

$$A = A^{(0)} = \begin{bmatrix} a_{11}^{(0)} & \dots & a_{1n}^{(0)} \\ \vdots & & \vdots \\ a_{m1}^{(0)} & \dots & a_{mn}^{(0)} \end{bmatrix}$$

If $a_{j1}^{(0)} \neq 0$ Pivots $\sigma: \left(a_{j1}^{(0)} \right) = \max_{k=1..n} \left(a_{k1}^{(0)} \right)$

Scambio righe j con riga 1

$$A^{(0)} \rightarrow A^{(1)} = \begin{bmatrix} a_{j1}^{(0)} & \dots & a_{jn}^{(0)} \\ a_{11}^{(0)} & \dots & a_{1n}^{(0)} \end{bmatrix}$$

$$A^{(1)} = P_1 A^{(0)}$$

P_1 matrice ottenuta dall'identità scambiando
colonna $\cdot 1$ e j

P_2 matrice di permutazione (proprietà: $P^{-1} = P^T$)

$$A^{(1)} = P_1 A^{(0)} \rightarrow A = P_2 P_1 A^{(0)}$$

$$A^{(1)} = \begin{bmatrix} a_{11}^{(1)} & \dots & a_{1n}^{(1)} \\ 0 & a_{22}^{(1)} & \dots \\ \vdots & \vdots & \vdots \\ 0 & a_{j2}^{(1)} & \dots \\ \vdots & \vdots & \vdots \end{bmatrix}$$

$$\exists j \geq 2 \text{ t.c. } a_{j2}^{(1)} \neq 0$$

L non è triviale

(“psicologicamente triviale” HoL & B)

$$[L, U] \subset \text{lin}(A) \quad / \quad X = A \setminus b$$

printing



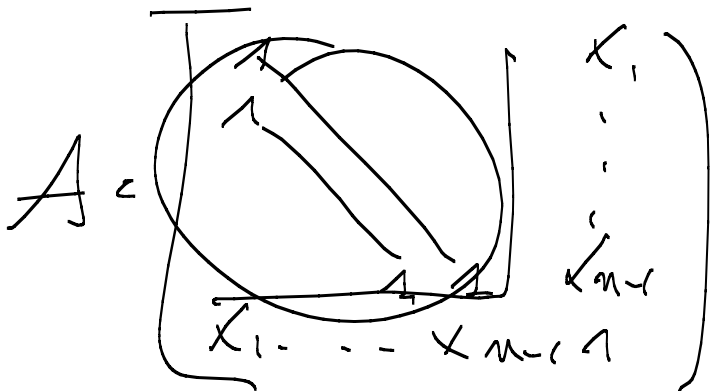
(partial printing)

① rende il processo applicabile
ad ogni istanza numerabile

② meglio B stabilizzato

③ con cura L e U computabili

④ più convenienti di $full-u$.



L, U sparse.

no row printing

U sparse.

$$A = \begin{bmatrix} \epsilon & 1-\epsilon \\ 1+\epsilon & -\epsilon \end{bmatrix}$$

$$b = \begin{bmatrix} 1 \\ 1 \end{bmatrix} \quad x = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

$$x_1 = A \setminus b \quad \text{OK}$$

$$U = \begin{bmatrix} 1 & 0 \\ -\frac{(1+\epsilon)}{\epsilon} & 1 \end{bmatrix}$$

$$U = I \cdot A$$

$$x_2 = U \setminus b$$

$$b = I \cdot b$$

NO

Metodi Iterativi

$Ax = b$ A invertibile A sparse

A sparse = molti elementi zero nulli

$$A = \begin{pmatrix} \times & \times & \times \\ \times & & \\ \times & & \end{pmatrix}$$

$$nnz(A) < cn^2$$

$$nnz(A) < \underbrace{O(n)}_p \leftarrow$$

$$O(n \log n)$$

$$O(n \sqrt{n})$$

Metodo di decomposizione gaussiana = metodo diretto
 = numero finito di passi determinati e soluzioni del sistema.

Metodo di Gauss-Jordan: costruiamo una successione
 $\{x^{(k)}\}_{k \in \mathbb{N}}$ di vettori tali che $x^{(k)} \rightarrow x$
 soluzioni del sistema lineare.

Critero di arresto: Quante m. dec. prec.

$$\{x^{(k)}\}_{k \in \mathbb{N}} \quad x^{(k)} \in \mathbb{R}^n$$

$$\text{Def: } \{x^{(k)}\}_{k \in \mathbb{N}} \quad \lim_{k \rightarrow +\infty} x^{(k)} = x$$

$$\Leftrightarrow \lim_{k \rightarrow +\infty} \|x^{(k)} - x\| = 0$$

Quali sono i v. m. \circledast (EQUIVALENZA TOPOLOGICA)

$$\|x^{(k)} - x\|_{\infty} = \max_{j=1, \dots, m} |x_j^{(k)} - x_j| \xrightarrow{k \rightarrow +\infty} 0$$

$$|x_j^{(k)} - x_j| \xrightarrow{k \rightarrow +\infty} 0 \quad j=1, \dots, m$$

$$0 \leq |x_j^{(k)} - x_j| \leq \max_{j=1, \dots, m} |x_j^{(k)} - x_j|$$

$$Ax = b \quad A \in \mathbb{R}^{m \times n}$$

M invertible.

$$Ax = b \Leftrightarrow (M - N)x = b$$

$$\Leftrightarrow Mx = Nx + b$$

$$\Leftrightarrow x = M^{-1}Nx + M^{-1}b$$

$$\Leftrightarrow x = Px + q$$

$$P = M^{-1}N$$

$$q = M^{-1}b$$

$$Ax = b$$

$$\Leftrightarrow$$

$$\begin{cases} x = Px + q \\ P = M^{-1}N & q = M^{-1}b \\ A = M - N \end{cases}$$

$$x = Px + q \rightsquigarrow \begin{cases} x^{(0)} \in \mathbb{R}^n \\ x^{(k+1)} = Px^{(k)} + q \end{cases}$$

Teorema: Se luo $x^{(k)} = x \in \mathbb{R}^n$

allora $x = Px + q$

(questo è il punto del sistema lineare)

Prüf: $x^{(k)} \rightarrow x$ (1 ptes.)

$$x^{(k+1)} = Px + q \quad (1 \text{ ptes.})$$

$$x = \lim_{k \rightarrow \infty} x^{(k+1)} = \lim_{k \rightarrow \infty} Px^{(k)} + q = Px + q$$

Frage: $A = \begin{bmatrix} 2 & 1 \\ 1 & 2 \end{bmatrix} \quad b = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$

$$Ax = b \quad x = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

① $M = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \quad N = \begin{bmatrix} -2 & 0 \\ 0 & -2 \end{bmatrix}$

$$A = M - N \quad M \text{ invertierbar}$$

$$P = M^{-1}N = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} -2 & 0 \\ 0 & -2 \end{bmatrix}$$

$$= \begin{bmatrix} 0 & -2 \\ -2 & 0 \end{bmatrix}$$

$$x^{(k+1)} = P x^{(k)} + \text{~~something~~}$$

$$\begin{bmatrix} x_1^{(k+1)} \\ x_2^{(k+1)} \end{bmatrix} = \begin{bmatrix} 0 & -2 \\ -2 & 0 \end{bmatrix} \begin{bmatrix} x_1^{(k)} \\ x_2^{(k)} \end{bmatrix}$$

Successione generata $\rightarrow \begin{bmatrix} 0 \\ 0 \end{bmatrix}$?

de parte dalla scelta del vettore iniziale.

$$\begin{bmatrix} x_1^{(0)} \\ x_2^{(0)} \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

$$\begin{bmatrix} x_1^{(k)} \\ x_2^{(k)} \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix} \quad \forall k$$

0 k konvergenz.

$$\begin{bmatrix} x_1^{(0)} \\ x_2^{(0)} \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \end{bmatrix} \quad \begin{bmatrix} x_1^{(1)} \\ x_2^{(1)} \end{bmatrix} = \begin{bmatrix} -2 \\ -2 \end{bmatrix}$$

$$\begin{bmatrix} x_1^{(2)} \\ x_2^{(2)} \end{bmatrix} = \begin{bmatrix} 4 \\ 4 \end{bmatrix} \quad \begin{bmatrix} x_1^{(3)} \\ x_2^{(3)} \end{bmatrix} = \begin{bmatrix} -8 \\ -8 \end{bmatrix}$$

$$\begin{bmatrix} x_1^{(k)} \\ x_2^{(k)} \end{bmatrix} = (-1)^{(k)} \begin{bmatrix} 2^k \\ 2^k \end{bmatrix}$$

$x^{(k)}$ diverge (non converg)

$$A = \begin{bmatrix} 2 & 1 \\ 1 & 2 \end{bmatrix} \quad H = \begin{bmatrix} 2 & 0 \\ 0 & 2 \end{bmatrix} \quad S = \begin{bmatrix} 0 & -1 \\ -1 & 2 \end{bmatrix}$$

$$P = H^{-1} N_2 \begin{pmatrix} \frac{1}{2} & 0 \\ 0 & \frac{1}{2} \end{pmatrix} \begin{pmatrix} 0 & -1 \\ -1 & 0 \end{pmatrix}$$

$$= \begin{pmatrix} 0 & -\frac{1}{2} \\ -\frac{1}{2} & 0 \end{pmatrix}$$

$$x^{(k+1)} = P x^{(k)}$$

$$x^{(0)} = \begin{pmatrix} 1 \\ 1 \end{pmatrix} \quad x^{(1)} = \begin{pmatrix} -\frac{1}{2} \\ -\frac{1}{2} \end{pmatrix} \quad x^{(2)} = \begin{pmatrix} \frac{1}{4} \\ \frac{1}{4} \end{pmatrix}$$

$$x^{(k)} \rightarrow 0 \text{ Steady state.}$$

Convergence in general depends on

zero - row - elements M e N

$$(A = M - N)$$

lezione 24/04

$$Ax = b \quad A \in \mathbb{R}^{M \times N} \quad M \text{ invertibile.}$$

$$(M-N) \times b \quad (\Leftrightarrow) \quad Mx = Nx + b$$

$$(\Leftrightarrow) \quad x = \underbrace{M^{-1}N}_P x + \underbrace{M^{-1}b}_q$$

$$\left\{ \begin{array}{l} x^{(0)} \in \mathbb{R}^n \\ x^{(k+1)} = Px^{(k)} + q \end{array} \right.$$

$$\left\{ \begin{array}{l} x^{(0)} \in \mathbb{R}^n \\ Mx^{(k+1)} = Nx^{(k)} + b \end{array} \right.$$

Fissata A la convergenza dipende in grado
dalla selezione di $\varepsilon \in \mathbb{R}, N$ (metodo)
o. vettore iniziale.

Usi della convergenza

Convergenza

Def: Un metodo iterativo $x^{(k+1)} = P x^{(k)} + q$ ($k \geq 1$)
per risolvere $Ax = b$ con $P = M^{-1}N$ $q = M^{-1}b$
 $A \in \mathbb{R}^{n \times n}$, si dice CONVERGENTE se

$x^{(0)} \in \mathbb{R}^m$
 $x^{(k)} \in \mathbb{R}^m$ $\lim_{k \rightarrow \infty} x^{(k)} = x$ soluzione del sistema
cioè la successione genera convergenza alla
soluzione del sistema lineare

$$x^{(k+1)} = P x^{(k)} + q \quad \text{metodo iterativo}$$

$$x = P x + q \quad \text{soluzione del sistema}$$

$$x^{(k+1)} - x = P (x^{(k)} - x)$$

$$e^{(k+1)} = P e^{(k)} - x$$

$$\boxed{e^{(k+1)} = P e^{(k)} \quad k \geq 0}$$

Lemma: IP reitete $x^{(k+1)} = P x^{(k)} + q$ e

Convergente $\Leftrightarrow \|P\| < 1$ wenn wir die Norm $\|\cdot\|$ wählen können, so dass $\|P\| < 1$.

$$\begin{aligned} \text{Dann: } e^{(k+1)} = P e^{(k)} &\Rightarrow \|e^{(k+1)}\| = \|P e^{(k)}\| \\ &\leq \|P\| \|e^{(k)}\| \leq \|P\|^2 \|e^{(k-1)}\| \leq \dots \\ &\leq \|P\|^{k+1} \|e^{(0)}\| \quad \underline{0 \leq \alpha < 1} \end{aligned}$$

$$0 \leq \|e^{(k+1)}\| \leq \|P\|^{k+1} \|e^{(0)}\|$$

$\downarrow 0$ $\downarrow 0$ $\downarrow 0$

Lemma also annehmen.

~~QED~~

Teorema: Se il vettore $x^{(k+1)} = P x^{(k)} + q$ è
 convergente allora $\rho(P) < 1$

Dim: Se \bar{x} è convergente la successione converge
 $\forall x^{(0)}$ e quindi $e^{(k)} \rightarrow 0 \quad \forall e^{(0)}$

Prendiamo $e^{(0)} = v$ di $Pv = \lambda v$ con
 $|\lambda| = \rho(P)$

$$e^{(k+1)} = P e^{(k)} = \dots = P^{k+1} e^{(0)} = P^{k+1} v =$$

$$\lambda^{k+1} v = \lambda^{k+1} e^{(0)}$$

$$\Rightarrow \|e^{(k+1)}\| = |\lambda|^{k+1} \|e^{(0)}\| \neq 0$$

$$\|e^{(k+1)}\| \rightarrow 0 \Leftrightarrow |\lambda| < 1$$



Teorema: Se $\rho(P) < 1$ allora il vettore
 $x^{(k+1)} = P x^{(k)} + q$ è convergente.

dim: (slo x vettori diagonalizzabili)

$$P = V \cdot D \cdot V^{-1}$$

$$P^{k+1} = V \cdot D^{k+1} \cdot V^{-1}$$

$$\|P^{k+1}\|_{\infty} \leq \|V\|_{\infty} \|V^{-1}\|_{\infty} \rho(P)^{k+1}$$

$$\begin{aligned} 0 < \epsilon &\Leftrightarrow 0 < \|e^{(k+1)}\|_{\infty} \leq \|P^{k+1}\|_{\infty} \|e^{(0)}\|_{\infty} \\ &\leq \|V\|_{\infty} \|V^{-1}\|_{\infty} \rho(P)^{k+1} \|e^{(0)}\|_{\infty} \\ &\downarrow 0 \end{aligned}$$



Teorema: Condizione necessaria e sufficiente x to

Convergenze e ok $\rho(P) < 1$.

$$X^{(k+1)} = P X^{(k)} + q \quad P = \begin{bmatrix} \frac{1}{2} & \frac{2}{3} \\ \frac{1}{3} & \frac{1}{4} \end{bmatrix}$$

Convergente?

$$\frac{11}{12} = \|P\|_1 = \max \left\{ \frac{1}{2} + \frac{1}{3}, \frac{2}{3} + \frac{1}{4} \right\} = \max \left\{ \frac{5}{6}, \frac{11}{12} \right\}$$

$$\|P\|_\infty = \frac{1}{2} + \frac{2}{3} = \frac{3+4}{6} = \frac{7}{6} > 1$$

$$\|P\|_\infty > 1 \quad \text{e} \quad \|P\|_1 < 1$$

Il vettore è convergente

$$A = \begin{bmatrix} 1 & & & 0 \\ & \ddots & & \\ & & 1 & \\ 0 & & & 1 \end{bmatrix}$$

$$A = I - N$$

$$M = \begin{bmatrix} 1 & & & \\ & \ddots & & \\ & & 1 & \\ 0 & & & 1 \end{bmatrix}$$

$$N = \begin{bmatrix} 0 & & & \\ & \ddots & & \\ & & 0 & \\ & & & 0 \end{bmatrix}$$

è convergente? A riduzion: le proprietà di \underline{P}

$$P = M^{-1}N = M^{-1} \begin{bmatrix} 0 & | & 0 & | & \dots & | & 0 & | & v \end{bmatrix}$$

$$= \begin{bmatrix} 0 & | & 0 & | & \dots & | & 0 & | & M^{-1}v \end{bmatrix}$$

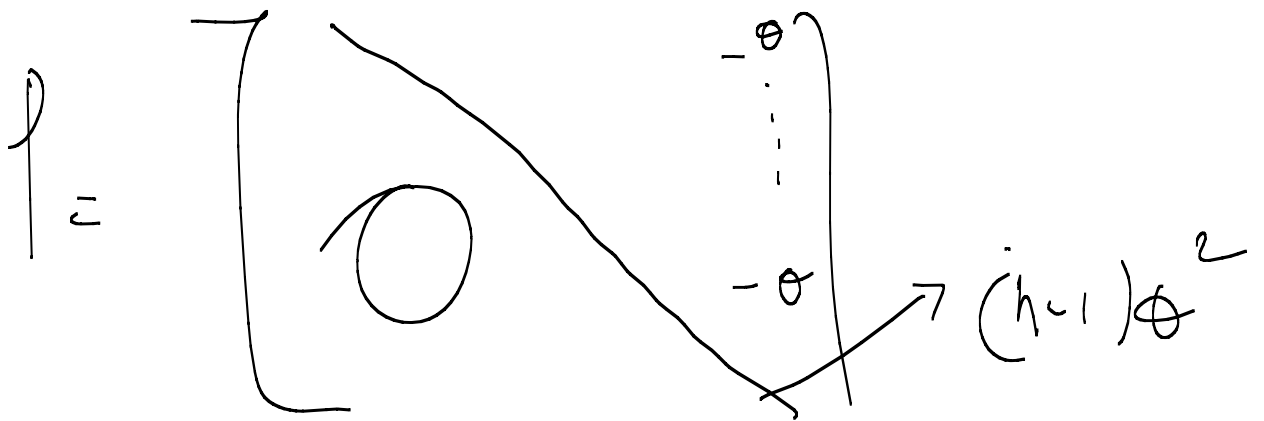
$$M^{-1}v = x \Leftrightarrow Mx = v$$

$$\begin{bmatrix} 1 & & & \\ & \ddots & & \\ & & 1 & \\ 0 & & & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ \vdots \\ x_m \end{bmatrix} = \begin{bmatrix} 0 \\ \vdots \\ 0 \\ v \end{bmatrix}$$

$$x_1 = x_2 = \dots = x_{n-1} = -\theta$$

$$\theta x_1 + \theta x_2 + \dots + \theta x_{n-1} + x_n = 0$$

$$x_n = \theta^2 + \theta^2 + \dots + \theta^2$$



$\varphi(P) = ?$, P è trapezoido superiore e piano.

1) Sui vertici, si usa il logaritmo principale

$$K = 0 \quad K = (n-1)\theta^2$$

$$\varphi(P) \leq (n-1)\theta^2$$

$$\text{è congettura } (\Leftrightarrow) \quad (n-1)\theta^2 < 1$$

$$\Leftrightarrow \quad \theta^2 < \frac{1}{n-1}$$

$$\Leftrightarrow -\frac{1}{\sqrt{n-1}} < \theta < \frac{1}{\sqrt{n-1}}$$

$$\|P\|_{\infty} = \max \left\{ |\theta|, (n-1)\theta^2 \right\} < 1$$

$$A = \begin{bmatrix} 1 & & & & \theta \\ & 1 & & & \theta \\ & & \ddots & & \theta \\ & & & 1 & \theta \\ \theta & \theta & & \theta & 1 \end{bmatrix} \quad M = \begin{bmatrix} 1 & & & & \\ & 1 & & & \\ & & \ddots & & \\ & & & 1 & \\ & & & & 1 \end{bmatrix}$$

$$N = \begin{bmatrix} & & & -\theta & \\ & 0 & & \vdots & \\ & & & \theta & \\ -\theta & \dots & \theta & & 0 \end{bmatrix}$$

$$M^{-1}N = P = \begin{bmatrix} 0 & & & -\theta & \\ & & & \vdots & \\ & & & \theta & \\ \theta & \dots & \theta & & 0 \end{bmatrix}$$

Abkürzung: mit zu θ hinreichend klein oder negativ

$$\|P\|_{\infty} = \|P\|_1 < 1$$

$$(n-1) |\theta| < 1$$

$$|\theta| < \frac{1}{n-1}$$

$$-\frac{1}{n-1} < \theta < \frac{1}{n-1}$$

Consequenter muss es möglich sein $\varphi(P)$?

$$\det \begin{pmatrix} \begin{matrix} x & & & 0 \\ & \ddots & & \vdots \\ & & & \theta \\ \theta & \dots & \theta & x \end{matrix} \end{pmatrix} = ?$$

$\Delta \neq 0$ bedeutet auch die Möglichkeit LU

$$\begin{bmatrix} x & & & 0 \\ & \ddots & & \vdots \\ & & & \theta \\ \theta & \dots & \theta & x \end{bmatrix} = \begin{bmatrix} 1 & & & \\ & \ddots & & \\ & & 1 & \\ \theta/x & \dots & \theta/x & 1 \end{bmatrix} \begin{bmatrix} x & & & 0 \\ & \ddots & & \vdots \\ & & & \theta \\ & & & x \end{bmatrix}$$

$$x + (n-1) \frac{\theta^2}{x} = x \quad x = x - (n-1) \frac{\theta^2}{x}$$

$$\det(\Delta J - P) = \det U \cdot \det L$$

$$= x^{n-1} \cdot \left[x - \frac{(n-1)\theta^2}{x} \right]$$

$$= x^n - (n-1)\theta^2 x^{n-2}$$

$$= x^{n-2} \cdot \left(x^2 - (n-1)\theta^2 \right) = 0$$

$x = 0$
 $x = \pm \sqrt{(n-1)\theta^2}$

$$\varphi.(P)_2 \quad \sqrt{n-1} |\theta| < 1$$

$$\Leftrightarrow |\theta| < \frac{1}{\sqrt{n-1}}$$