

$Ax = y$ $A \in \mathbb{R}^{m \times m}$ large and sparse

Direct methods

Iterative methods

- general $A = LU$
- symm. $A = LDL^T$
- posdef Cholesky $A = R^T R$

- GMRES
- MINRES
- Conjugate gradient

Effective when the factors are sparse

Effective when eigenvalues are clustered

("large-world models", e.g. road networks)

Trick: reordering: we can apply a symmetric permutation

$A \rightarrow P^T A P$

$$\begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

Without reordering:

$A = R^T R$

$Ax = y \Leftrightarrow R^T \underbrace{Rx}_z = y$

$$\begin{cases} R^T z = y \\ Rx = z \end{cases}$$

① $R = \text{chol } A$

② $z = R^T \setminus y$



forward subst.

③ $x = R \setminus z$



back-substitution

With reordering:

$R^T R = P^T A P$

$P^T A P P^T x = P^T y$

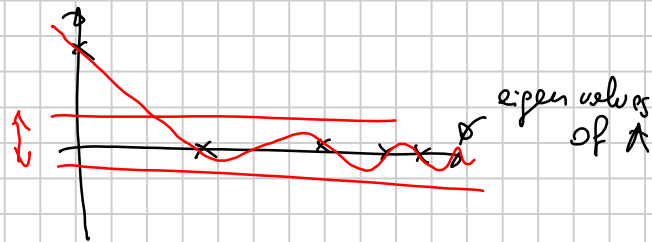
$R^T \underbrace{R P^T x}_z = P^T y \Leftrightarrow \underbrace{R^T z}_w = P^T y$

$$\begin{cases} R^T w = P^T y \\ R z = w \\ P^T x = z \end{cases} \quad \begin{matrix} \text{forward} \\ \text{back} \\ \text{forward} \end{matrix} \quad x = P z$$

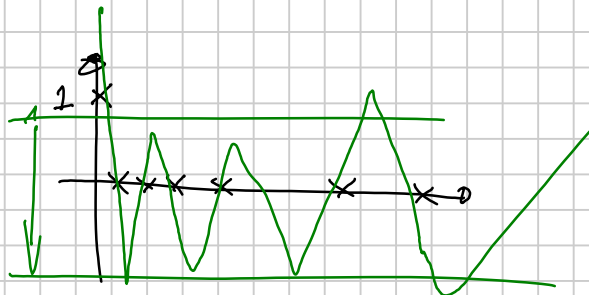
~ 30% cpu savings in our example, due to sparser R.

"fill-in" is reduced.

small
min max $|q(\lambda)|$
 $\rho \quad \Delta$



large
min max $|q(\lambda)|$
 $\rho \quad \Delta$



⚠ Don't use $\kappa(A)$

• slower

• less accurate:

$$A^{-1} = \left[A^{-1}e_1 \mid A^{-1}e_2 \mid \dots \mid A^{-1}e_n \right], \text{ so when}$$

$$\text{computing } A^{-1}y = A^{-1}e_1 y_1 + \dots + A^{-1}e_n y_n$$

we are as accurate as the worse of these linear systems!

Let $A \in \mathbb{R}^{n \times n}$, non symmetric, with "bad" sparsity pattern and "bad" eigenvalues.

Transforming $Ax=y$ to $PAx=Py$, for any P invertible, gives an equivalent linear system but different eigenvalue location!

Preconditioning: P is a preconditioner for the problem (or P^{-1})

Idea: if $PA \approx I$, its eigenvalues are close to 1.

$P \approx A^{-1}$ improves eigenvalue location.

We want P such that (i) cheap to compute

~~(1) cheap to invert~~

(3) $PA \approx I$

EX: $P = \begin{bmatrix} a_{11} & a_{12} & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & \dots & \dots & a_{nn} \end{bmatrix}^{-1}$ (effective only on some problems)

Better strategy: incomplete LU factorization

$$\begin{pmatrix} \times & \times & \times \\ \otimes & \times & \times \\ \otimes & \times & \times \\ \otimes & \times & \times \end{pmatrix}$$

Start computing an LU factorization, and set small entries to \otimes artificially along the way!

$$A \approx LU$$

\uparrow not equal anymore!

\rightarrow much more sparse!

$$P = U^{-1}L^{-1} \approx A$$

For iterative methods, we do not need the entries of A , but just a way to compute $w \mapsto Aw$

Instead of passing $P \cdot A$, we pass the function

$$w \mapsto U^{-1}L^{-1}Aw = U \setminus (L \setminus (Aw))$$

Better to use this "callback form" of guess/pcg when working with preconditioners.

⚠ the residual $A\tilde{x} - y$ is not equal in general to the "preconditioned residual" $PA\tilde{x} - Py$.

This is called "left preconditioning". There is also "right preconditioning":

$$(AP)(P^{-1}x) = y$$

⤴

Is there a way to use preconditioning for symmetric systems? Note that even if A is symmetric, PA usually is not \Rightarrow $PCG(PA, Py)$ fails!

Idea: multiply from both sides to keep symmetry!

Look for P such that

$$P^T A P \approx I$$

and solve with an iterative method

$$\underbrace{P^T A P}_{\hat{A}} \underbrace{P^{-1}x}_{\hat{x}} = \underbrace{P^T y}_{\hat{y}} \Leftrightarrow Ax = y$$

Symmetric preconditioning

A symm. pos. def. $\Leftrightarrow P^T A P$ symm. pos. def.

General preconditioner for A symm. pos. def.:
incomplete Cholesky

$\Rightarrow R = \text{ichol}(A)$ returns a sparse R s.t. $A \approx RR^T$

$$A \approx RR^T \quad R^{-1}A(R^{-1})^{-1} \approx I$$

$$\underbrace{R^{-1}AR^{-1}}_{\hat{A}} \underbrace{R^{-1}x}_{\hat{x}} = \underbrace{R^{-1}y}_{\hat{y}}$$

$$\hat{A}w = R^{-1}AR^{-1}w$$