

Es 1:

$$f(x) = x - 1 + \frac{1}{x+1} \quad x \neq -1$$

2) Condizioni di convergenza

$$\varepsilon_n = \frac{x}{f(x)} \quad f'(x) \in \mathbb{R}$$

$$f(x) = x - 1 + \frac{1}{x+1} = \frac{x^2}{x+1}$$

$$f'(x) = 1 - \frac{1}{(x+1)^2} = \frac{x^2+2x}{(x+1)^2}$$

$$\Rightarrow \varepsilon_n = \frac{x}{\cancel{x^2}} \cdot \cancel{(x+1)} \cdot \frac{x(x+2)}{(x+1)^2} \underset{x \rightarrow \infty}{\approx} \frac{(x+2)}{x+1} \varepsilon_x$$

$$|\varepsilon_n| = \left| \frac{x+2}{x+1} \right| |\varepsilon_x| \leq \frac{|x+2|}{|x+1|} u$$

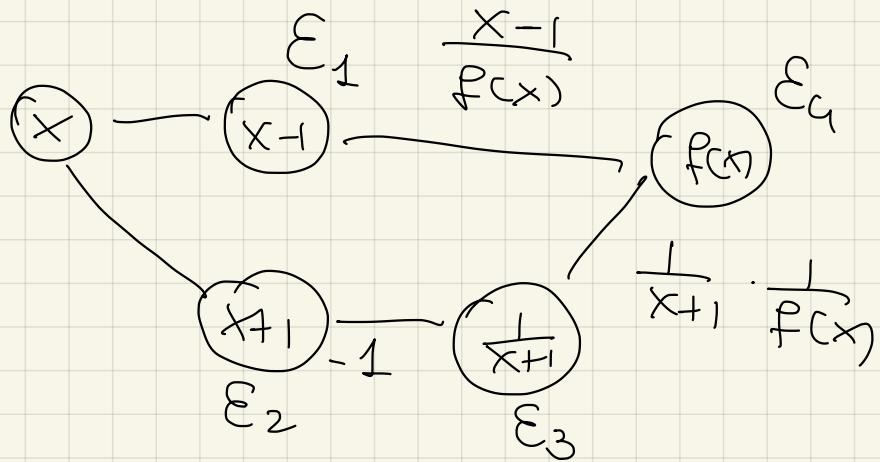
per $x \rightarrow -1$ il problema risulta malcondizionato
per $x \rightarrow \pm \infty$ il problema è invece ben condizionato

b) Stabilità: errore algoritmo.

$$g_1(x) = x - 1 + \frac{1}{x+1}$$

$$g_2(x) = \frac{x^2}{x+1}$$

①



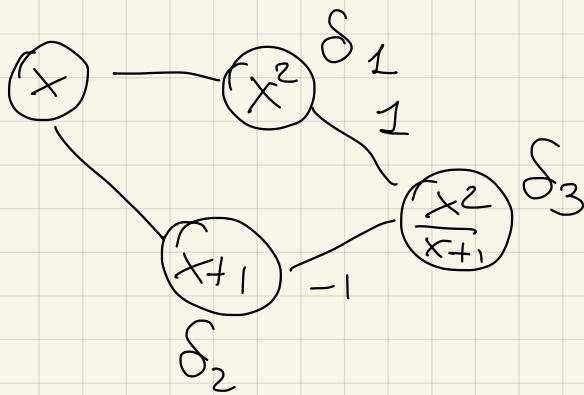
$$\varepsilon_{\text{ALG}_4}^{(1)} = \varepsilon_4 + \frac{x-1}{f(x)} \varepsilon_1 + \frac{1}{(x+1)} \cdot \frac{1}{f(x)} (\varepsilon_3 - \varepsilon_2)$$

$$\begin{aligned}
 |\varepsilon_{\text{ALG}_4}^{(1)}| &= \left| \varepsilon_4 + \frac{x-1}{f(x)} \varepsilon_1 + \frac{1}{(x+1)} \cdot \frac{1}{f(x)} (\varepsilon_3 - \varepsilon_2) \right| \\
 &\leq |\varepsilon_4| + \frac{|x-1|}{|f(x)|} |\varepsilon_1| + \frac{1}{|x+1|} \cdot \frac{1}{|f(x)|} (|\varepsilon_3| + |\varepsilon_2|) \\
 &\leq \alpha \left(1 + \frac{1}{|f(x)|} \left(|x-1| + \frac{2}{|x+1|} \right) \right) \\
 &= \alpha \left(1 + \frac{(x+1)}{x^2} \cdot \left(|x-1| + \frac{2}{|x+1|} \right) \right) \\
 &= \alpha \left(1 + \frac{|x^2-1|}{x^2} + \frac{2}{x^2} \right)
 \end{aligned}$$

Questo algoritmo è instabile per $x \rightarrow 0$, mentre è stabile per $|x| \rightarrow \infty$ con le stesse riferenze.

②

$$g_2(x) = \frac{x^2}{x+1}$$



$$\mathcal{E}_{ACG}^{(2)} = S_3 + S_1 - S_2$$

$|\mathcal{E}_{ACG}^{(2)}| \leq 3u$ Questo algoritmo è sempre preferibile perché stabile per ogni soluzione di x .

Ese 2:

$$A = I + \lambda e_n e_n^T + \lambda e_n u^T \quad \lambda \in \mathbb{R}$$

$$u = [0, 1 \dots 1]$$

$$e_n e_n^T = \begin{bmatrix} 0 \\ \vdots \\ 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & \dots & 0 \end{bmatrix} = \begin{bmatrix} & & & \\ & & & \\ & & 0 & \\ & & & \\ 1 & & & \end{bmatrix}$$

$$e_n u^T = \begin{bmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{bmatrix} \begin{bmatrix} 0, 1 & \dots & 1 \end{bmatrix} = \begin{bmatrix} 0 & 1 & \dots & 1 \\ 0 & & & \end{bmatrix}$$

$$\Rightarrow A = \begin{bmatrix} 1 & \alpha & \dots & \alpha \\ \alpha & 1 & \dots & \alpha \\ \vdots & \vdots & \ddots & \vdots \\ \alpha & \alpha & \dots & 1 \end{bmatrix}$$

2) Possono usare la pred. diagonale per colonne.
 Se $|\alpha| < 1 \Rightarrow$ la matrice è a predominanza diagonale per colonne col è puram. invertibile

$$\text{col. } 1: \quad 1 > |\alpha|$$

$$\text{col. } 2: \dots n: \quad |\alpha_{jj}| = 1 > \sum_{\substack{i=1 \\ i \neq j}}^n |\alpha_{ij}| = p_{1j} = |\alpha|$$

b) Possono calcolare il det $\overset{\text{def}}{A}$ con esempio

utilizzando le formule di Laplace, sviluppando rispetto alle prime colonne

$$\det A = 1 \cdot \det \begin{bmatrix} 1 & \alpha & \dots & \alpha \\ \alpha & 1 & \dots & \alpha \\ \vdots & \vdots & \ddots & \vdots \\ \alpha & \alpha & \dots & 1 \end{bmatrix} + (-1)^{n+1} \cdot \alpha \cdot \det \begin{bmatrix} 1 & \alpha & \dots & \alpha \\ \alpha & 1 & \dots & \alpha \\ \vdots & \vdots & \ddots & \vdots \\ \alpha & \alpha & \dots & 1 \end{bmatrix}$$

$$= 1 + (-1)^{n+1} \alpha \cdot \alpha (-1)^n \det (I_{n-2}) =$$

$$= 1 - 1 \cdot \alpha^2 = 0 \quad (\Leftrightarrow \alpha^2 = 1 \Leftrightarrow |\alpha| = 1)$$

$$c) \|A\|_1 = \max_j \sum_{i=1}^n |\alpha_{ij}| = 1 + |\alpha|$$

$$\|A\|_\infty = \max_i \sum_{j=1}^n |a_{ij}| = \max \{ (n-1)|\lambda| + 1, 1, 1 + |\lambda| \}$$

perché $n \geq 2$

$$\|A\|_\infty = (n-1) |\lambda| + 1 \geq 1 + |\lambda| = \|A\|_1$$

d) per $\lambda = 1$ $A^T A = I + ee^T + e_n e_n^T + e_n e_n^T$

$$A^T = (I + e_1 e_1^T + ue_n^T)$$

$$A^T A = (I + e_1 e_1^T + ue_n^T)(I + e_n e_n^T + e_n ue^T)$$

$$= I + e_1 e_1^T + ue_1^T + e_n e_n^T + \underbrace{e_1 e_1^T e_n e_n^T e_1^T}_1 + \\ + e_1 ue^T + \underbrace{e_n e_n^T e_1 u}_0 + \underbrace{ue_n^T e_1 ue^T}_1$$

$$= I + \underbrace{e_1 e_n^T}_x + \underbrace{ue_1^T}_x + e_n e_1^T + e_1 e_1^T + \underbrace{e_n ue^T + ue^T u}_1$$

$$= I + e_1 e_n^T + (u + e_1) e_1^T + (e_n + u) ue^T + e_n e_1^T$$

$$= I + e_1 e_n^T + e (e_n^T + u^T) + e_n e_1^T$$

$$= I + ee^T + e_1 e_n^T + e_n e_1^T$$

$$= \begin{bmatrix} 1 & & & \\ & \ddots & & \\ & & 1 & \\ & & & 1 \end{bmatrix} + \begin{bmatrix} 1 & \cdots & 2 \\ & \ddots & 1 \\ & & 1 \\ 2 & \cdots & -1 \end{bmatrix} = \begin{bmatrix} 2 & 1 & \cdots & 1 & 2 \\ 1 & \ddots & \ddots & 1 \\ \vdots & & \ddots & \vdots \\ 2 & 1 & \cdots & -1 & 2 \end{bmatrix}$$

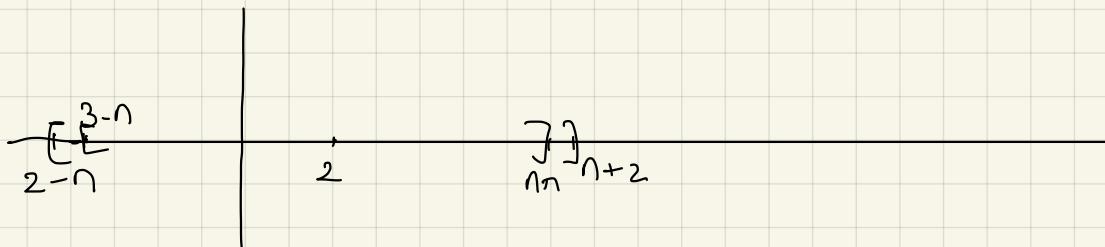
$$\|A\|_2 = \sqrt{\rho(A^T A)}$$

Quindi per dimostrare che $\|A\|_2 \leq \sqrt{n+2}$ occorre dimostrare che $\rho(A^T A) \leq n+2$

Possiamo guardare con i cerchi di Gershgorin se sono possibili concentri al coppia spettrale.

$$K_1 \equiv K_n \quad \text{centro 2 e coppia } (n-2)+2 = n$$

$$K_2 \equiv K_3 \equiv \dots \equiv K_{n-1} \quad \text{centro 2 e coppia } n-1$$



$$\Rightarrow \rho(A^T A) \leq \max\{n+2, |2-n|\} = n+2$$

$$\Rightarrow \|A\|_2 \leq \sqrt{n+2}$$

c) Fait LU

$$A = \begin{bmatrix} 1 & \alpha & -\alpha \\ & \ddots & \\ & & 1 \end{bmatrix}$$

$$L_{n-1} = \begin{bmatrix} I_{n-1} \end{bmatrix}$$

$$U_{n-1} = \begin{bmatrix} 1 & \alpha & -\alpha \\ & \ddots & \\ & & 1 \end{bmatrix}$$

$$\begin{pmatrix} d \\ 0 \\ \vdots \\ 0 \end{pmatrix} = U_{n-1} \cdot y \quad y = \begin{pmatrix} d \\ 0 \\ \vdots \\ 0 \end{pmatrix}$$

$$x^T U_{n-1} = (d \dots d)$$

$$U_{n-1}^T x = \begin{pmatrix} d \\ 0 \\ \vdots \\ 0 \end{pmatrix} \quad \begin{pmatrix} d \\ d \\ d \\ \vdots \\ d \end{pmatrix} \begin{pmatrix} x_1 \\ \vdots \\ x_{n-1} \end{pmatrix} = \begin{pmatrix} d \\ 0 \\ \vdots \\ 0 \end{pmatrix}$$

Matrice elementos
di gauss

$$(U_{n-1}^T)^{-1} = \begin{pmatrix} 1 & & & \\ -d & 1 & & \\ -d & -d & 1 & \\ -d & -d & -d & 1 \end{pmatrix}$$

$$\begin{pmatrix} x_1 \\ \vdots \\ x_{n-1} \end{pmatrix} = (U_{n-1}^T)^{-1} \begin{pmatrix} d \\ 0 \\ \vdots \\ 0 \end{pmatrix} = \begin{pmatrix} 1 & & & \\ -d & 1 & & \\ -d & -d & 1 & \\ -d & -d & -d & 1 \end{pmatrix} \begin{pmatrix} d \\ 0 \\ \vdots \\ 0 \end{pmatrix} = \begin{pmatrix} d \\ -d^2 \\ \vdots \\ -d^{n-2} \end{pmatrix}$$

$$x^T y + \beta = 1$$

$$\beta = 1 - x^T y = 1 - (d - d^2 - \dots - d^{n-2}) \begin{pmatrix} d \\ 0 \\ \vdots \\ 0 \end{pmatrix}$$

$$= 1 - d^{n-2}$$

$$L = \begin{pmatrix} 1 & & & \\ d-d^2 & 1 & & \\ d-d^2 & -d^2 & 1 & \\ \vdots & \vdots & \ddots & 1 \end{pmatrix}$$

$$U = \begin{pmatrix} 1 & & & \\ d & -d & & \\ d & d & -d & \\ \vdots & \vdots & \ddots & 1 \\ d & d & \dots & 1-d^{n-2} \end{pmatrix}$$

(F)

$$\begin{cases} L y = b \\ U x = y \end{cases}$$

$$y_{1:n-1} = b_{1:n-1}$$

$$d y_n - \sum_{i=2}^{n-1} d^2 y_i + y_n = 0$$

$$y_n = \sum_{i=2}^{n-1} \alpha^2 y_i - \alpha y_1$$

$$x_n = \frac{y_n}{1-\alpha^2}$$

$$x_n = y_n \quad i = n-1 : -1 : 2$$

$$x_1 + \alpha x_2 + \dots + \alpha^{n-1} x_n = y_n$$

$$x_1 = y_n - \sum_{i=2}^{n-1} \alpha x_i$$

function $x = \text{ssysolve}(b, \alpha f)$

$n = \text{length}(b);$

$y = b;$

$$y(n) = -\alpha f * y(1);$$

for $i = 2 : n-1$

$$y(n) = y(n) + \alpha f^2 * y(i)$$

end

$x = y;$

$$x(n) = y(n) / (1 - \alpha f^2);$$

for $i = 2 : n$

$$x(i) = x(1) - \alpha f * x(i);$$

end

end