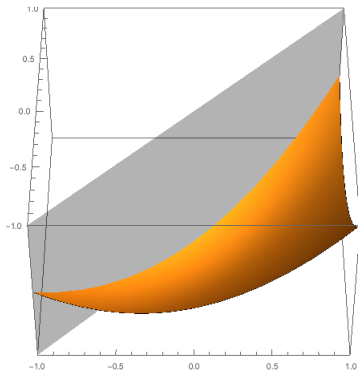
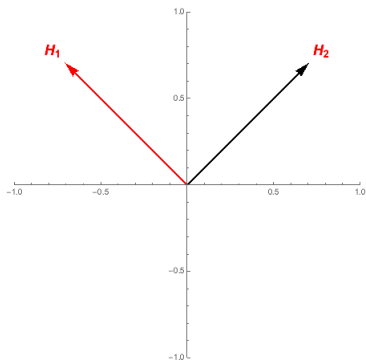
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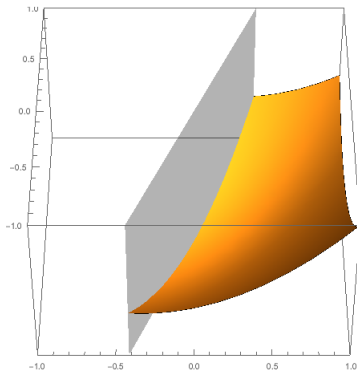
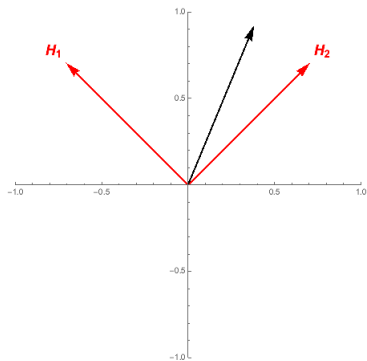
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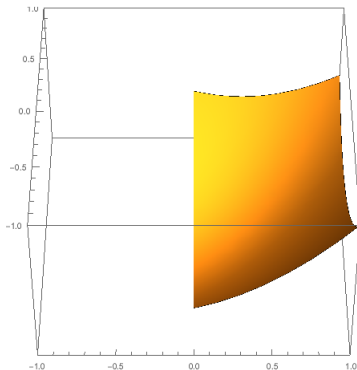
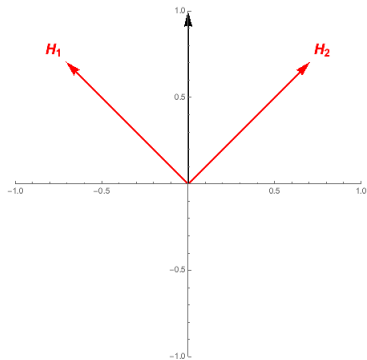
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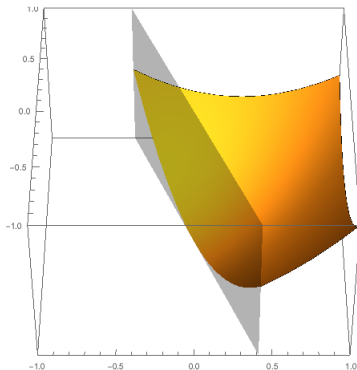
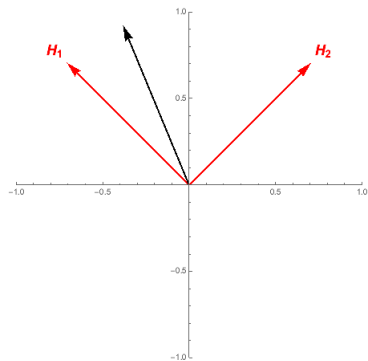
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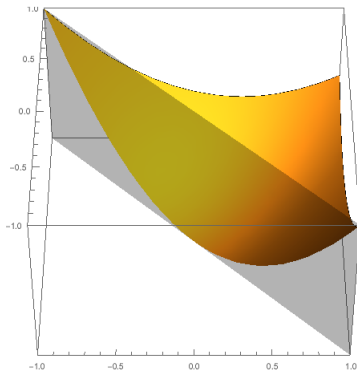
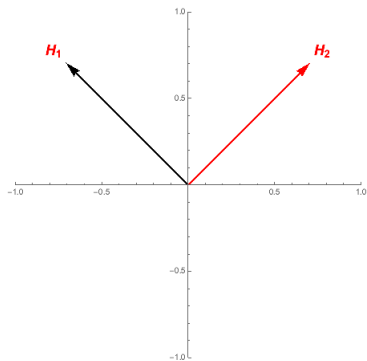
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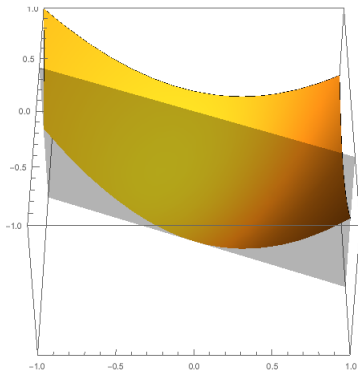
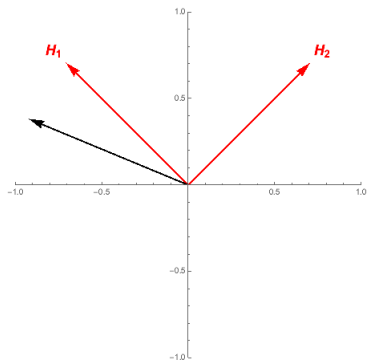
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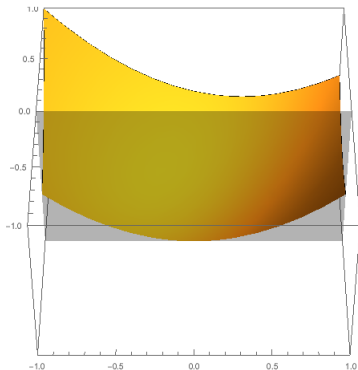
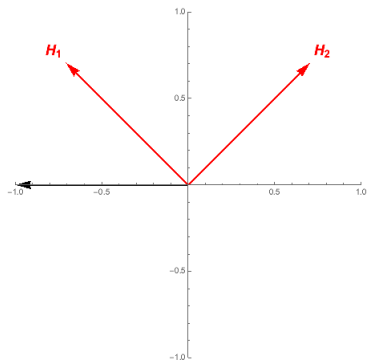
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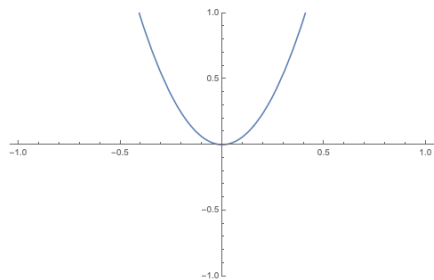
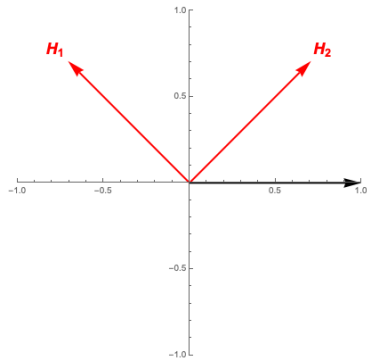
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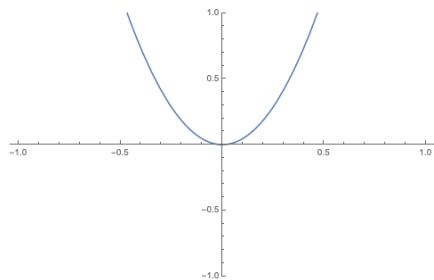
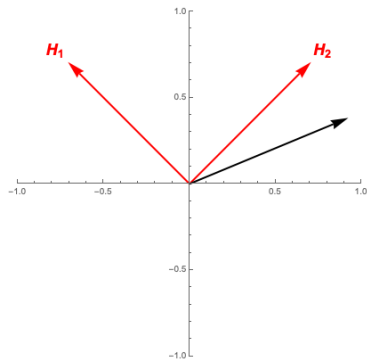


► $d^T Q d > 0 \forall d$, steepness change with d

►

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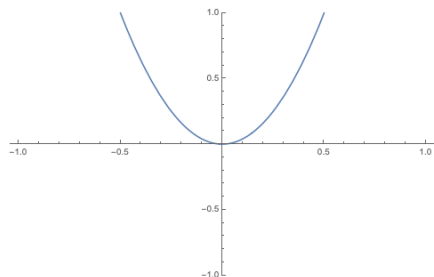
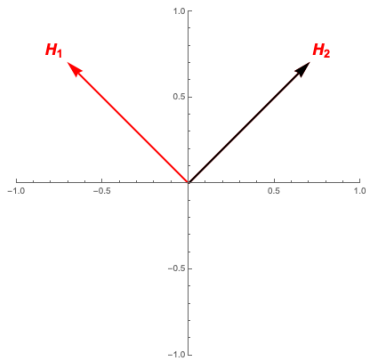


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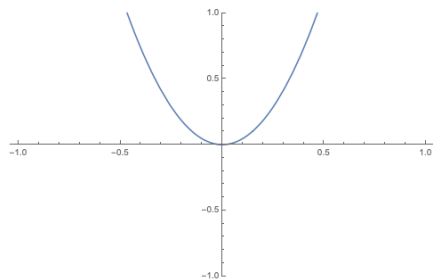
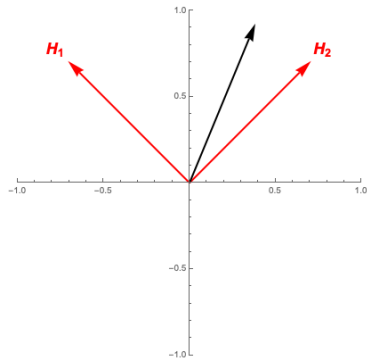


► $d^T Q d > 0 \forall d$, steepness change with d

► least steep along H_2 ($\lambda_2 = 4$)

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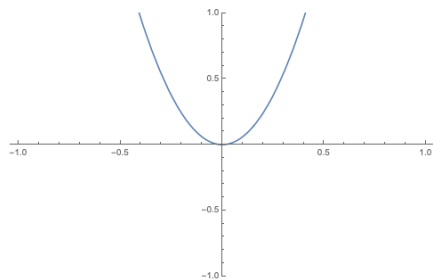
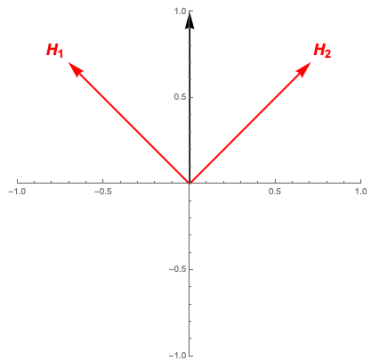


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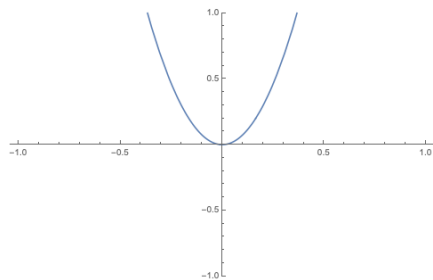
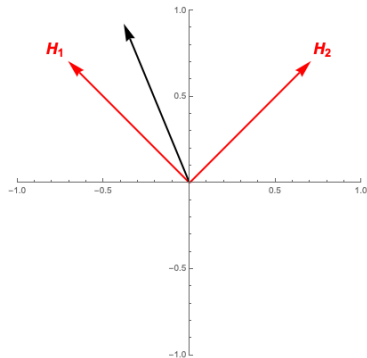


► $d^T Q d > 0 \forall d$, steepness change with d

►

► Fundamental relation: $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$ (check)

► $Q = \begin{bmatrix} 6 & -2 \\ -2 & 6 \end{bmatrix} \succ 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} 8 \\ 4 \end{bmatrix}$

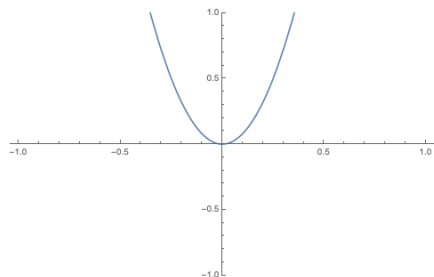
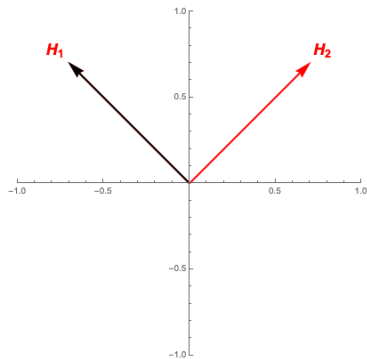


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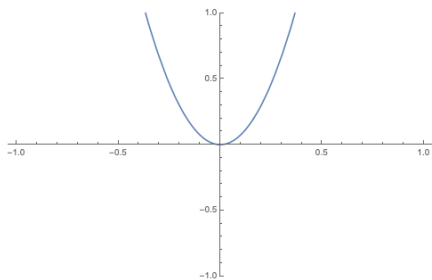
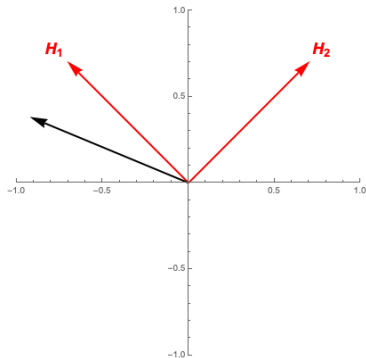


► $d^T Q d > 0 \forall d$, steepness change with d

► steepest along H_1 ($\lambda_1 = 8$)

► Fundamental relation: $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$ (check)

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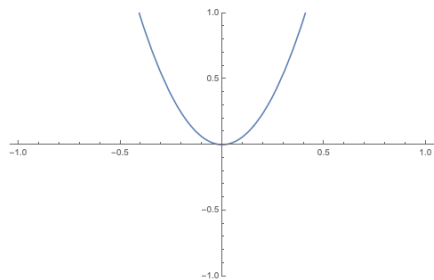
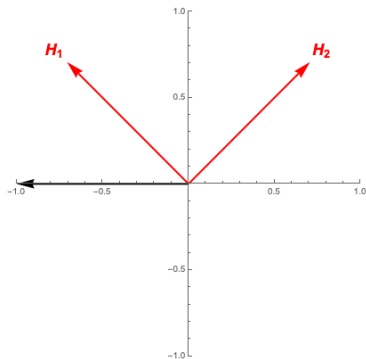


► $d^T Q d > 0 \forall d$, steepness change with d

► intermediate steepness “in between”

► Fundamental relation: $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$ (check)

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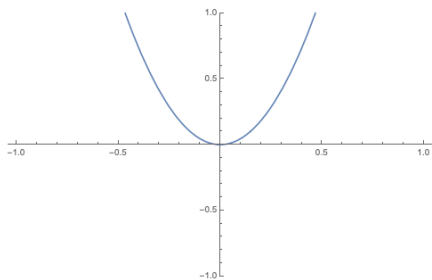
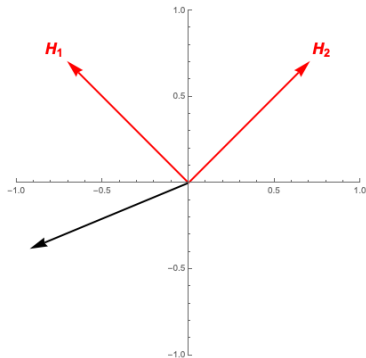


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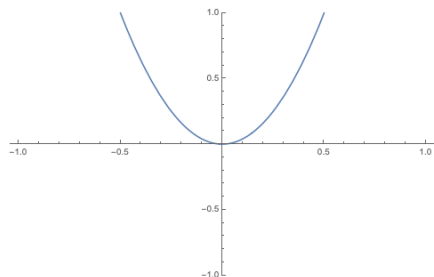
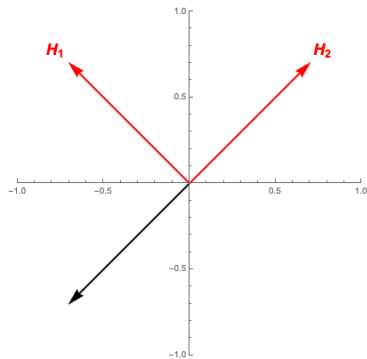


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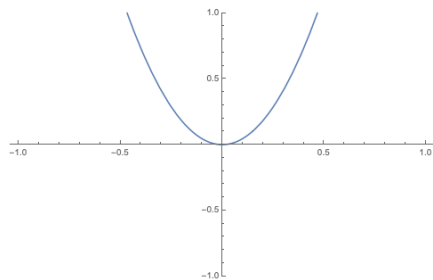
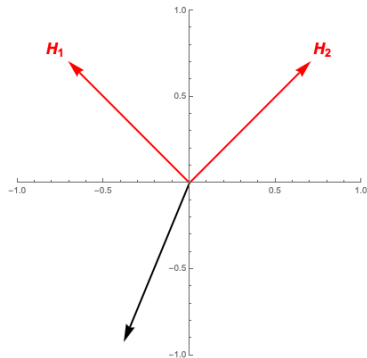


► $d^T Q d > 0 \forall d$, steepness change with d

► least steep along $-H_2$ ($\lambda_2 = 4$)

► Fundamental relation: $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$ (check)

► $Q = \begin{bmatrix} 6 & -2 \\ -2 & 6 \end{bmatrix} \succ 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} 8 \\ 4 \end{bmatrix}$

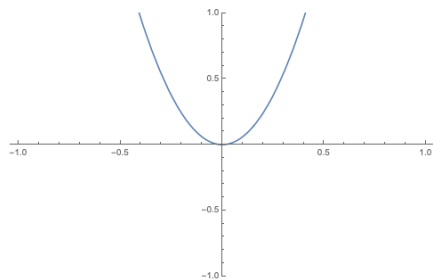
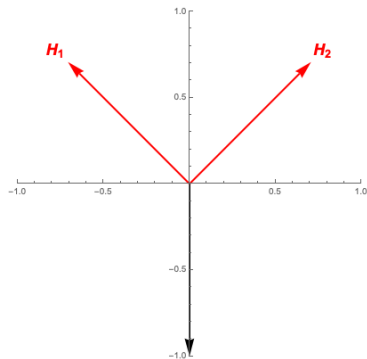


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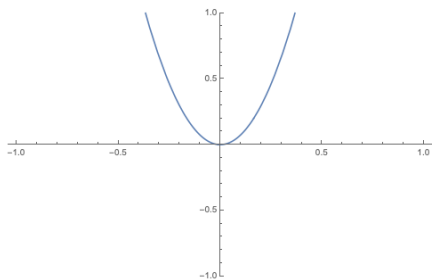
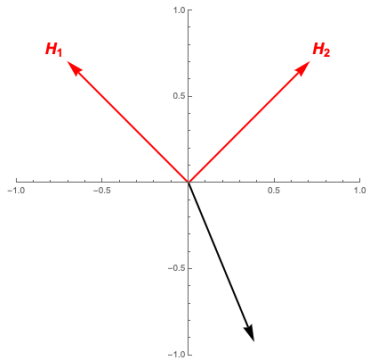


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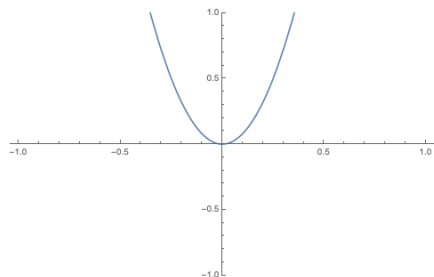
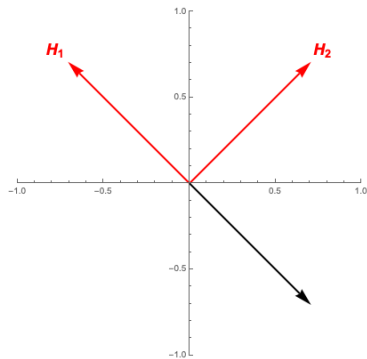


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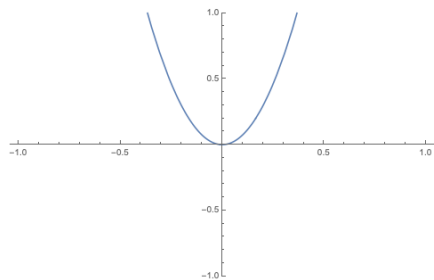
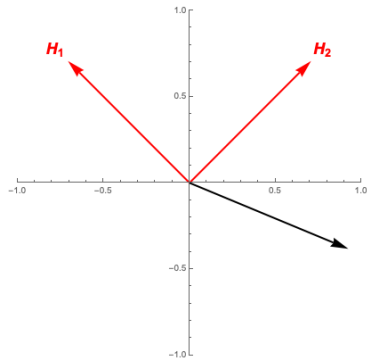


► $d^T Q d > 0 \forall d$, steepness change with d

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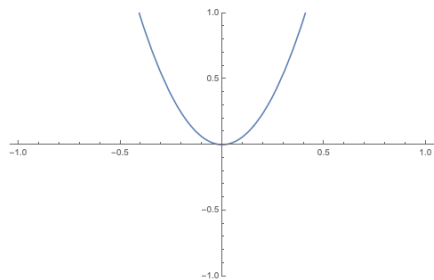
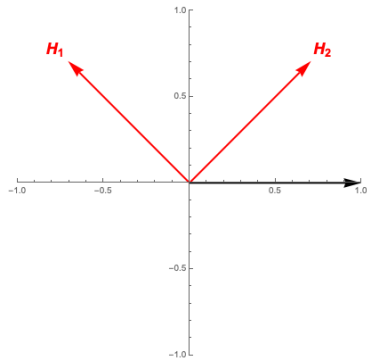


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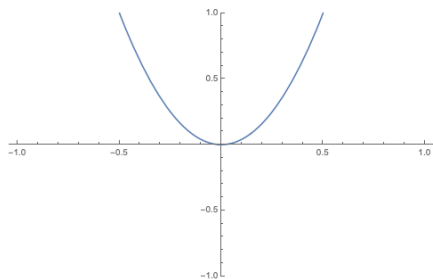
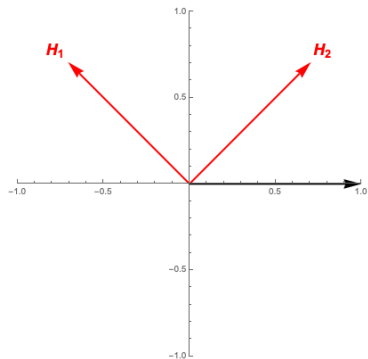


► $d^T Q d > 0 \forall d$, steepness change with d

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► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

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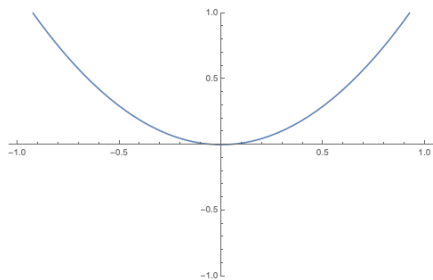
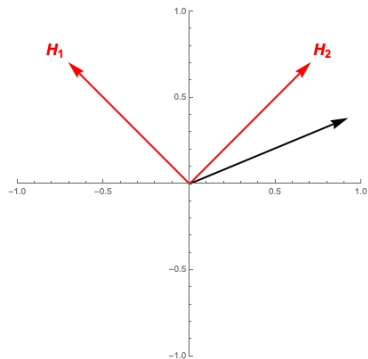


► $d^T Q d \geq 0 \forall d$, but $\exists d$ s.t. $d^T Q d = 0$

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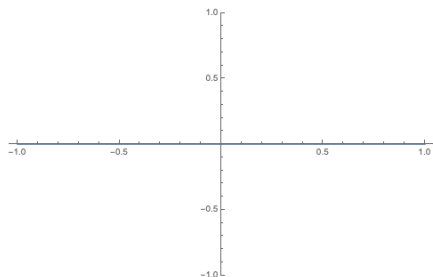
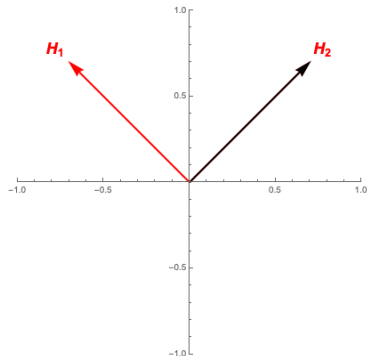


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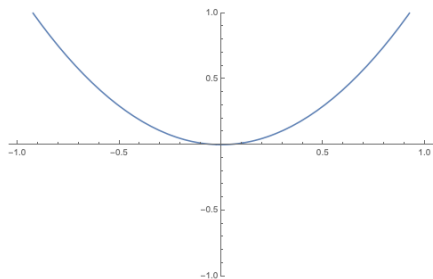
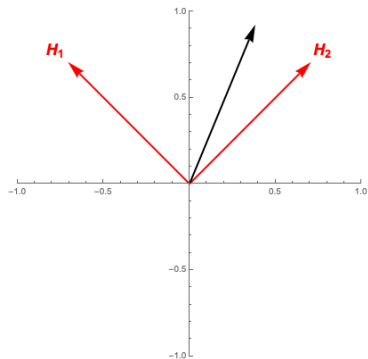


► $d^T Q d \geq 0 \forall d$, but $\exists d$ s.t. $d^T Q d = 0$

► completely flat along H_2 ($\lambda_2 = 0$)

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

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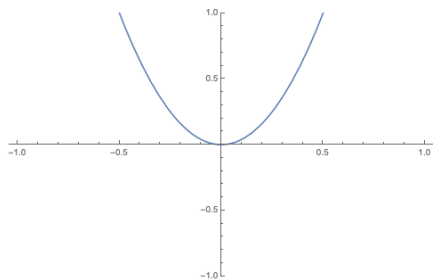
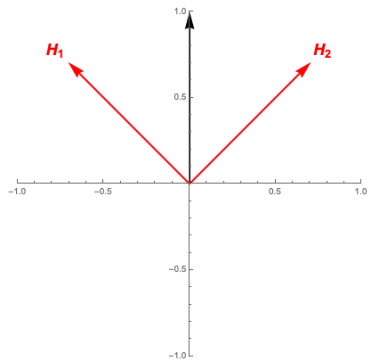


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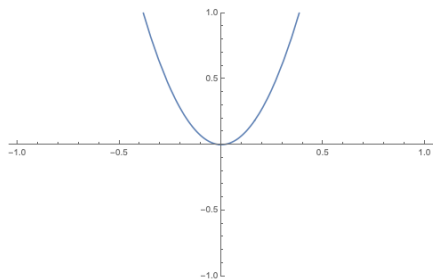
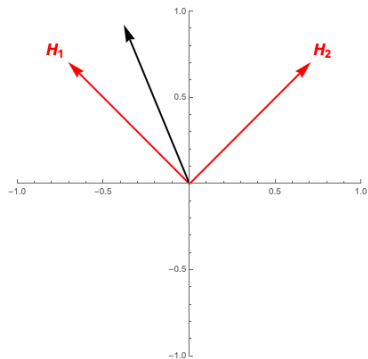


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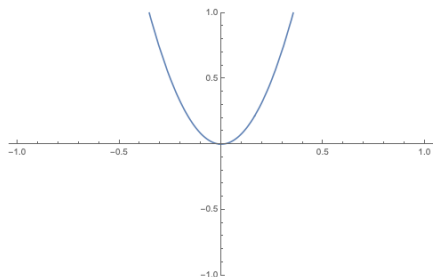
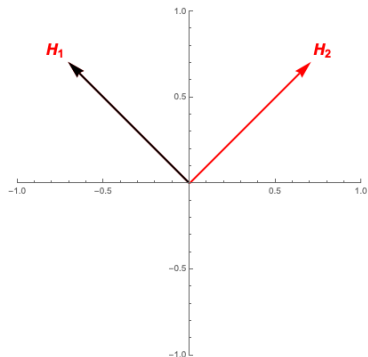


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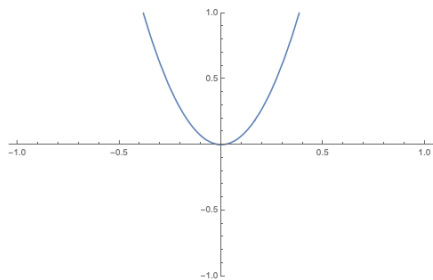
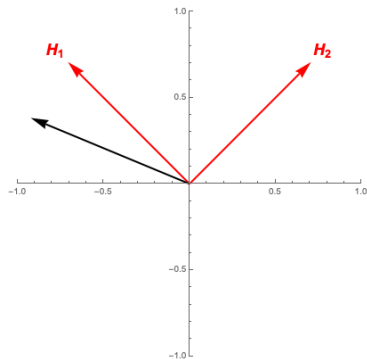


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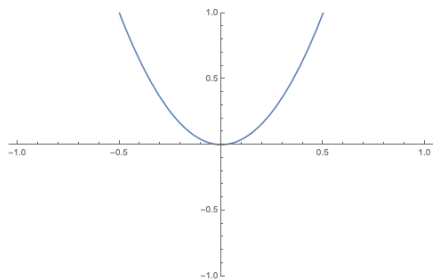
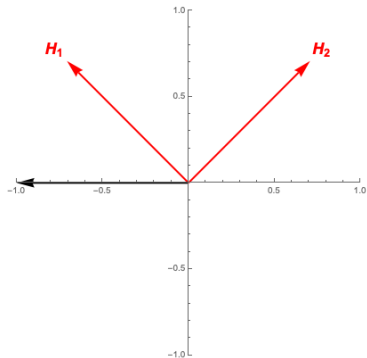


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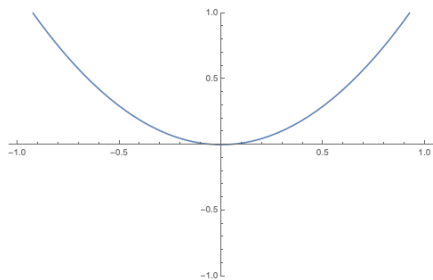
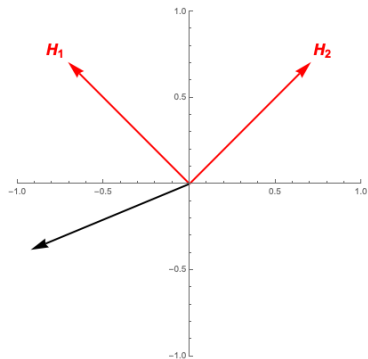


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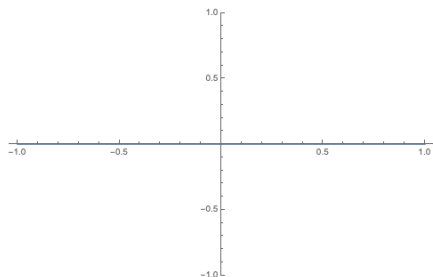
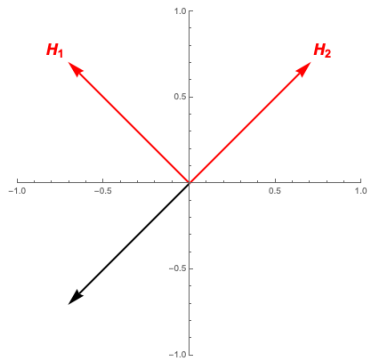


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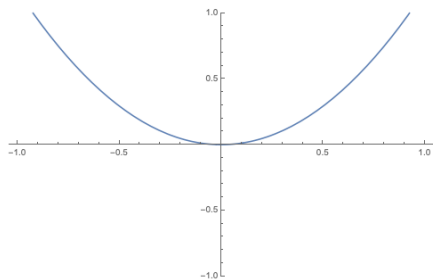
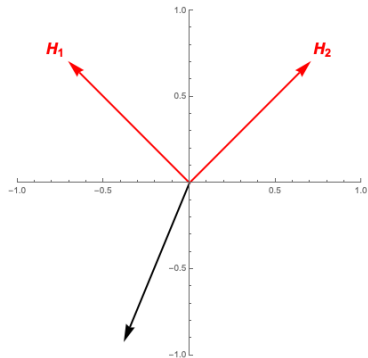


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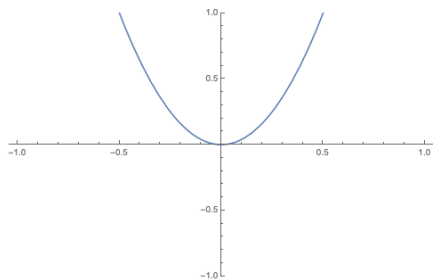
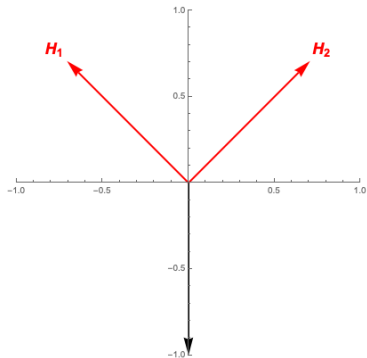


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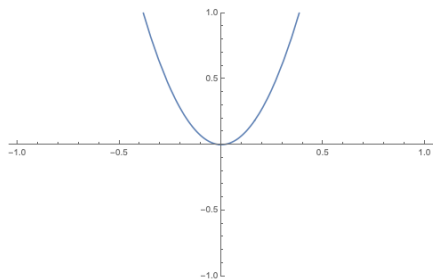
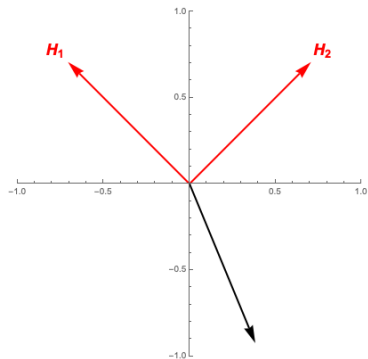


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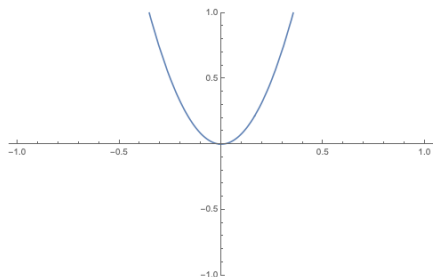
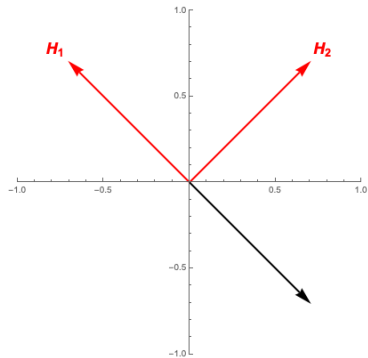


► $d^T Q d \geq 0 \forall d$, but $\exists d$ s.t. $d^T Q d = 0$

► intermediate steepness “in between”

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

► $Q = \begin{bmatrix} 4 & -4 \\ -4 & 4 \end{bmatrix} \succeq 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} 8 \\ 0 \end{bmatrix}$

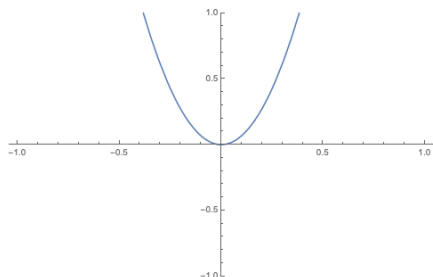
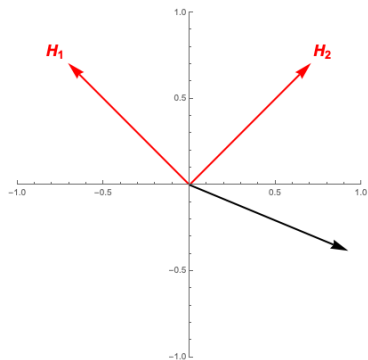


► $d^T Q d \geq 0 \forall d$, but $\exists d$ s.t. $d^T Q d = 0$

► steepest along $-H_1$ ($\lambda_1 = 8$)

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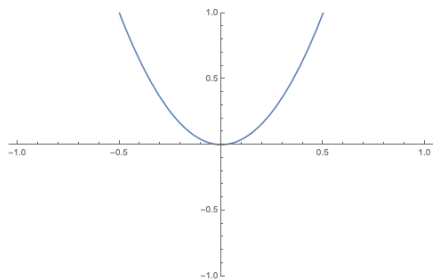
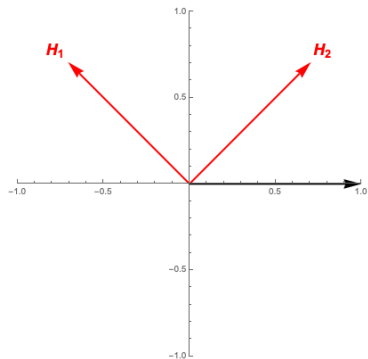


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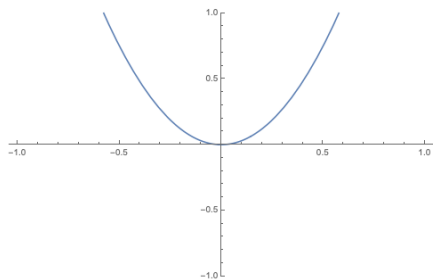
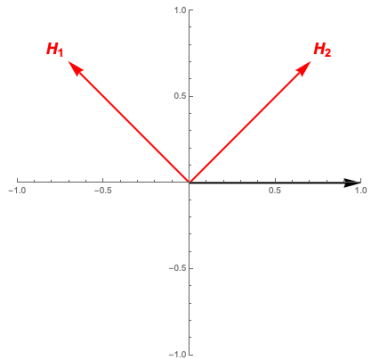


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► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

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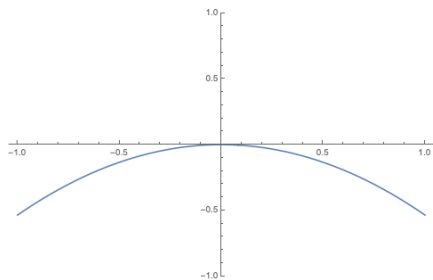
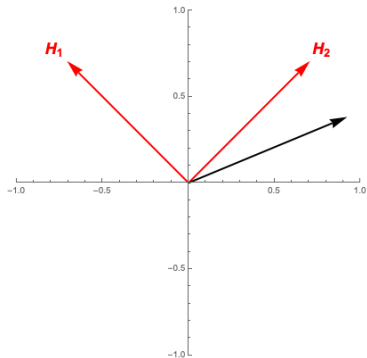


► $d^T Q d$ can be both > 0 and < 0

►

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

► $Q = \begin{bmatrix} 3 & -5 \\ -5 & 3 \end{bmatrix} \succ 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} 8 \\ -2 \end{bmatrix}$

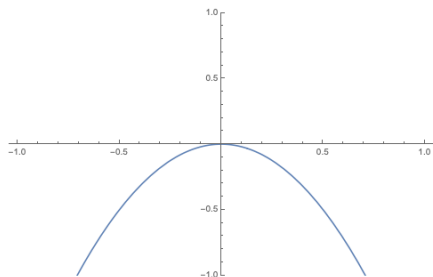
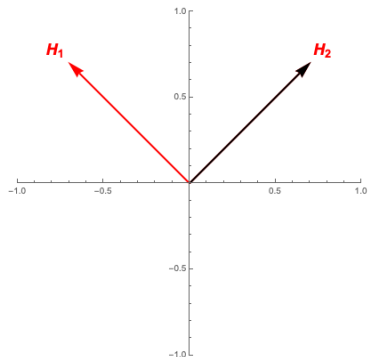


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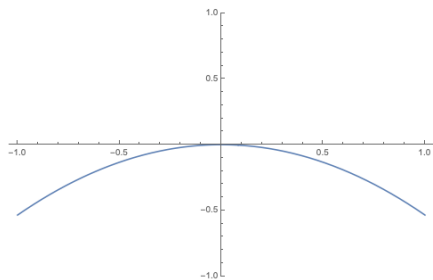
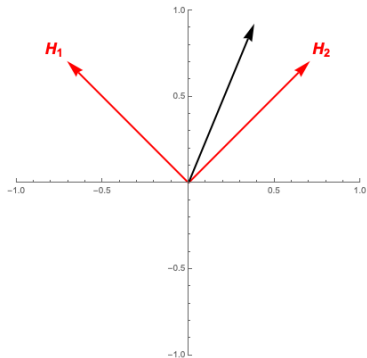


► $d^T Q d$ can be both > 0 and < 0

► steepest **negative** along H_2 ($\lambda_2 = -2$)

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

► $Q = \begin{bmatrix} 3 & -5 \\ -5 & 3 \end{bmatrix} \succ 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} 8 \\ -2 \end{bmatrix}$

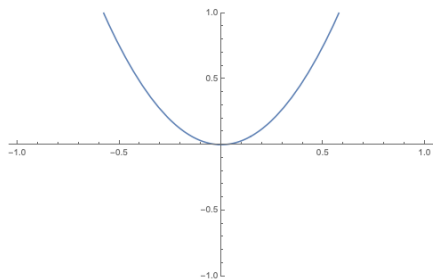
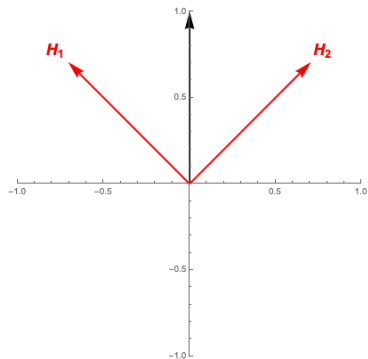


► $d^T Q d$ can be both > 0 and < 0

►

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

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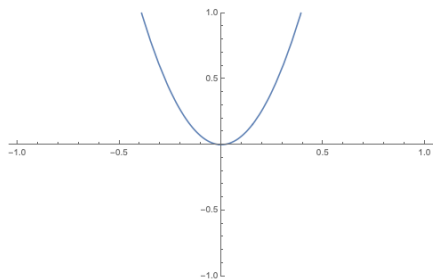
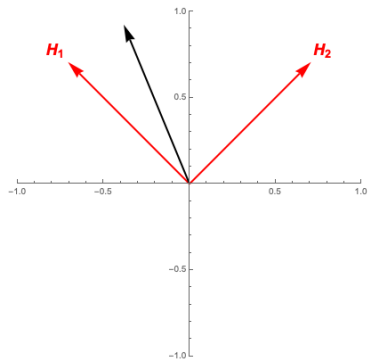


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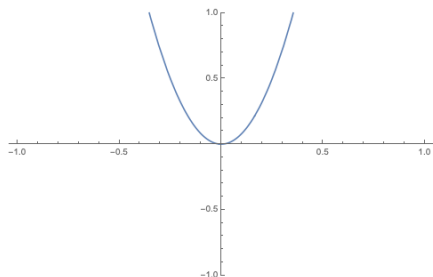
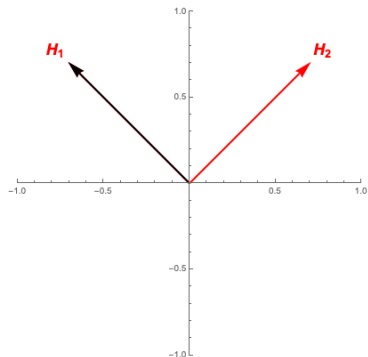


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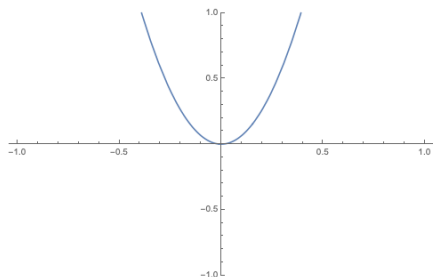
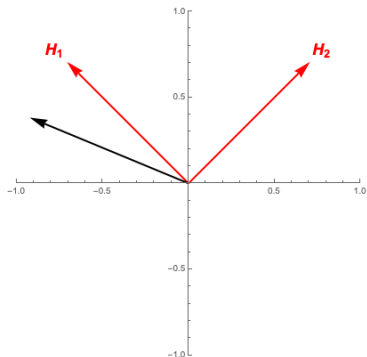


► $d^T Q d$ can be both > 0 and < 0

► steepest **positive** along H_1 ($\lambda_1 = 8$)

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

► $Q = \begin{bmatrix} 3 & -5 \\ -5 & 3 \end{bmatrix} \succ 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} 8 \\ -2 \end{bmatrix}$

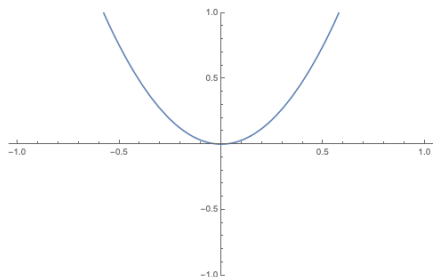
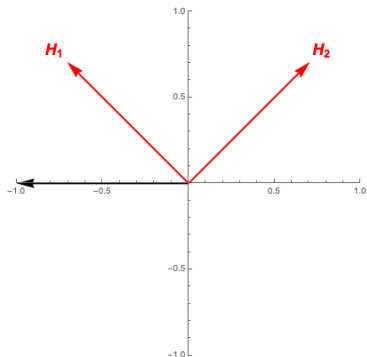


► $d^T Q d$ can be both > 0 and < 0

► intermediate steepness (positive or negative) “in between”

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

► $Q = \begin{bmatrix} 3 & -5 \\ -5 & 3 \end{bmatrix} \succ 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} 8 \\ -2 \end{bmatrix}$

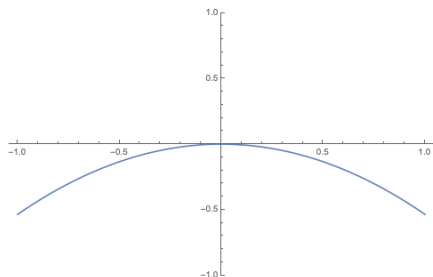
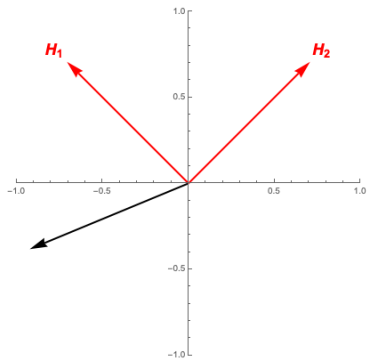


► $d^T Q d$ can be both > 0 and < 0

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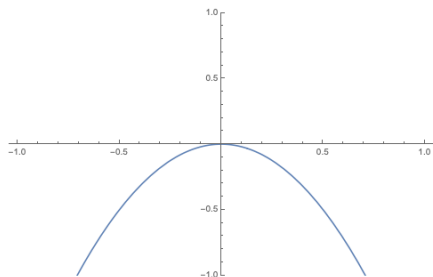
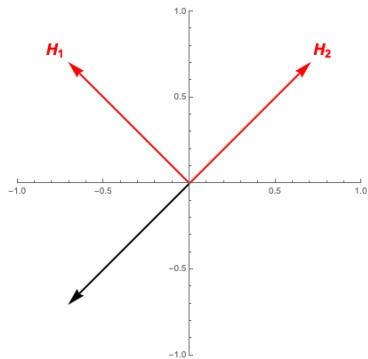


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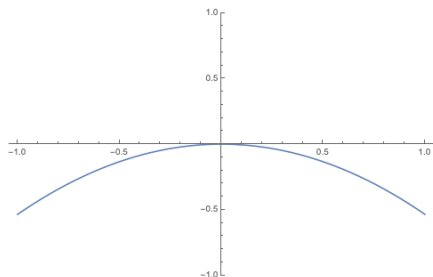
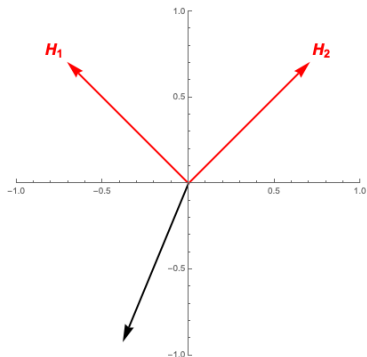


► $d^T Q d$ can be both > 0 and < 0

► steepest **negative** along $-H_2$ ($\lambda_2 = -2$)

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

► $Q = \begin{bmatrix} 3 & -5 \\ -5 & 3 \end{bmatrix} \succ 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} 8 \\ -2 \end{bmatrix}$

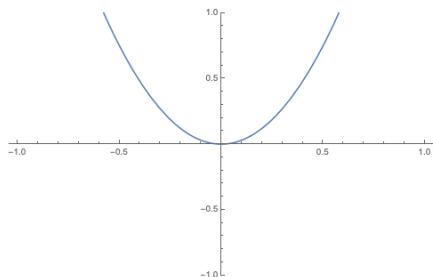
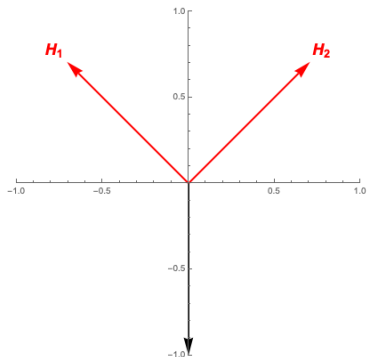


► $d^T Q d$ can be both > 0 and < 0

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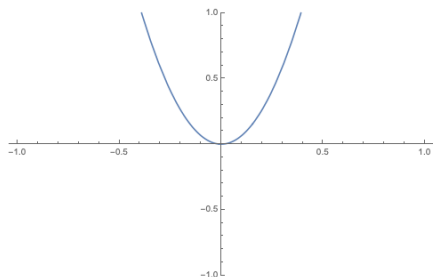
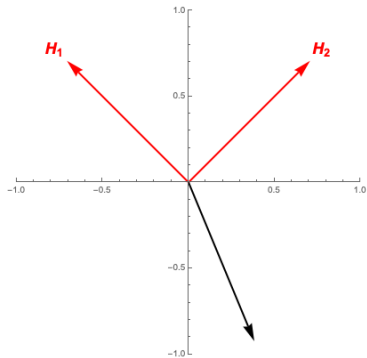


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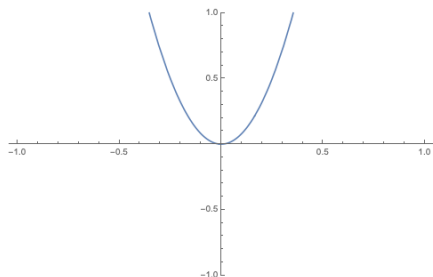
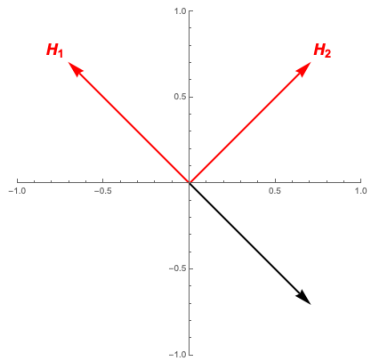


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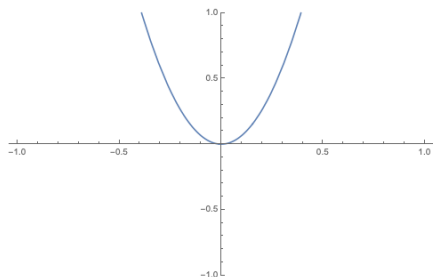
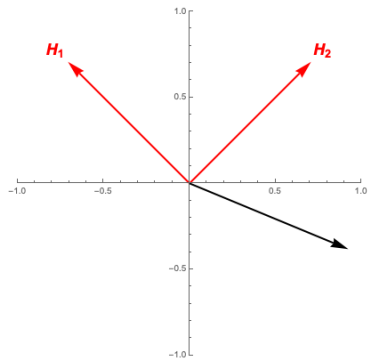


► $d^T Q d$ can be both > 0 and < 0

► steepest **positive** along $-H_1$ ($\lambda_1 = 8$)

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

► $Q = \begin{bmatrix} 3 & -5 \\ -5 & 3 \end{bmatrix} \succ 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} 8 \\ -2 \end{bmatrix}$

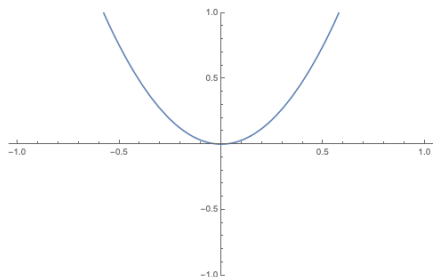
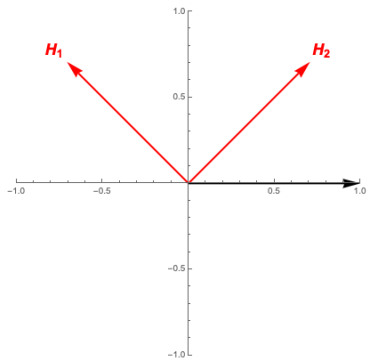


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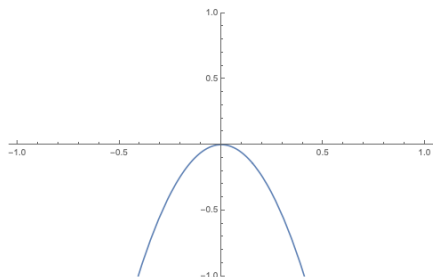
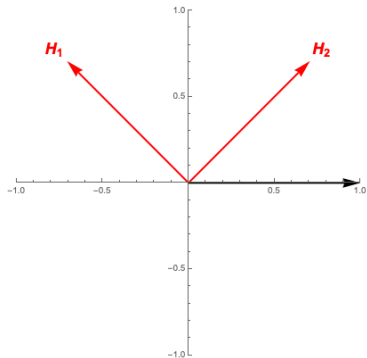


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► intermediate steepness (positive or negative) “in between”

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

► $Q = \begin{bmatrix} -6 & -2 \\ -2 & -6 \end{bmatrix} \prec 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} -4 \\ -8 \end{bmatrix}$

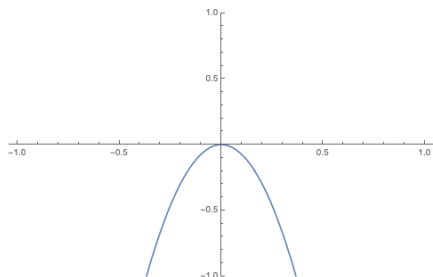
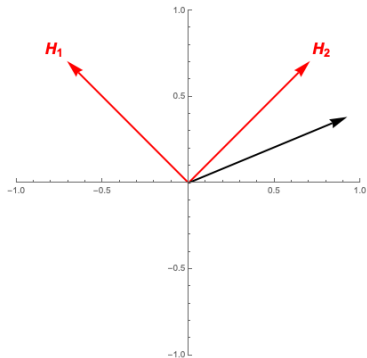


► $d^T Q d < 0 \forall d$, steepness change with d

►

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

► $Q = \begin{bmatrix} -6 & -2 \\ -2 & -6 \end{bmatrix} \prec 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} -4 \\ -8 \end{bmatrix}$

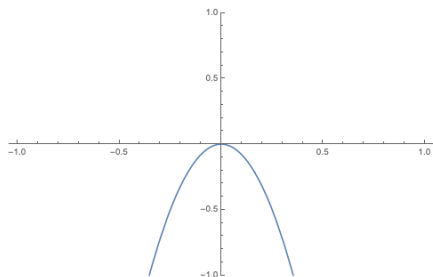
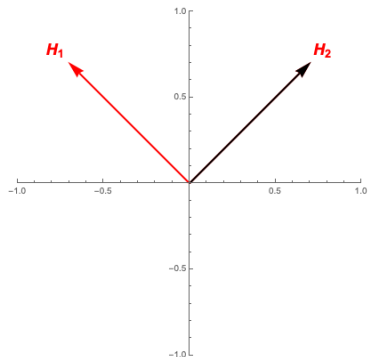


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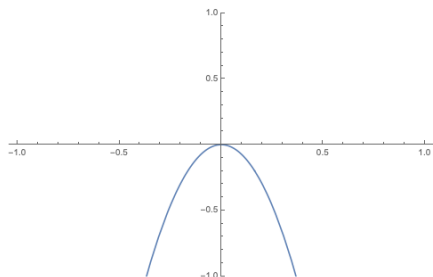
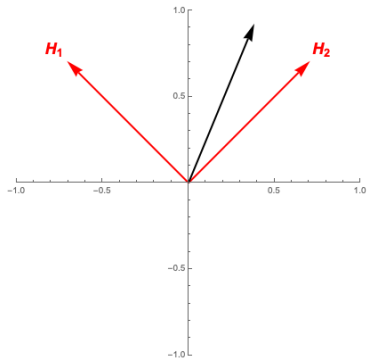


► $d^T Q d < 0 \forall d$, **steepness change with d**

► steepest negative along H_2 ($\lambda_2 = -8$)

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

► $Q = \begin{bmatrix} -6 & -2 \\ -2 & -6 \end{bmatrix} \prec 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} -4 \\ -8 \end{bmatrix}$

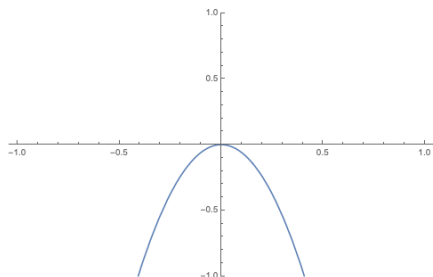
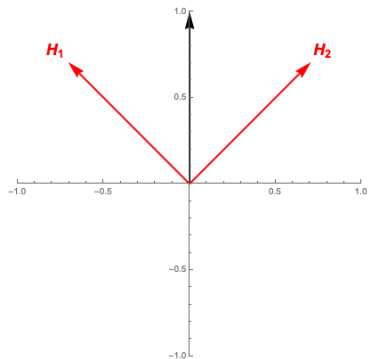


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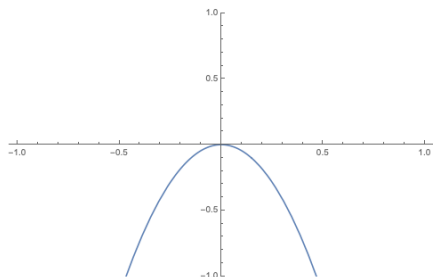
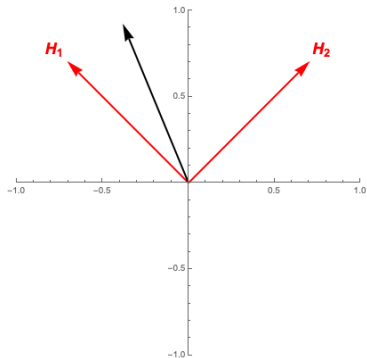


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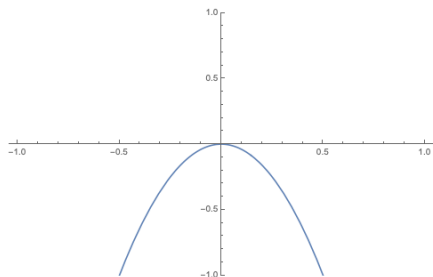
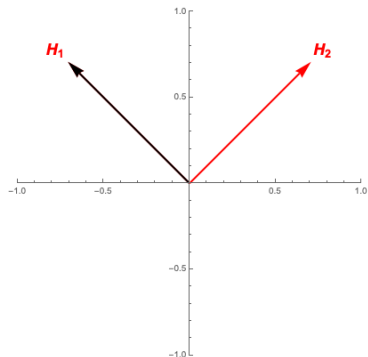


► $d^T Q d < 0 \forall d$, steepness change with d

►

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

► $Q = \begin{bmatrix} -6 & -2 \\ -2 & -6 \end{bmatrix} \prec 0 \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix} \quad \lambda = \begin{bmatrix} -4 \\ -8 \end{bmatrix}$

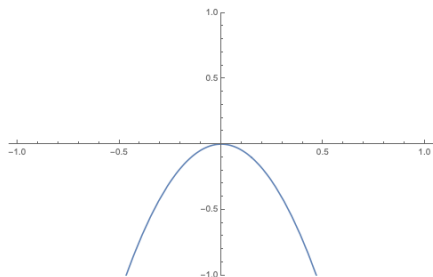
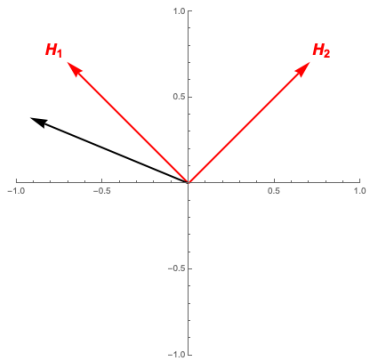


► $d^T Q d < 0 \forall d$, steepness change with d

► least steep negative along H_1 ($\lambda_1 = -4$)

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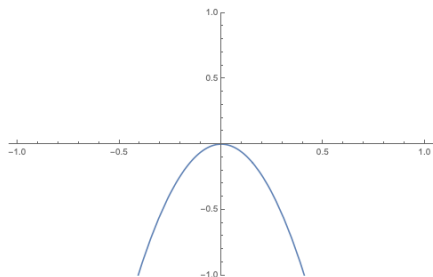
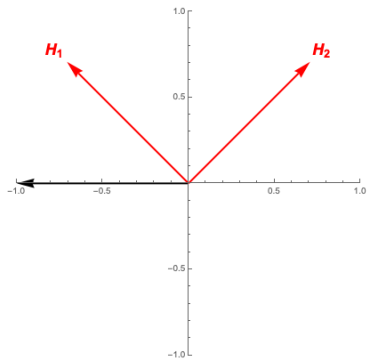


► $d^T Q d < 0 \forall d$, steepness change with d

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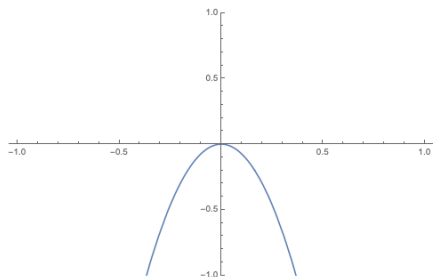
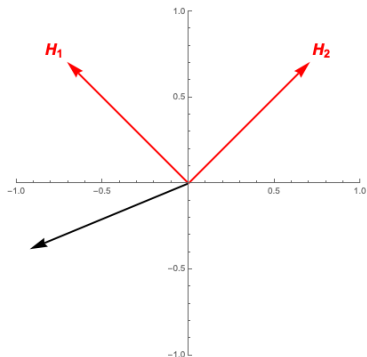


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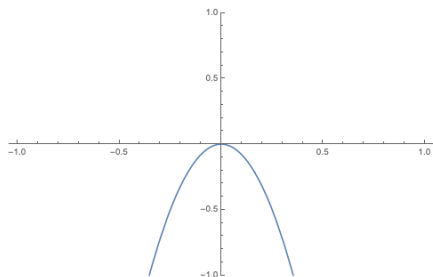
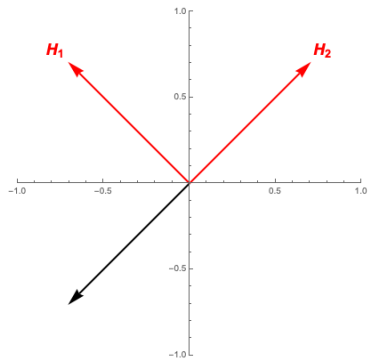


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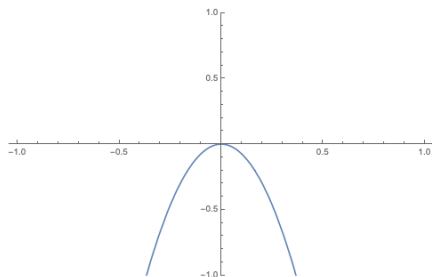
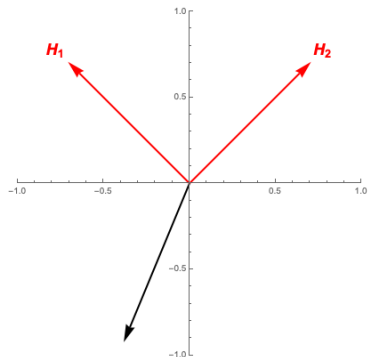


► $d^T Q d < 0 \forall d$, **steepness change with d**

► steepest negative along $-H_2$ ($\lambda_2 = -8$)

► Recall $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

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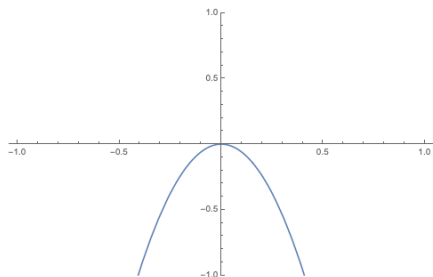
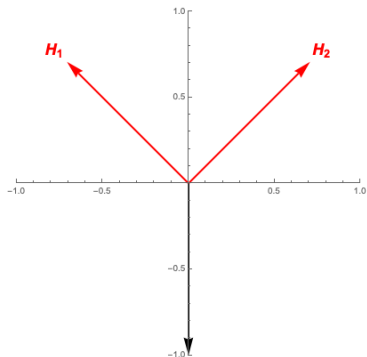


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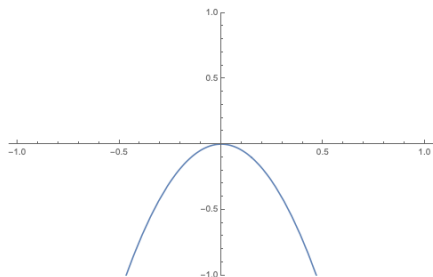
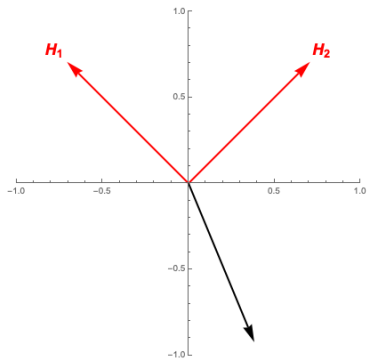


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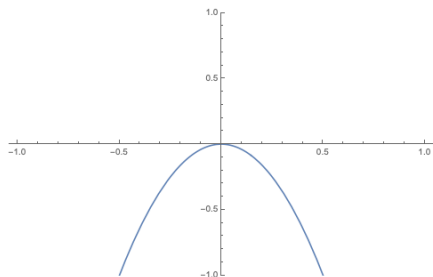
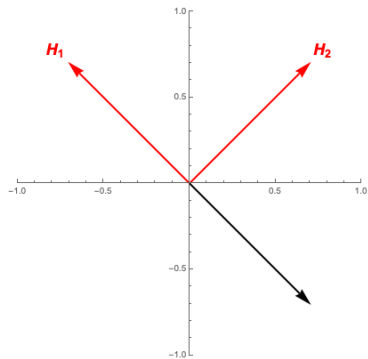


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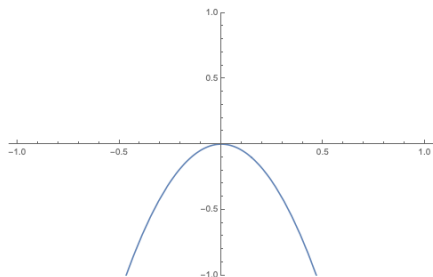
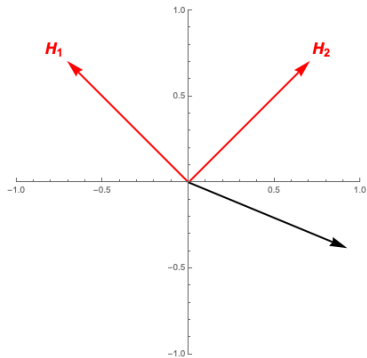


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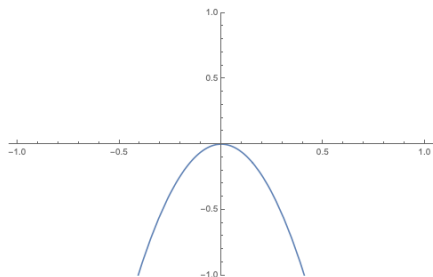
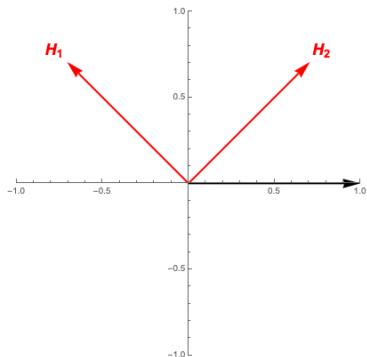


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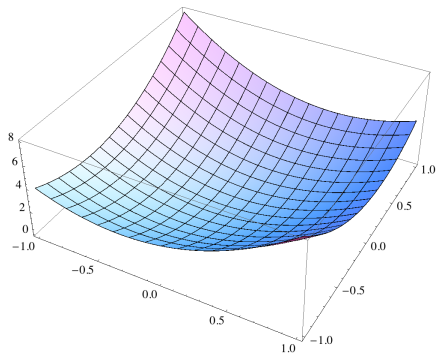


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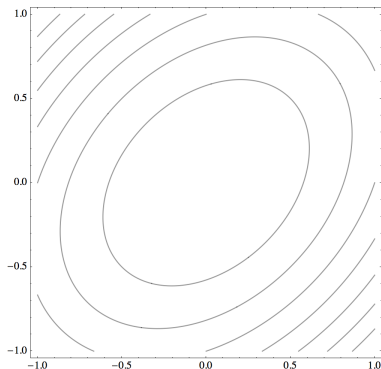
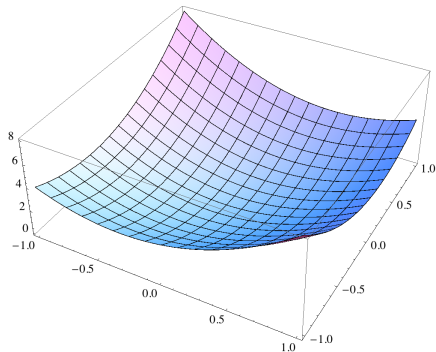
- ▶ All level sets centred in $x = 0$ by symmetry

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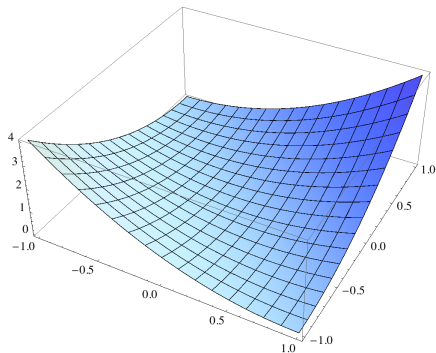
- ▶ $Q = \begin{bmatrix} 6 & -2 \\ -2 & 6 \end{bmatrix} \succ 0$ graph is a (convex) paraboloid

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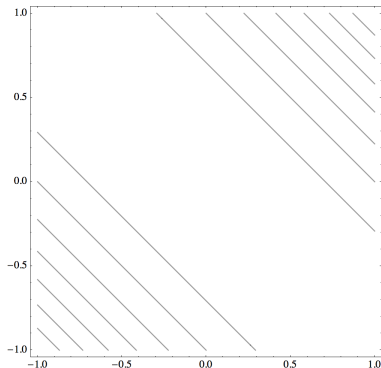
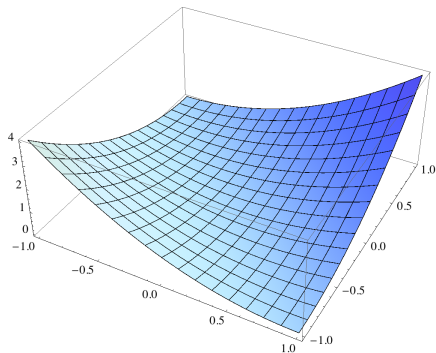
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level sets are **ellipsoids**

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- ▶ $Q = \begin{bmatrix} 2 & 2 \\ 2 & 2 \end{bmatrix} \preceq 0$ graph is a degenerate paraboloid

- All level sets centred in $x = 0$ by symmetry

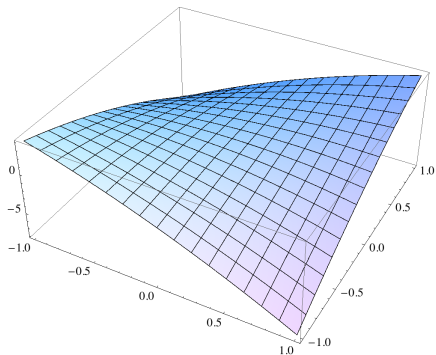


► $Q = \begin{bmatrix} 2 & 2 \\ 2 & 2 \end{bmatrix} \succcurlyeq 0$

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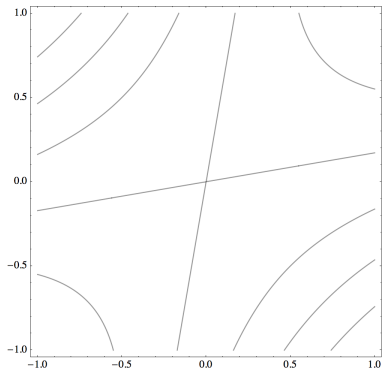
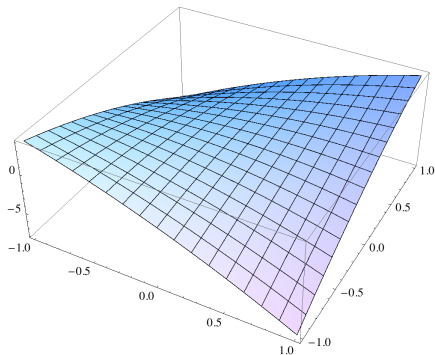
level sets are **degenerate ellipsoids**

- ▶ All level sets centred in $x = 0$ by symmetry



- ▶ $Q = \begin{bmatrix} 3 & -5 \\ -5 & -3 \end{bmatrix} \succ 0$ graph saddle-shaped (0 is a saddle point)

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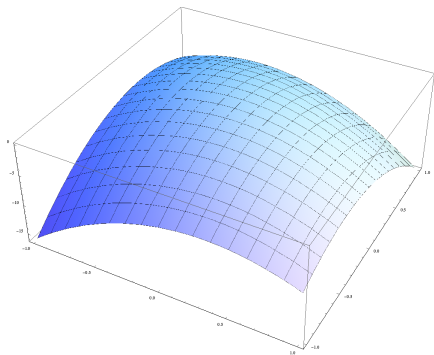


► $Q = \begin{bmatrix} 3 & -5 \\ -5 & -3 \end{bmatrix} \chi^2 = 0$

graph saddle-shaped (0 is a saddle point)

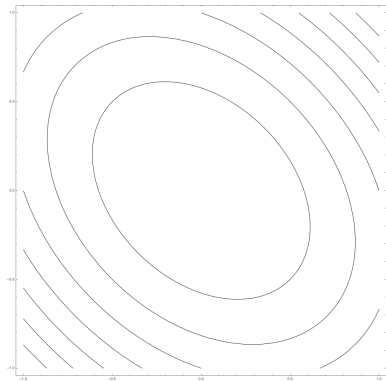
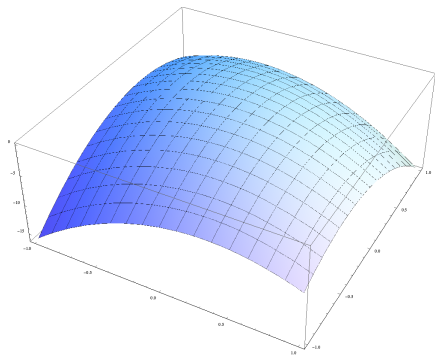
level sets are **hyperboloids**

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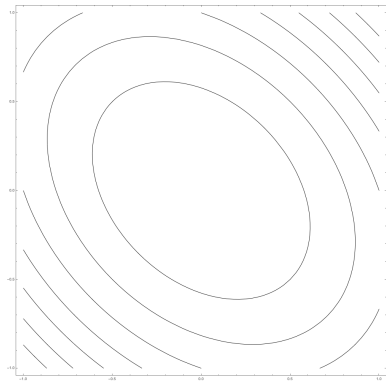
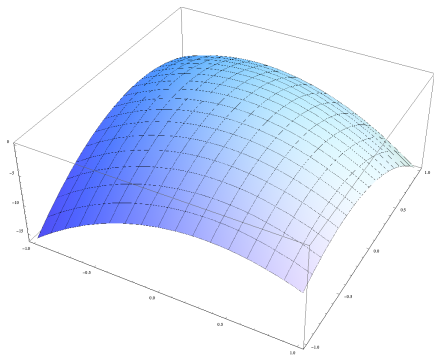
- ▶ $Q = \begin{bmatrix} -6 & -2 \\ -2 & -6 \end{bmatrix} \prec 0$ graph a (concave, i.e., “upside-down”) paraboloid

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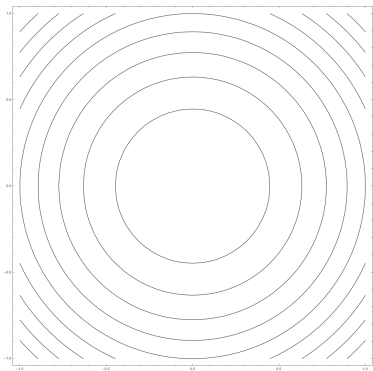


- $Q = \begin{bmatrix} -6 & -2 \\ -2 & -6 \end{bmatrix} \prec 0$ graph a (concave, i.e., “upside-down”) paraboloid
level sets are **ellipsoids** again

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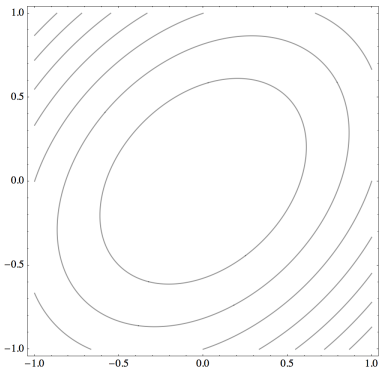


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level sets are **ellipsoids** again
- ▶ Level sets can be precisely described in terms of H_i, λ_i



► $\|x\|_2^2 \equiv Q = H = \Lambda = I$: perfect circles

$$Q = H = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}, \quad \lambda = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

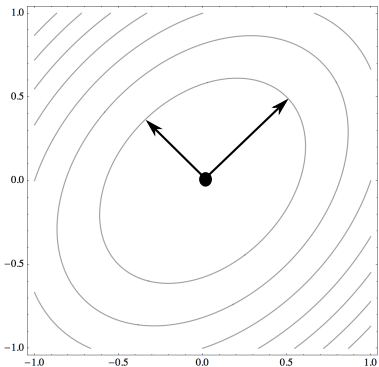


▶ Recall again $\varphi_{H_i}(\alpha) = \alpha^2 \lambda_i$

▶ $L(f, 1) \cap H_i \equiv \varphi_{H_i}(\alpha) = 1 \implies \lambda_i > 0$

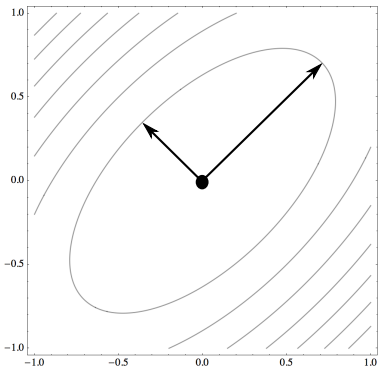
▶ $\varphi_{H_i}(\alpha) = 1 \equiv \alpha = \sqrt{1/\lambda_i} \implies$

$$Q = \begin{bmatrix} 6 & -2 \\ -2 & 6 \end{bmatrix}, \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix}, \quad \lambda = \begin{bmatrix} 8 \\ 4 \end{bmatrix}$$



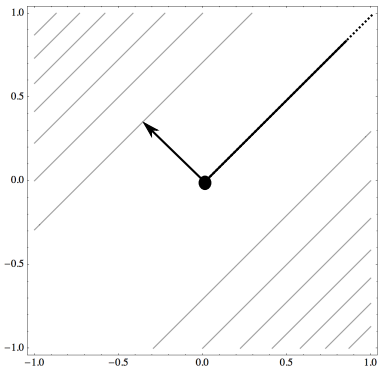
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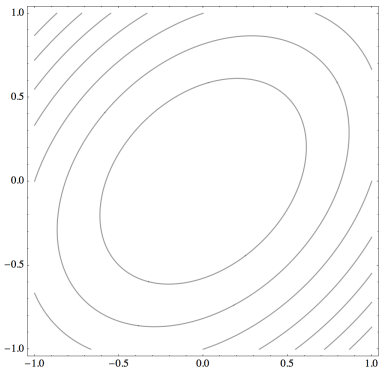
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$$Q = \begin{bmatrix} 5 & -3 \\ -3 & 5 \end{bmatrix}, \quad H = \frac{\sqrt{2}}{2} \begin{bmatrix} -1 & 1 \\ 1 & 1 \end{bmatrix}, \quad \lambda = \begin{bmatrix} 8 \\ 2 \end{bmatrix}$$



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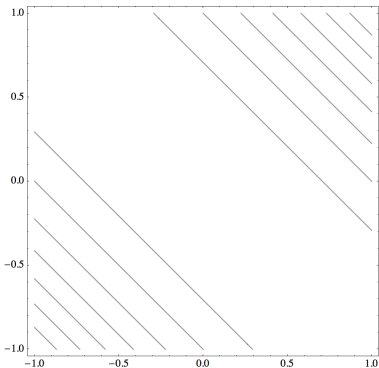
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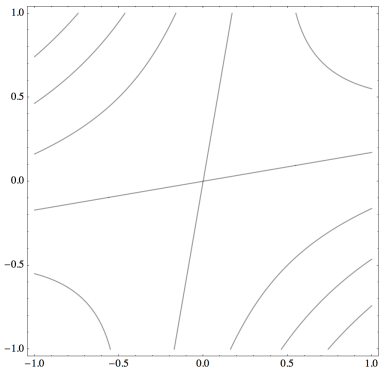
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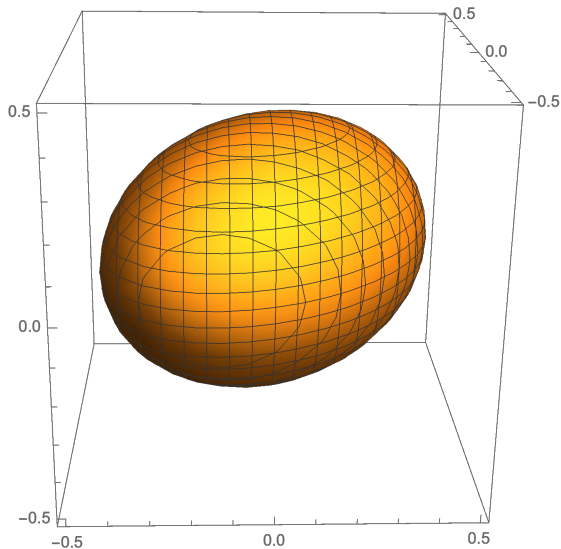
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- ▶ All λ_i have the same sign: $f(x)$ either ≥ 0 or $\leq 0 \implies$ ellipsoids
- ▶ Some $\lambda_i = 0 \implies$ “degenerate” ellipsoids (∞ axis)
- ▶ $\lambda_i > 0$ and $\lambda_j < 0$: $\exists \alpha_i, \alpha_j$ s.t. $\varphi_{H_i}(\alpha_i) + \varphi_{H_j}(\alpha_j) = 0 \implies$ hyperboloids

Level sets homogeneous quadratic functions, 3D example

37

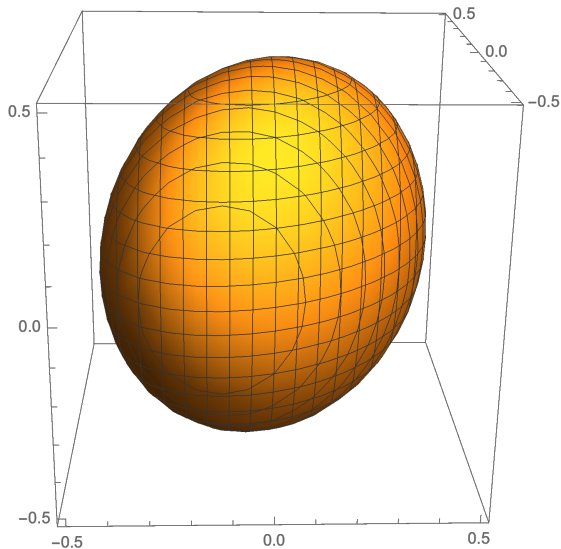
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 $L(f, 1)$ 

Level sets homogeneous quadratic functions, 3D example

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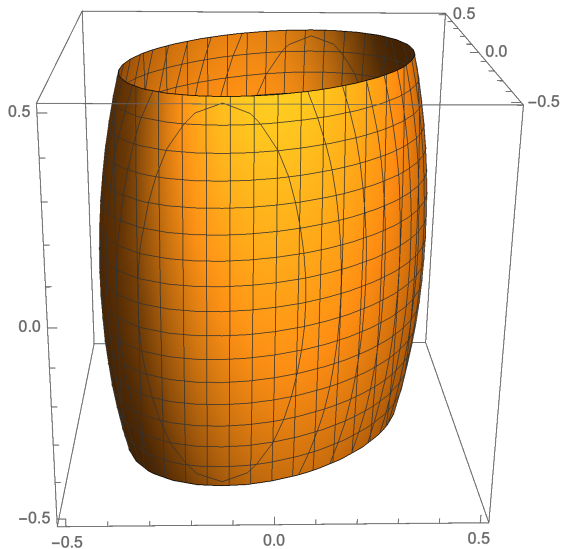
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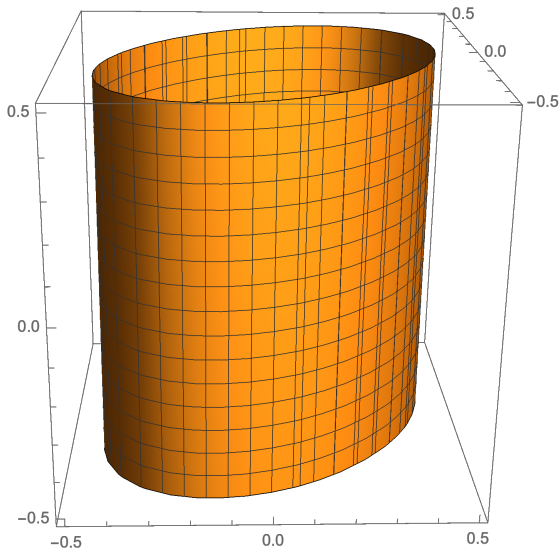
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- ▶ Clearly depends **sign of eigenvalues of $Q \equiv$ definiteness**:
 - ▶ $Q \succeq 0 \wedge Q \preceq 0 \equiv \lambda_1 = \lambda_n = 0 \equiv Q = 0 \implies \min = \max = 0$ (constant)
 - ▶ $Q \succeq 0 \implies \min = 0, \operatorname{argmin} = 0, \max = +\infty$
 - ▶ $Q \preceq 0 \implies \max = 0, \operatorname{argmax} = 0, \min = -\infty$
 - ▶ $Q \not\preceq 0 \implies \max = +\infty, \min = -\infty$

analogous to univariate case, but “many more ways to be $> 0 / < 0$ ”

Exercise: Formally prove all the unboundedness results

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Exercise: Formally prove all the unboundedness results

- ▶ Box-constrained optimization on (closed) hyperrectangle X absolutely **not** analogous to the univariate case:
 - ▶ \mathcal{NP} -hard in most cases [3]
 - ▶ min with $Q \succeq 0$ and max with $Q \preceq 0$ **polynomial** but **nontrivial** (will see)
- ▶ \mathcal{NP} -hardness due to \mathbb{R}^n “big” (X has 2^n vertices), issue also in \mathcal{P} case
- ▶ $\max\{f(x)\}$ and $\min\{f(x)\}$ **very very different**

- ▶ $f(x) = \frac{1}{2}x^T Qx + \langle q, x \rangle$: a homogeneous quadratic plus a linear
- ▶ $q \neq 0$ but Q nonsingular $\equiv \lambda_i \neq 0 \forall i$ (regardless of the sign)
- ▶ Then $f(x) = g(z) = \frac{1}{2}z^T Qz + f(\bar{x})$ for $z = x - \bar{x}$ and $\bar{x} = -Q^{-1}q$

Exercise: Prove the result, but it should look familiar

▶ Optimizing a quadratic non-homogeneous function

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▶ Optimizing a quadratic non-homogeneous function

- ▶ $\bar{x} (\neq 0)$ centre of the level sets: repeat
 for $g(z)$, translate the results back in x -space
- ▶ Box-constrained case remains hard / nontrivial
- ▶ Analogous to univariate case, but many more ways for (pieces of) Q to be 0 and therefore the result not be applicable
- ▶ More complicated analysis needed, coming right next

▶ Optimizing a homogeneous quadratic multivariate function

- ▶ $Q \in \mathbb{R}^{n \times n}$, eigenvalue decomposition (H, Λ) , $I = \{1, 2, \dots, n\}$
- ▶ $I^0 = \{i \in I; \lambda_i = 0\}$, $I^+ = I \setminus I^0$, **nonempty** ($k = |I^0| > 0$, $h = |I^+| > 0$)
- ▶ $\ker(Q) = \{v \in \mathbb{R}^n : \exists \eta \in \mathbb{R}^k \text{ s.t. } v = \sum_{i \in I^0} \eta_i H_i\}$
- ▶ $Qv = 0 \forall v \in \ker(Q) [\supset \{0\}]$ (check)
- ▶ $\text{im}(Q) = \{w \in \mathbb{R}^n : \exists \mu \in \mathbb{R}^h \text{ s.t. } w = \sum_{i \in I^+} \mu_i H_i\}$:
- ▶ $\forall w \in \text{im}(Q) \exists x \in \mathbb{R}^n \text{ s.t. } Qx = w$, $\text{im}(Q) = \text{im}(-Q)$

Exercise: Prove the result (recall $Q = \lambda_1 H_1 H_1^T + \dots + \lambda_n H_n H_n^T$, use [16])

- ▶ $\mathbb{R}^n = \text{im}(Q) + \ker(Q)$, $\text{im}(Q) \perp \ker(Q)$ (H is a hortonormal base of \mathbb{R}^n)
- ▶ $q = q^+ + q^0$, $q^+ \perp q^0$, with $q^0 \in \ker(Q) \equiv Qq^0 = 0$, and $q^+ \in \text{im}(Q) = \text{im}(-Q) \equiv \exists \bar{x} \text{ s.t. } (-Q)\bar{x} = q^+$
- ▶ Then $f(x) = g(z) = \frac{1}{2}z^T Qz + q^0 z + f(\bar{x})$ for $z = x - \bar{x}$

Exercise: Prove the result, but it should look very very familiar

- ▶ f is “truly quadratic” on $\text{im}(Q)$ but actually linear on $\ker(Q)$
- ▶ No surprise: $v \in \ker(Q) \implies f(v) = qv$
- ▶ Assume $Q \succeq 0$: f has minimum $\iff q^0 = 0 \equiv Q\bar{x} = -q$ has solution $\equiv q \in \text{im}(Q)$
- ▶ \bar{x} is not unique, in fact ∞ -ly many of them: “all are centres”
- ▶ \bar{x} solution $\implies \bar{x} + v$ solution $\forall v \in \ker(Q)$, all have the same objective value \equiv they are all and only the minima of f

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- ▶ $q_0 \neq 0 \equiv q \notin \text{im}(Q) \implies \min = -\infty, \max = +\infty$
- ▶ Box-constrained version \mathcal{P} (but nontrivial) if $Q \succeq 0 / Q \preceq 0$, hard otherwise

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- ▶ All in all: solving system $Q\bar{x} = -q$ (or proving no solutions) required

Outline

Optimization Problems

Optimization is difficult

Simple Functions, Univariate case

Simple Functions, Multivariate case

Multivariate Quadratic case: Gradient Method

Wrap up & References

Solutions

- ▶ If one is lucky, optimising a quadratic function \equiv solving $Q\bar{x} = -q$
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- ▶ Iterative procedures: start from **initial guess** x^0 , some process $x^i \rightsquigarrow x^{i+1}$
 \implies a **sequence** $\{x^i\}$ that should “go towards an optimal solution”
- ▶ The natural way: $\{f^i = f(x^i)\}$ **sequence of values** “go towards f_* ”
- ▶ Typically we **can't get f_* in finite time** ($\exists i \ v_i = f_*$), but we can
 “get as close as we want”: **there in the limit**
- ▶ Recall: (infinite) sequence $\{v_i\} = \{v_1, v_2, \dots\}$,
 $\{v_i\} \rightarrow v \equiv \lim_{i \rightarrow \infty} v_i = v \equiv \forall \varepsilon > 0 \exists h \text{ s.t. } |v_i - v| \leq \varepsilon \forall i \geq h$
 $\lim_{i \rightarrow \infty} v_i = +\infty \iff \forall M > 0 \exists h \text{ s.t. } v_i \geq M \forall i \geq h$

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- ▶ $\{x^i\}$ s.t. $\{f^i\} \rightarrow f_*$ a **minimizing sequence**
- ▶ note that $\{f^i\} \rightarrow -\infty \implies f_* = -\infty \implies$ minimizing sequence
- ▶ A sequence **may not have limit**: are we “not converging”?
- ▶ **Any monotone sequence has a limit** (monotone algorithms are good)

- ▶ We generally assume minimization, but maximization is equivalent
- ▶ Given x^i , necessarily compute $g^i = Qx^i + q$: if $g^i = 0$ then stop
- ▶ “ $g^i = 0$ ” not doable in floating point arithmetic $\implies \|g^i\| \leq \varepsilon$ (which ε ?)
- ▶ Idea: if $\|g^i\| > [\varepsilon >] 0$, produce a x^{i+1} “better” than x^i
- ▶ How? Consider the **tomography** $\varphi_{x^i, -g^i}(\alpha) = f(x^i - \alpha g^i) - f(x^i)$

$$= \frac{1}{2}(x^i - \alpha g^i)^T Q(x^i - \alpha g^i) + q(x^i - \alpha g^i) - f(x^i)$$

$$= \frac{1}{2}\alpha^2 (g^i)^T Q g^i - \alpha [(g^i)^T Q x^i + q g^i] = \frac{1}{2}\alpha^2 (g^i)^T Q g^i - \alpha \|g^i\|^2$$

positive negative
- ▶ For **some** $\alpha > 0$, $\varphi_{x^i, -g^i}(\alpha) < 0 \implies f(x^i - \alpha g^i) < f(x^i)$

Exercise: Check all the above (recall ▶ Optimizing a quadratic non-homogeneous function)

- ▶ The **same information** (called **gradient**, we'll see why) saying “you **can't stop**” is at the same time saying “you can **get a better solution than x^i** over there”
- ▶ This immediately suggests a (**monotone**, $f^{i+1} < f^i$) algorithm

- ▶ In fact it is easy to minimize $\varphi_{x^i, -g^i}(\alpha)$ (Optimizing a quadratic non-homogeneous function)
 $\alpha^i = \|g^i\|^2 / ((g^i)^T Q g^i)$ [$1/\lambda_1 \leq \alpha \leq 1/\lambda_n$ (check)]
- ▶ Computing g^i and the optimal value of α is $O(n^2) \implies$
 n “large” \implies “we can do many iterations before hitting $O(n^3)$ ”

```

procedure  $x = SDQ(Q, q, x, \varepsilon)$ 
  do forever
     $g \leftarrow Qx + q;$ 
    if ( $\|g\| \leq \varepsilon$ ) then break;
     $\alpha \leftarrow \text{stepsize}(); x \leftarrow x - \alpha g;$ 

```

- ▶ $\text{stepsize}() \{ \text{return}(\|g\|^2 / (g^T Q g)); \}$, others possible

Exercise: something can go wrong with that formula \uparrow : what does it mean?
 Improve the pseudo-code to take that occurrence into account.

Exercise: what happens if $Q \not\preceq 0$? Does the (improved) code need be fixed?

Exercise: Discuss how to change the code to solve $\max\{f(x)\}$ instead

Exercise: Rewrite the code with one product with Q per iteration

- ▶ It is very simple, but does it work? And is it efficient?

- ▶ Optimal stepsize $\implies g^{i+1} \perp g^i$ (check)
- ▶ “Homogeneous form of the error”: $A(x) = \frac{1}{2}(x - x_*)^T Q(x - x_*)$ (check)
- ▶ The above for $x = x^{i+1}$, $Q \succ 0$ and some algebra [5, Lm. 8.6.1] gives

$$A(x^{i+1}) = \left(1 - \frac{\|g^i\|^4}{((g^i)^T Q g^i)((g^i)^T Q^{-1} g^i)} \right) A(x^i) \quad \text{(check)[tedious]}$$

- ▶ Easy to derive an estimate using $\kappa = \lambda_1 / \lambda_n \geq 1$ condition number of Q

$$\frac{\|x\|^4}{(x^T Q x)(x^T Q^{-1} x)} \geq \frac{\lambda_n}{\lambda_1} = \frac{1}{\kappa} \quad \text{(check)} \implies A(x^{i+1}) \leq \left(1 - \frac{1}{\kappa} \right) A(x^i)$$

- ▶ This means the algorithm converges: $A(x^i) \leq r^i A(x^0)$ (check) with $r \leq (\kappa - 1) / \kappa < 1 \implies A(x^i) \rightarrow 0$ exponentially fast as $i \rightarrow \infty$

- ▶ Kantorovich inequality [5, 8.6.(34)] gives better estimate

$$\frac{\|x\|^4}{(x^T Q x)(x^T Q^{-1} x)} \geq \frac{4\lambda_1 \lambda_n}{(\lambda_1 + \lambda_n)^2} \implies r \leq \left(\frac{\lambda_1 - \lambda_n}{\lambda_1 + \lambda_n} \right)^2 = \left(\frac{\kappa - 1}{\kappa + 1} \right)^2$$

- ▶ Let's see it in practice

- ▶ Crucial sequences: $\{x^i\} / \{d^i = \|x^i - x_*\|\}$ [iterates / distance from x_*]
 $\{f^i = f(x^i)\} / \{a^i = A(x^i)\} / \{r^i = R(x^i)\}$ [f -values / A/R gaps]
- ▶ Complexity as a function of prescribed accuracy ε :
 max number of iterations k such that $d^i / a^i / r^i \leq \varepsilon \forall i \geq k$
- ▶ General formula: $v^k \leq r^k v^1 \leq \varepsilon$ for $k \geq \lceil 1 / \log(1/r) \rceil \log(v^1 / \varepsilon)$ (check)
- ▶ $r \approx 1 \implies k \in O(\lceil r / (1-r) \rceil \log(v^1 / \varepsilon))$ (check)
- ▶ Good news: dimension independent (n not there) \implies very-large-scale
- ▶ $O(\log(1/\varepsilon))$ (good), but the constant $\uparrow \infty$ as $r \rightarrow 1$ (bad)
- ▶ $v^1 = f(x^1) - f_*$: starting closer to f_* helps (would be strange if not)
- ▶ “ $\|x^i - x_*\| \leq \varepsilon$ ” and “ $f(x^i) - f_* \leq \varepsilon$ ” not the same (ε):
 $a^i = \frac{1}{2}(x^i - x_*)^T Q(x^i - x_*) \leq \varepsilon \implies \lambda_n \|x^i - x_*\|^2 \leq \varepsilon \implies$
 $d^i = \|x^i - x_*\| \leq \sqrt{\varepsilon / \lambda_n}$

Exercise: Cook up the other direction ($d^i \leq \varepsilon \implies \dots$)

► Converge: $\{f^i\} \rightarrow f_* \approx \equiv \{a^i\} \rightarrow 0 \equiv \{r^i\} \rightarrow 0 \iff \{d^i\} \rightarrow 0$ (\nRightarrow)

Exercise: Discuss why $\{f^i\} \rightarrow f_*$ is only $\approx \equiv$ to $\{a^i\} \rightarrow 0$ and why the \nRightarrow

► But how rapidly does it (“in the tail”)? Rate/order of convergence

$$\lim_{i \rightarrow \infty} \left[\frac{f^{i+1} - f_*}{(f^i - f_*)^p} = \frac{a^{i+1}}{(a^i)^p} \approx \frac{r^{i+1}}{(r^i)^p} \right] = r \quad \left[\begin{array}{l} x^p \rightarrow 0 \text{ faster than} \\ x \rightarrow 0 \text{ when } p > 1 \end{array} \right] \quad (\text{check})$$

► $p = 1$, $r = 1 \equiv$ **sublinear**: important examples

error	$O(1/i)$	$O(1/i^2)$	$O(1/\sqrt{i})$
i	$O(1/\varepsilon)$ (bad)	$O(1/\sqrt{\varepsilon})$ (a bit better)	$O(1/\varepsilon^2)$ (horrible)

► $p = 1$, $r < 1 \equiv$ **linear**: $r^i \implies i \in O(\log(1/\varepsilon))$ (good unless $r \approx 1$)

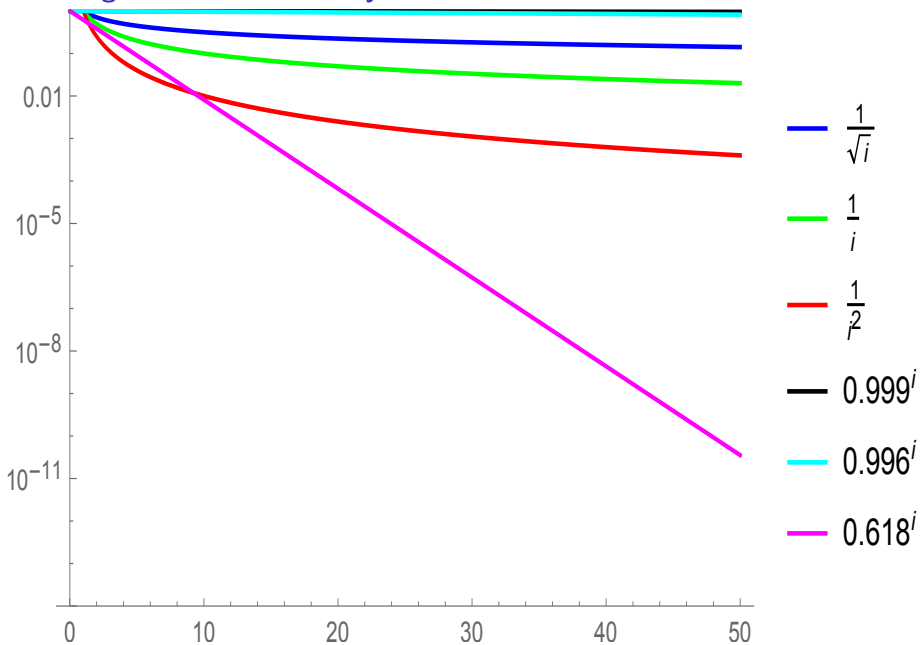
► $p = 2$, $r > 0 \equiv$ **quadratic (!!!)**: $\approx 1/2^{2^i} \implies i \in O(\log(\log(1/\varepsilon)))$
in practice $O(1)$ (correct digits **double** at each iteration)

► $p \in (1, 2) \equiv p = 1$, $r = 0 \equiv$ **superlinear (!)**: “something in the middle”

► $p = 2$ the best you can reasonably hope for: **possible** but **not easy**

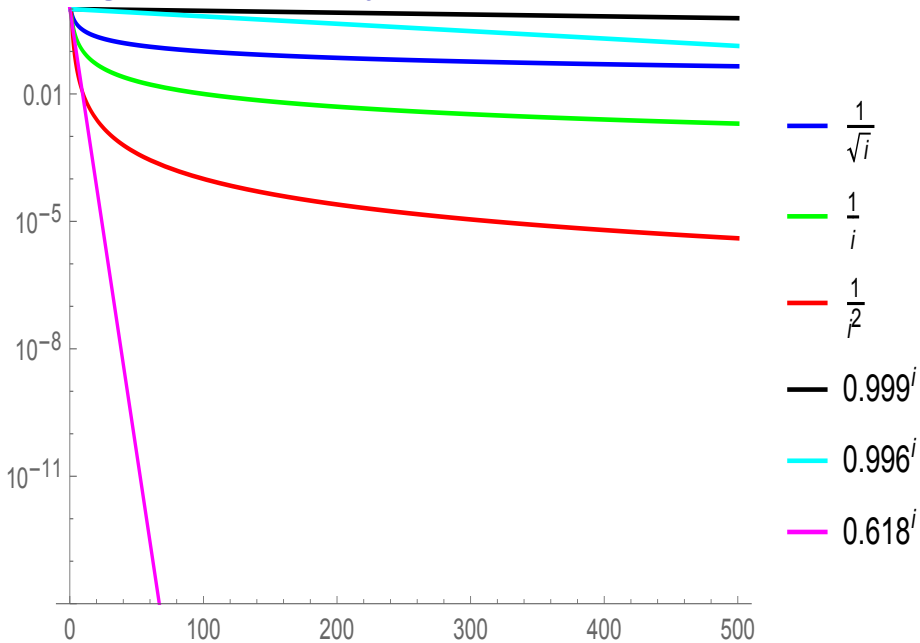
Convergence Rates Pictorially

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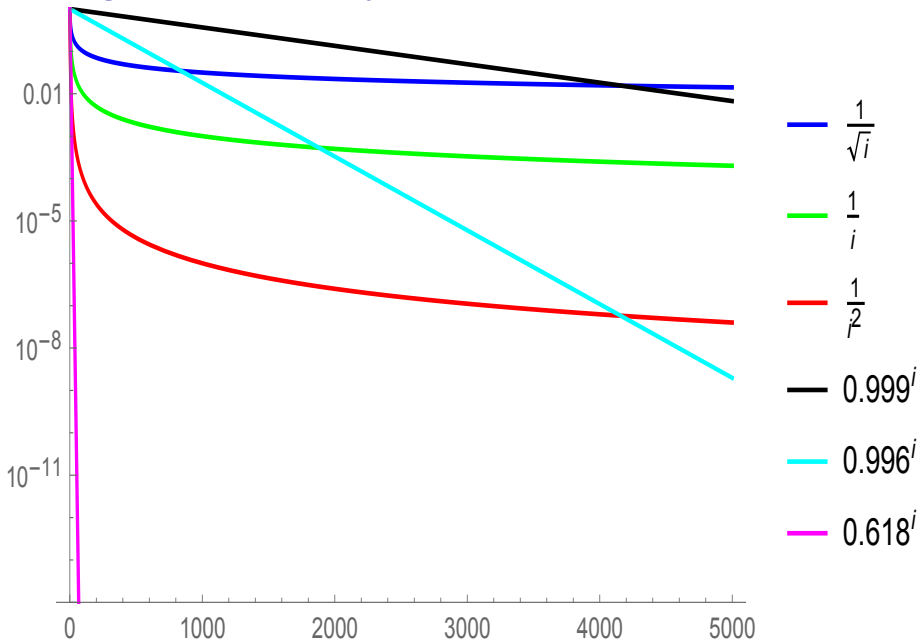
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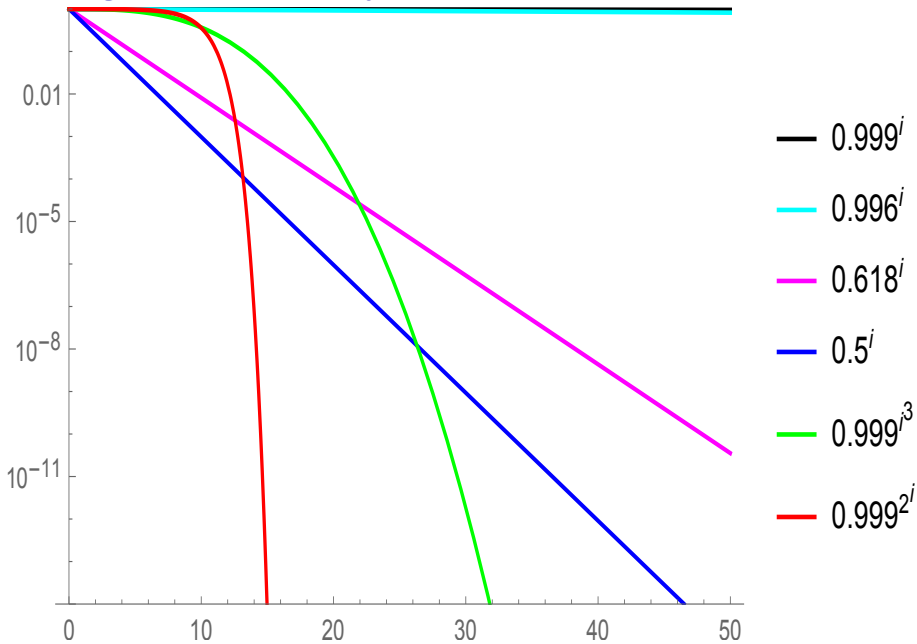
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- ▶ The stopping criterion one **would want**: $A(x^i) \leq \varepsilon / R(x^i) \leq \varepsilon$
- ▶ Issue: f_* **typically unknown**, cannot be used on-line
- ▶ $\|g^i\|$ “proxy” of $A(x^i)$: **hopefully** $\|g^i\|$ “small” $\implies A(x^i)$ “small”
but **exact relationship nontrivial** \implies choosing ε non obvious
- ▶ $\|g^i\| = Q(x^i - x_*) \implies \|g^i\| \leq \lambda_1 \|x^i - x_*\| \dots (??)$ **wrong inequality**:
 $\|g^i\| \leq \varepsilon \not\Rightarrow \|x^i - x_*\|$ “small”
- ▶ $a^i = \frac{1}{2}(x^i - x_*)^T Q(x^i - x_*) = \frac{1}{2} \langle x^i - x_*, g^i \rangle \leq \frac{1}{2} \|g^i\| \|x^i - x_*\|$;
if we knew $\delta \geq \|x^i - x_*\|$, which we **don't**, then $\|g^i\| \leq 2\varepsilon / \delta \implies a^i \leq \varepsilon$
- ▶ If we knew $\lambda_n > 0$, which we **don't**, $\|g^i\| \leq \sqrt{2\lambda_n \varepsilon} \implies a^i \leq \varepsilon$ (**check**)
- ▶ All in all, **exact** control on final a^i / r^i not obvious (not always really needed)

- ▶ Convergence **fast** if $\lambda_1 \approx \lambda_n$ (one iteration for $\|x\|^2$), **rather slow** if $\lambda_1 \gg \lambda_n$:
 $\kappa = \lambda_1 / \lambda_n \rightarrow \infty$ (Q ill conditioned) $\implies r \rightarrow 1 \implies$ **slow** in practice
- ▶ $g^{i+1} \perp g^i$ + level sets very elongated \implies lots of “zig-zags” \implies **slow**
- ▶ Ex.: $\kappa = 1000 \implies r \approx 0.996 \implies r / (1 - r) \approx 250$
 $f(x^1) - f_* = 1, \varepsilon = 10^{-6} \implies k \geq 3450$ for $n = 2$

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 $f(x^1) - f_* = 1, \varepsilon = 10^{-6} \implies k \geq 3450$ for $n = 2 \dots$ but also for $n = 10^8$
- ▶ Note: with coarser formula $r = 0.999 \equiv r / (1 - r) \approx 1000 \implies k \geq 13800$
- ▶ In other words: $0.996^{10} \approx 0.96071$ $0.999^{10} \approx 0.99004$

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- ▶ Ex.: $\kappa = 1000 \implies r \approx 0.996 \implies r / (1 - r) \approx 250$
 $f(x^1) - f_* = 1, \varepsilon = 10^{-6} \implies k \geq 3450$ **for $n = 2$... but also for $n = 10^8$**
- ▶ Note: with coarser formula $r = 0.999 \equiv r / (1 - r) \approx 1000 \implies k \geq 13800$
- ▶ In other words: $0.996^{100} \approx 0.66978$ $0.999^{100} \approx 0.90479$

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- ▶ More bad news, “hidden dependency”:
 λ_1 and λ_n **may depend on n** , κ may grow as $n \rightarrow \infty$
- ▶ More bad news: **the behaviour in practice is close to the bound**
- ▶ Even more bad news: **$\lambda_n = 0 \equiv \kappa = \infty$ happens**

What if $\lambda_n = 0$?

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- ▶ $\lambda_n = 0 \implies$ not converging? No, just can't prove it this way
- ▶ In fact we can prove convergence (in a more general setting) [2, Theorem 3.3]:
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- ▶ $O(1 / \varepsilon)$ vs. $O(\log(1 / \varepsilon))$: sublinear convergence, exponentially slower
- ▶ One further digit of accuracy \equiv 10 times more iterations \implies typically unfeasible to get more than $1e-3 / 1e-4$ accuracy
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- ▶ Is it bad? Rather. Can it be worse? Yes (in general, will see)
- ▶ If $\lambda_n > 0$, can we do better than $O(\log(1 / \varepsilon))$? Yes – @Federico
- ▶ Fundamental idea, will see more than once: changing the space

Outline

Optimization Problems

Optimization is difficult

Simple Functions, Univariate case

Simple Functions, Multivariate case

Multivariate Quadratic case: Gradient Method

Wrap up & References

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- ▶ Solving (simple) optimization problems requires linear algebra, and vice-versa
- ▶ We now know all we need about simple problems, time to step up the game
- ▶ Will keep following an incremental approach: next step is **more complicated functions** but **only one variable**

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- ▶ Use $\max\{|f_*|, 1\}$ instead; this corresponds to $\min\{f(x) + 1\}$ [back]
- ▶ $b > 0$ and $x - z > 0 \implies b(x - z) > 0 \equiv bx > bz$;
the others are analogous (or simpler) [back]
- ▶ If $x_+ = +\infty$, obviously $x_* = +\infty = x_+$
If $x_+ < +\infty$, since $f(x)$ is increasing, $f(x) < f(x_+) \forall x < x_+$
The treatment of x_- is analogous.
If $b < 0$, the role of x_+ and x_- reverses ($x_+ = \operatorname{argmin}$, $x_- = \operatorname{argmax}$)
If $b = 0$, every point in X is an optimal solution [back]
- ▶ $x > z$, $a > 0$ and $x > 0 \implies ax^2 > axz > az^2$. Since $f(x)$ is symmetric ($ax^2 = a(-x)^2$), increasing for $x > 0 \equiv$ decreasing for $x < 0$. When $a < 0$ the sign of the inequalities is inverted (the function is reflected upon the x axis).
The case $a = 0$ is trivial [back]

- ▶ $f(x)$ has a minimum in 0 , is decreasing for $x < 0$ and increasing for $x > 0$. If $x_- > 0$ then $f(x)$ is increasing along all X , hence x_- is the minimum and x_+ the maximum. The argument is symmetric if $x_+ < 0$. Obviously, if $0 \in X$ then it is the minimum; for the maximization, since the function is increasing when x moves away from 0 in both directions, the maximum has to be one of the two extremes but we don't know which until we test. The rest is too trivial [back]
- ▶ No, this is both too trivial and didactic [back]
- ▶ $f(x) = (ax + b)x$, hence the roots are $x = 0$ and $x = x_p = -b/a$. Clearly, $\bar{x} = -b/2a$ is always in the middle of the interval defined by the roots. If a and b have the same sign then $x_p < \bar{x} < 0$, otherwise $x_p > \bar{x} > 0$ [back]
- ▶ $\varphi_{x,(\beta d)}(\alpha) = f(x + \alpha(\beta d)) = f(x + (\alpha\beta)d) = \varphi_{x,d}(\alpha\beta)$ [back]

- ▶ We assume that i. and ii. hold for f and we want to show that $f(x) = \langle b, x \rangle$ for some $b \in \mathbb{R}^n$. Let u_i , $i = 1, \dots, n$, the i -th vector of the canonical base of \mathbb{R}^n (having 1 in the i -th position and 0 otherwise), and $b_i = f(u_i)$. For any $x \in \mathbb{R}^n$, $x = \sum_{i=1}^n x_i u_i$, hence $f(x) = f(\sum_{i=1}^n x_i u_i) = \sum_{i=1}^n f(x_i u_i)$ (using ii. recursively n times) $= \sum_{i=1}^n x_i f(u_i)$ (using i. on each individual term) $= \sum_{i=1}^n b_i x_i$ (using the definition of b_i) $= \langle b, x \rangle$ (using the definition of scalar product). The results clearly breaks in the affine case ($c \neq 0$):
 $f(x) = x + 1 \implies f(2x) = 2x + 1 \neq 2(x + 1) = 2f(x)$ [back]
- ▶ By contradiction, $\exists \gamma \in \mathbb{R}^n \setminus \{0\}$ s.t. $H\gamma = 0 \implies$
 $0 = \|H\gamma\|^2 = \gamma^T [H^T H] \gamma = \|\gamma\|^2 > 0$ [$\gamma \neq 0$] \neq [back]

- ▶ This is based on a general result: for $[A^1, A^2, \dots, A^n] = A \in \mathbb{R}^{m \times n}$ (not necessarily square) written by columns, $AA^T = M \in \mathbb{R}^{m \times m}$ (symmetric, prove it using $[AB]^T = B^T A^T$) can be written as the sum of the n rank-one matrices corresponding to the columns, i.e., $M = \sum_{i=1}^n [D^i = A^i(A^i)^T]$. In fact, the h -th row of A is $A_h = [A_h^1, A_h^2, \dots, A_h^n]$ and the k -th column of A^T is the k -th row of A , thus $M_{hk} = \langle A_h, A_k \rangle = \sum_{i=1}^n A_h^i A_k^i$. But $D_{hk}^i = A_h^i A_h^i$, hence $M_{hk} = \sum_{i=1}^n D_{hk}^i$ for all h and k .
- To complete the result, for $\Lambda = \text{diag}([\lambda_1, \lambda_2, \dots, \lambda_n]) \in \mathbb{R}^{n \times n}$, $L = A\Lambda = [\lambda_1 A^1, \lambda_2 A^2, \dots, \lambda_n A^n]$. In fact, the h -th row $A_h = [A_h^1, A_h^2, \dots, A_h^n]$ and the k -th column of Λ , i.e., $\lambda_k u_k$ (u_k being the k -th vector of the canonical base) give $L_{hk} = \langle A_h, \lambda_k u_k \rangle = \lambda_k A_h^k$ [back]
- ▶ $\varphi_{H_i}(\alpha) = (\alpha H_i)^T Q(\alpha H_i) = \alpha^2 [H_i^T (\lambda_i H_i)] = \lambda_i \alpha^2$ [back]
- ▶ $\lambda_n < 0 \implies \varphi_{H_n}(\alpha) [= \lambda_n \alpha^2]$ unbounded below $\implies f(x)$ unbounded below
 $\lambda_1 > 0 \implies \varphi_{H_1}(\alpha)$ unbounded above $\implies f(x)$ unbounded above [back]

▶ $x = z + \bar{x} \implies \frac{1}{2}x^T Qx + qx = \frac{1}{2}(z + \bar{x})^T Q(z + \bar{x}) + q(z + \bar{x}) =$
 $\frac{1}{2}z^T Qz + z^T(Q\bar{x} + q) + [\frac{1}{2}\bar{x}^T Q\bar{x} + q\bar{x}] = \frac{1}{2}z^T Qz + f(\bar{x})$
 as $Q\bar{x} + q = Q(-Q^{-1}q) + q = -q + q = 0$ [back]

▶ $Qv = Q[\sum_{i \in Z} \eta_i H_i] = \sum_{i \in Z} \eta_i QH_i = \sum_{i \in Z} \eta_i \lambda_i H_i = 0$ [back]

▶ $Q = H\Lambda H^T = \sum_{i=1}^n \lambda_i H_i H_i^T = \sum_{i \in Z} \lambda_i H_i H_i^T [= 0] + \sum_{i \in N} \lambda_i H_i H_i^T$

We want to prove $\exists x$ s.t. $(\sum_{i \in N} \lambda_i H_i H_i^T)x = \sum_{i \in N} \mu_i H_i = w$

True if $\lambda_i H_i^T x = \mu_i \quad i \in N \equiv H_i^T x = \gamma_i = \mu_i / \lambda_i \quad i \in N,$

a linear system of $k \leq n$ equations in n variables (likely underdetermined)

All H_i linearly independent, $H_N = [H_i]_{i \in N} \in \mathbb{R}^{n \times k} \implies \text{rank}(H_N) = k$

$\implies [H_N^T, \gamma] \in \mathbb{R}^{k \times n+1}$ has rank k (rank \leq number of rows) \implies

by [16] the system has a solution x (∞ -ly many if $k < n$) [back]

▶ $\frac{1}{2}x^T Qx + qx = \frac{1}{2}(z + \bar{x})^T Q(z + \bar{x}) + q(z + \bar{x}) =$
 $\frac{1}{2}z^T Qz + z^T(Q\bar{x} + q) + f(\bar{x}) = \frac{1}{2}z^T Qz + q^0 z + f(\bar{x})$ [back]

- ▶ We know that $f(z) = z^T Qz + f(\bar{x})$, with $z = x - \bar{x}$. For $x \in \bar{x} + v$, with $v \in \ker(Q)$, $z = x - \bar{x} = \bar{x} + v - \bar{x} = v$. Hence $f(z) = f(\bar{x})$. On the other hand, $f(z) \geq f(\bar{x})$ for all z since $Q \succeq 0$, thus any such point is a minimum. Any point $x \in \bar{x} + v$ with $v \notin \ker(Q)$ has $f(x) = v^T Qv + f(\bar{x}) > f(\bar{x})$ since $v^T Qv > 0$ [back]

- ▶ No, this is both too trivial and didactic [back]

- ▶ $\varphi(\alpha) = a\alpha^2 + b\alpha$ quadratic non-homogeneous with $a = (g^i)^T Qg^i \geq 0$ and $b = -\|g^i\|^2 < 0$. If $a > 0$, then $\varphi(\bar{\alpha}) < \varphi(0) = f(x^i) \forall \bar{\alpha} \in (0, -b/a)$; in particular, $\bar{\alpha} = \|g^i\|^2 / (2(g^i)^T Qg^i)$ is the minimum of φ . If $a = 0$ then φ is decreasing and $\varphi(\bar{\alpha}) < \varphi(0) = f(x^i) \forall \bar{\alpha} > 0$ [back]

- ▶ The variational characterization of the eigenvalues implies that $\lambda_1 \geq d^T Qd / \|d\|^2 \geq \lambda_n$ for all $d \neq 0$; this immediately gives $1/\lambda_1 \leq \|d\|^2 / d^T Qd \leq 1/\lambda_n$ for all d , and therefore in particular $d = g^i$ (knowing that $g^i \neq 0$ otherwise the algorithm would have stopped) [back]

- ▶ The issue clearly is $g^T Qg = 0$ (very small), which means that $\varphi_{x,-g}$ is (almost) linear, and therefore f is unbounded below. One should therefore add a line
if($g^T Qg \leq \delta$) then break;
for a “very small” δ , but also add a proper way for the algorithm to signal that the returned x is not optimal, e.g., by also returning a “status code” **[back]**
- ▶ Having added the extra check above, the code just works: if $g^T Qg < 0$ then $(-g)$ is direction where φ has negative curvature, which still implies f is unbounded below. Note that this is not guaranteed to happen **[back]**
- ▶ Because $a < 0$, the step α will be negative, which basically means one is going in direction g rather than $-g$. The algorithm remains the same, except that the extra check above has to become $g^T Qg \geq -\delta$ **[back]**

- ▶ Assuming the gradient is computed in the “natural way” as $g = Q * x + q$ before the algorithm starts (i.e., with x the initial guess x^0), both quantities depending from matrix-vector products can be recovered by computing the vector $v = Q * g$. In fact, $a = g^T Q g = \langle g, v \rangle$. Then, with $x' = x - \alpha g$ one has $g' = Qx' + q = Q(x - \alpha g) + q = (Qx + q) - \alpha Qg = g - \alpha v$. Hence, the gradient at the next iteration can be computed in $O(n)$ out of that of the previous iteration and the vector v . As for the objective function, $1/2x^T Qx + \langle q, x \rangle = 1/2(x^T Qx + 2\langle q, x \rangle) = 1/2x^T (Qx + q + q) = 1/2\langle q + g, x \rangle$, i.e., it can be computed in $O(n)$ once g is known **[back]**
- ▶ $g^i = Q(x^i - x_*) = Qx^i + q$, $\alpha^i = \|g^i\|^2 / [(g^i)^T Qg^i]$
 $g^{i+1} = Qx^{i+1} + q = Q(x^i - \alpha^i g^i) + q = (I - \alpha^i Q)g^i \implies$
 $\langle g^{i+1}, g^i \rangle = \|g^i\|^2 - \alpha^i [(g^i)^T Qg^i] = 0$ **[back]**

- ▶ All arguments boil down to the crucial $Qx^* + q = 0$. This first of all gives that $f(x^*) = \frac{1}{2}(x^*)^T Qx^* + \langle x^*, q \rangle = (x^*)^T Qx^* + \langle x^*, q \rangle - \frac{1}{2}(x^*)^T Qx^* = (x^*)^T (Qx^* + q) - \frac{1}{2}(x^*)^T Qx^* = -\frac{1}{2}(x^*)^T Qx^*$. Then, $\frac{1}{2}(x - x^*)^T Q(x - x^*) = \frac{1}{2}x^T Qx + \frac{1}{2}(x^*)^T Qx^* - x^T(Qx^*) = \frac{1}{2}x^T Qx - \langle x, q \rangle + \frac{1}{2}(x^*)^T Qx^* = f(x) - f(x^*)$ (in the penultimate step we have used $Qx^* = -q$) **[back]**
- ▶ Just induction: obvious for $i = 0$, if it holds for $i - 1$ then $A(x^i) \leq rA(x^{i-1}) \leq r(r^{i-1}A(x^0))$ **[back]**
- ▶ Q nonsingular $\implies x^i - x_* = Q^{-1}g^i \implies a^i = \frac{1}{2}(x^i - x_*)^T Q(x^i - x_*) = \frac{1}{2}(g^i)^T Q^{-1}g^i \implies a^{i+1} = \frac{1}{2}(x^{i+1} - x_*)^T Q(x^{i+1} - x_*) = \frac{1}{2}(x^i - \alpha^i g^i - x_*)^T g^{i+1} = \frac{1}{2}(x^i - x_*)^T g^{i+1}$
 [using $\langle g^{i+1}, g^i \rangle = 0$] $= \frac{1}{2}(x^i - x_*)^T Q(x^i - \alpha^i g^i - x_*)$
 $= \frac{1}{2}(x^i - x_*)^T Q(x^i - x_*) - \frac{1}{2}\alpha^i (x^i - x_*)^T Qg^i = a^i - \frac{1}{2}\alpha^i \|g^i\|^2$
 [using $Q(x^i - x_*) = g^i$] $= a^i - \frac{1}{2}\|g^i\|^4 / (g^i)^T Qg^i$

$$= a^i - \frac{\|g^i\|^4}{((g^i)^T Q g^i)((g^i)^T Q^{-1} g^i)} = a^i \left(1 - \frac{\|g^i\|^4}{((g^i)^T Q g^i)((g^i)^T Q^{-1} g^i)} \right) \quad \text{[back]}$$

- ▶ Recall $1/\lambda_n \geq \dots \geq 1/\lambda_1 > 0$ eigenvalues of Q^{-1} ; from the usual $\lambda_n \|x\|^2 \leq x^T Q x \leq \lambda_1 \|x\|^2$ (applied to Q^{-1} as well) one has $\|g\|^2 / g^T Q g \geq 1/\lambda_1$ and $\|g\|^2 / g^T Q^{-1} g \geq 1/[1/\lambda_n]$ [back]
- ▶ $r^k v_1 \leq \varepsilon \equiv r^k \leq \varepsilon / v^1 \equiv \log(r^k) \leq \log(\varepsilon / v^1)$ (log monotone) $\equiv k \log(r) \leq \log(\varepsilon / v^1)$ (property of log); since $r < 1$, $\log(r) < 0$, giving $k \geq \log(\varepsilon / v^1) / \log(r) = [-\log(\varepsilon / v^1)] / [-\log(r)] = \log(v^1 / \varepsilon) / \log(1/r) = \log(v^1 / \varepsilon)[1 / \log(1/r)]$ [back]
- ▶ This requires a bit of elementary calculus. The derivative of $\ln(x)$ is $1/x$. The first-order Taylor approximation is $f(x + \delta) \approx f(x) + f'(x)\delta$ for $\delta \approx 0$. Applied to $\ln(\cdot)$ with $x = 1$ gives $\ln(1 + \delta) \approx \delta$, whence $1 / \ln(1/r) = 1 / \ln(1 + (1-r)/r) = r / (1-r)$. But $\log_a(x) = \log_b(x) / \log_b(a)$, hence $\ln(x) = \log_e(x) = \log_{10}(x) / \log_{10}(e) \approx \log(x) / 0.43 \approx 2.3 \log(x)$, i.e., $\ln(x) \in O(\log(x))$ [back]

- ▶ $\lambda_1 \|x^i - x_*\|^2 \geq (x_i - x_*)^T Q (x_i - x_*) = 2a^i \equiv \|x^i - x_*\| \geq \sqrt{2a^i / \lambda_1}$,
hence $d^i \leq \varepsilon \implies a^i \leq \lambda_1 \varepsilon^2 / 2$ **[back]**

- ▶ $a^i = \frac{1}{2} (x^i - x_*)^T Q (x^i - x_*) = \frac{1}{2} \langle g^i, x^i - x_* \rangle \leq \frac{1}{2} \|g^i\| \|x^i - x_*\|$. On the other hand, $\|g^i\|^2 = (x^i - x_*)^T Q^T Q (x^i - x_*) \geq \lambda_n^2 \|x^i - x_*\|^2$ (recall λ_n^2 eigenvalue of Q^2 , clearly the smallest), i.e., $\|g^i\| \geq \lambda_n \|x^i - x_*\|$. Hence, $\|g^i\| \leq \sqrt{2\lambda_n \varepsilon} \implies \varepsilon \geq \frac{1}{2\lambda_n} \|g^i\|^2 \geq \frac{1}{2} \|g^i\| \|x^i - x_*\| \geq a^i$ **[back]**

- ▶ If $f_* = -\infty$, $f_i \rightarrow -\infty$ is OK (minimising sequence) but $a^i = a^{i+1} = \infty$ and therefore their ratio is not well-defined. Since f is continuous, $\{d^i\} \rightarrow 0 \implies \{a^i\} \rightarrow 0$, but the converse need not happen in general: say, $\{x^{2i}\} \rightarrow x'_*$ and $\{x^{2i+1}\} \rightarrow x''_*$ with $x'_* \neq x''_*$ optimal solutions **[back]**

- ▶ Simply, $\lim_{x \rightarrow 0} x^p / x = \lim_{x \rightarrow 0} x^{p-1} = 0$: the numerator goes to 0 faster than the denominator **[back]**